

ctf-实验吧-crypto

原创

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实验吧

2 篇文章 0 订阅

订阅专栏



crypto

20 篇文章 1 订阅

订阅专栏

1.trivial

下载文件解压, 得到一个py文件, 里面又一个加密算法, 这是一个类移位加密, 加密算法的key和加密之后的密文已知, 然后对每个明文在不同范围内进行移位加密, 直接解密即可

```
#!/usr/bin/env python
import sys

alphaL = "abcdefghijklmnopqrstuvwxyz"
alphaU = "ABCDEFGHIJKLMNOPQRSTUVWXYZ"
num = "0123456789"
keychars = num+alphaL+alphaU
key="T0pS3cre7key"
plaintext="Bot kmws mikferuigmzf rmfrrwqe abs perudsfl Nvm kda ut ab8bv_w4ue0_ab8v_DDU"
ciphertext=""
for i in range(len(plaintext)):
    rotate_amount = keychars.index(key[i%len(key)])
    if plaintext[i] in alphaL:
        enc_char = ord('a') + (ord(plaintext[i])-ord('a')-rotate_amount)%26
    elif plaintext[i] in alphaU:
        enc_char = ord('A') + (ord(plaintext[i])-ord('A')-rotate_amount)%26
    elif plaintext[i] in num:
        enc_char = ord('0') + (ord(plaintext[i])-ord('0')-rotate_amount)%10
    else:
        enc_char = ord(plaintext[i])
    ciphertext = ciphertext + chr(enc_char)
print ciphertext
```

2.rsarsa

rsa加密, 已知p, q, e, c, 算出n, d, 直接解密, 提交十进制的明文即可

```

import gmpy2
p=964842302901051567659055174001042653494573763923573980064398935203985250729849139956103500916342705037010
q=118748438379802970320924058486536568527609101545433809076500401907042833589092085782510630477324439922306
e=65537
c=832082989951746041747735902982036393605400248712561268928896613457424033149298619391004926666056473166465
n=p*q
d=int(gmpy2.invert(e,(p-1)*(q-1)))
mingwenint=pow(c,d,n)
print mingwenint

```

3.Marvin is plain Jane

根据题目提示这是Menezes-Vanstone加密（题中提到他的朋友叫menezes或van-stone），这里简单介绍一下Menezes-Vanstone加密，有利于理解下面的内容，Menezes-Vanstone加密是基于椭圆曲线的加密方案，在椭圆曲线上取一个基点 $G(x,y)$ ， A ， B 选择随机数 da ， db 作为密钥，然后就可以得到公钥 $DA=da*G$ ， $DB=db*G$ 。 A 将明文 M 转化为二元组 $m(m_1,m_2)$ ，然后计算 $(x_1,y_1)=da*Db$ ，最后计算 $x_2=(m_1*x_1)modp$ ， $y_2=(m_2*y_1)modp$ 组成二元组 (x_2,y_2) 即为密文，而 B 拿到密文后计算 $(x_1,y_1)=db*Da$ ，则可以得到 $m_1=(x_2*x_1^{-1})modp$ ， $m_2=(y_2*y_1^{-1})modp$ ，将 m_1 和 m_2 组合即可。在这道题中提到的brainpool p256r1经查询是椭圆曲线中的一个标准(在rfc5639中)，可查到该椭圆曲线的一些标准，而在这里我们已知了二元组 m 中 m_1 （文中反复提到的"Marvin is"），以及 m 加密后的 (x_2,y_2) ，而 (x_1,y_1) 所在的椭圆曲线已知(brainpool p256r1)，那么首先计算 $x_1=x_2*m_1^{-1}$ ，根据椭圆曲线得到 x_1 对应点 y_1 ，在计算 $m_2=y_2*y_1^{-1}$ 即可得到答案。

```

def txt(istr):
    return int(istr.encode("hex"),16)
x1 = txt("Marvin is")
y1 = 71164450240897430648972143714791734771985061339722673162401654668605658194656
y2 = 12951693517100633909800921421096074083332346613461419370069191654560064909824
p = 0xA9FB57DBA1EEA9BC3E660A909D838D726E3BF623D52620282013481D1F6E5377
A = 0x7D5A0975FC2C3057EEF67530417AFFE7FB8055C126DC5C6CE94A4B44F330B5D9
B = 0x26DC5C6CE94A4B44F330B5D9BBD77CBF958416295CF7E1CE6BCCDC18FF8C07B6
#这里使用Tonelli-shanks算法寻找整数x使得a=(x^2)modp成立
def modular_sqrt(a, p):
    if legendre_symbol(a, p) != 1:
        return 0
    elif a == 0:
        return 0
    elif p == 2:
        return p
    elif p % 4 == 3:
        return pow(a, (p + 1) / 4, p)
    s = p - 1
    e = 0
    while s % 2 == 0:
        s /= 2
        e += 1
    n = 2
    while legendre_symbol(n, p) != -1:
        n += 1
    x = pow(a, (s + 1) / 2, p)
    b = pow(a, s, p)
    g = pow(n, s, p)
    r = e
    while True:
        t = b
        m = 0
        for m in xrange(r):

```

```

    if t == 1:
        break
    t = pow(t, 2, p)

    if m == 0:
        return x

    gs = pow(g, 2 ** (r - m - 1), p)
    g = (gs * gs) % p
    x = (x * gs) % p
    b = (b * g) % p
    r = m
#这里使用Euler's criterion判别是否存在整数x使得a=(x^2)modp成立
def legendre_symbol(a, p):
    ls = pow(a, (p - 1) / 2, p)
    return -1 if ls == p - 1 else ls
def egcd(a, b):
    if a == 0:
        return (b, 0, 1)
    else:
        g, y, x = egcd(b % a, a)
        return (g, x - (b // a) * y, y)
def modinv(a, m):
    g, x, y = egcd(a, m)
    if g != 1:
        raise Exception('modular inverse does not exist')
    else:
        return x % m
def gety(x):
    y = modular_sqrt(x**3 + A * x + B, p) % p
    y2 = -y % p
    return y,y2
def hextotext(nbr):
    s = hex(nbr)[2:-1]
    if len(s) % 2 ==1:
        s = "0"+s
    return s.decode("hex")
x1_inv = modinv(x1, p)
c1 = (y1 * x1_inv) % p
c2_1, c2_2 = gety(c1)
print repr(hextotext(y2*modinv(c2_1, p) % p))
print repr(hextotext(y2*modinv(c2_2, p) % p))

```

5.RSAROLL

根据题意可知该题是使用了RSA加密，而在文件中给出了n，e和一些c，这里n比较小，很容易就可以分解成p，q，然后即可解出d，再对c解密发现他们都在128以内，使用ascii表转化组合即可得到答案


```

'I':"CDCCC",
'J':"CDCCD",
'K':"CDCDC",
'L':"CDCDD",
'M':"CDDCC",
'N':"CDDCD",
'O':"CDDDC",
'P':"CDDDD",
'Q':"DCCCC",
'R':"DCCCD",
'S':"DCCDC",
'T':"DCCDD",
'U':"DCDCC",
'V':"DCDCD",
'W':"DCDDC",
'X':"DCDDD",
'Y':"DDCCC",
'Z':"DDCCD",
}
def decode(s):
    cipher=""
    i=0
    while i<len(s):
        c=s[i:i+5]
        sign=True
        for k in codebook3.keys():
            if codebook3[k] == c:
                cipher+=k
                sign=False
                break
        if sign:
            #cipher+=c
            pass
        i=i+5
    return(cipher)
cipher=""
ss = s.split(" ")
for c in ss:
    for k in codebook.keys():
        if codebook[k] == c:
            cipher+=k
            break
    else:
        cipher+=c
ch=cipher[34:]
list2=ch.split("/")
mingwen=""
for i in list2:
    mingwen+=decode(i)
    mingwen+=" "
print mingwen

```

打开文件发现顶部写着n, e, c, 说明这是rsa加密, 而再仔细看这里的所有n相同, 猜测这里是相同明文在不同公钥加密而来的结果, 而n有相同, 所以这里可以使用rsa共模攻击, 只需找到两个互质的e即可

```
import gmpy2
def egcd(a, b):
    if a == 0:
        return (b, 0, 1)
    else:
        g, y, x = egcd(b % a, a)
        return (g, x - (b // a) * y, y)

def shuchu(mingwenstr):
    if mingwenstr[len(mingwenstr)-1]=='L':
        mingwenstr=mingwenstr[2:len(mingwenstr)-1]
    else:
        mingwenstr=mingwenstr[2:len(mingwenstr)]
    if not len(mingwenstr)%2==0:
        mingwenstr='0'+mingwenstr
    i=len(mingwenstr)
    mingwen=""
    while i>=1:
        str1=mingwenstr[i-2:i]
        if int(str1,16)>33 and int(str1,16)<126:
            mingwen=chr(int(str1,16))+mingwen
        else :
            mingwen=" "+mingwen
        i=i-2
    print mingwen

list1=[{'e':0x1614984a0df,'c':0x7ded578992900e4d7799f910fdb615824d04b055336de784e88ba2d119f0c708c3b21e9d5}
{'e':0x15ef25e10f54a3,'c':0x7c5b756b500801e3ad68bd4f2d4e1a3ff94d049774bc9c37a05d4c18d212c5b223545444e7015a7}
{'e':0x1da0ca25f5a8d,'c':0x65af8559c93c05efecb6a3029dce7e831787878d5539f7b20fc7645ef4892cee23f53384377180a8}
{'e':0xc2eac4c2b,'c':0x711892a29a738e3ac3b996427e4188f23d1c63d9d9c962b6f65b675698e432f27f0ce4e42101576dacaf}
{'e':0x1a6c23,'c':0x57ab8d4c79a58718c0db0dd62a8ba97883e03cd7d14cc3366108a37e8998fc55abd555ca54f81fc975c64e1}
{'e':0x2beccafd,'c':0x6630d2faf104547351da66f760fde920203a041b82f07c0db9034148f9dd17c1f14c2c8ec95ae64e8d0b5}
{'e':0x280554063943,'c':0x5bcc5f5435fd087c9615bf04864a82a8fed19576fd311cfde565ca340303cd72d3842ad7a8de9c712}
{'e':0x23b0d,'c':0x48bd06fba0691da8883286c21cd49e02eb65d0e3b6ee12b2113940cc64d9f6b921fcb6a8aa82aac592e6a955}
{'e':0x6b8a5ae7,'c':0x6fdbcfb5cd2cacd032ef7200fd49b9f304a6bdb8399f4a91a72d1d9150f97b3b513f44dfc56f6f7c8ec41}
{'e':0x360f1c91fed,'c':0xa149bb3969479d5b9efff15099ce863d36899d1146c731a91db91ef15869358df4dbe82eaca128d5cd9}
{'e':0xefe30ec7dabb,'c':0x6fce7de911abe59864e01b9b306c167bdde17da28dfcfb7b3c768ec47d0ae4160cebedae9e482468c}
{'e':0x753fdb5,'c':0xa2403b99c19c2a882bbca51ee414486e1d60db003d16fdf8f30290bae586aaa5500c74b6e8dfe7a3081092}
{'e':0x12546aff963f4b6b35fL,'c':0x7293cec4d46c8073ce78b4ec8d97c086124376cf75fbcc4c0a57159b02e9c7a8545d4fa73}
{'e':0x11d2843e693,'c':0x9d96bc542c7afe105b6415d7f0f6d5114d81761f46bfaebbf36188f9fbc3759c4693645f4605d17}
{'e':0x9d540226f,'c':0x4c17528a0d1e36030f882d9c1060ccf974e48178cb7c4c8630968846ca668773881e41a780ba686315cc}
{'e':0xee4c39df4ed4c0f,'c':0x9fdf693d41e020a9efff786f87f12dc2e2b518ecaa178c8991d06a3ef2e8e136aba94441bc8dbd5}
{'e':0x213901ef4052b8b251c3L,'c':0xefa14d35d75b629673e8d983f1253134e4584ef16fd13618b23ea4e281f775942d370b38}
{'e':0xe93f,'c':0x2e1622191a5d6092ef23dabb82bdf0f5f9eb018f27184c05512679a38be06fc23ca57c1bd4129720e5d562ff}
{'e':0x4042c3955,'c':0x8caeeaa7d272f9606fee9222efd1d922143db738b95bd64746b27bc4c0fd979a2c57b4735131a4391a81b}
{'e':0x61553816b407935,'c':0x9eade5cb88d453b00c0558f76ab78dc76537588ed1212ffdfdc4ecff98c55457a4b581d1579011}
n=0xa5f7f8aaa82921f70aad9ece4eb77b62112f51ac2be75910b3137a28d22d7ef3be3d734dabb9d853221f1a17b1afb956a50236a}
gcdlist=[]
for i in xrange(0,len(list1)):
    j=i+1
    while j<len(list1):
        g,x,y=egcd(list1[i]['e'],list1[j]['e'])
        if g==1:
            gcdlist.append(x)
            gcdlist.append(y)
            c1 = list1[i]['c']
            c2 = list1[j]['c']
```

```

        break
    else:
        j=j+1

if gcdlist[0]<0:
    c1=int(gmpy2.invert(c1,n))
    gcdlist[0]=-gcdlist[0]
if gcdlist[1]<0:
    c2=int(gmpy2.invert(c2,n))
    gcdlist[1]=-gcdlist[1]
miwenint=(pow(c1,gcdlist[0],n)*pow(c2,gcdlist[1],n))%n
shuchu(hex(miwenint))

```

9.唯密文攻击

根据后面作者提示可知，在200个密文中有一个明文与n有公因子，找到这个明文，求出其与n的公因子，即可找到n的素因子中最小的那个，求其md5值按要求填入答案即可

```

import hashlib
def gcd(a,b):
    if b>a:
        temp=a
        a=b
        b=temp
    c=1
    while not c==0:
        c=a%b
        a=b
        b=c
    return a
list1=[]
file = open("ciphertext.txt")
line = file.readline()
n=file.readline()
n=int(n)
while 1:
    line = file.readline()
    if not line:
        break
    elif line[0]>='0' and line[0]<='9':
        list1.append(line)
file.close()
for i in list1:
    b=int(i)
    if not gcd(n,b)==1:
        g=gcd(n,b)
        break
ga=n/g
if ga>g:
    print (hashlib.md5(str(g)).hexdigest()).upper()[:6]
else:
    print (hashlib.md5(str(ga)).hexdigest()).upper()[:6]

```

10.RSA

下载文件打开，发现开头有“pk”，猜测只是zip压缩文件(可用binwalk检测)，添加后缀后解压，发现enc文件和pem文件，pem文件打开时openssl的公钥文件，而enc不是数据包文件，所以enc文件应该是公钥文件加密后的密文文件，使用openssl得到RSA的n, e, 用yafu对n分解获得p, q, 从而得到d, 用密钥(n,d)对密文解密即可

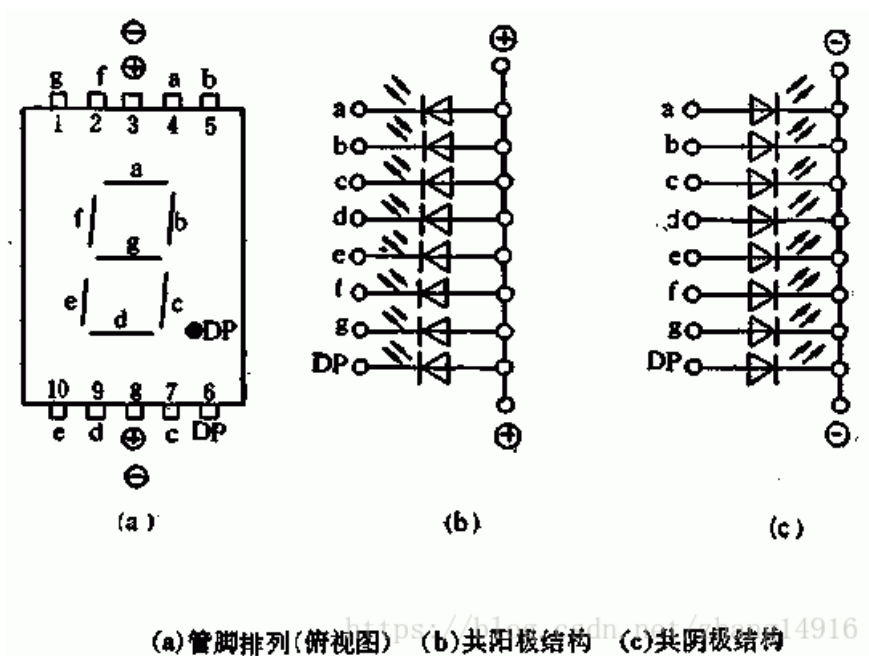
```
import sys
import os
import gmpy2
def shuchu(mingwenstr):
    if mingwenstr[len(mingwenstr)-1]=='L':
        mingwenstr=mingwenstr[2:len(mingwenstr)-1]
    else:
        mingwenstr=mingwenstr[2:len(mingwenstr)]
    if not len(mingwenstr)%2==0:
        mingwenstr='0'+mingwenstr
    i=len(mingwenstr)
    mingwen=""
    while i>=1:
        str1=mingwenstr[i-2:i]
        if int(str1,16)>33 and int(str1,16)<126:
            mingwen=chr(int(str1,16))+mingwen
        else :
            mingwen=" "+mingwen
        i=i-2
    print mingwen
p=258631601377848992211685134376492365269
q=286924040788547268861394901519826758027
e=65537
n=p*q
d=int(gmpy2.invert(e,(p-1)*(q-1)))
with open("flag.enc" , "rb") as f:
    s=f.read()
    miwen=long(s.encode('hex'),16)
    mingwenint=pow(miwen,d,n)
    mingwenstr=hex(mingwenint)
    shuchu(mingwenstr)
```

11.NSCTF crypto200

这个题准确的说应该是一道隐写题，下载得到一张图片，放入stegsolve中，发现在blue plane 0时这是个二维码，对其反色扫描即可得到答案

12.数码管

打开文件发现一张数码管的图片，查询有关数码管的资料，得知一个数码管需要8个电位(即0, 1)控制，而且有共阴显示和共阳显示两种（即发光对应高电位还是低电位），这里的8个电位可组成一个ascii值，而颜色的不同是表示其是否发光所对电位不同，尝试求解即可（从dp到a为从二进制高位到低位的方向）



13. 压缩的问题

打开链接，发现一串16进制字符串，看开头就知道这大概是rar文件，用winhex将16进制写入文件，生成rar文件，解压，发现需要密码，根据题目提示取出65h-71h字符输入，解压成功，获得txt文件，计算文件sha1值输入即可(可直接使用linux下命令)

```
from hashlib import sha1
sha1obj = sha1()
with open("秒破测试By天易love.txt", 'rb') as f:
    sha1obj.update(f.read())
print sha1obj.hexdigest()[:8]
```



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