

# buuctf easyRE

原创

菜逼的ctf之路 于 2020-10-18 10:37:22 发布 592 收藏

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本文链接：[https://blog.csdn.net/weixin\\_45701079/article/details/109141962](https://blog.csdn.net/weixin_45701079/article/details/109141962)

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## buuctf easyRE[2019红帽杯]

简单总结一下，一般说是easy的都不简单

加载ida查看字符串，找到主要函数，

```
.97    v52 = 76;
.98    v53 = 64;
.99    v54 = 69;
.00    v55 = 67;
.01    memset(v56, 0, sizeof(v56));
.02    v57 = 0;
.03    v58 = 0;
.04    v3 = v56;
.05    sub_4406E0(0LL, (__int64)v56);
.06    v58 = 0;
.07    v4 = v56;
.08    LODWORD(v5) = sub_424BA0((const __m128i *)v56);
.09    if ( v5 == 36 )
.10    {
.11        for ( i = 0; ; ++i )
.12        {
.13            v4 = v56;
.14            LODWORD(v7) = sub_424BA0((const __m128i *)v56);
.15            if ( i >= v7 )
.16                break;
.17            if ( (unsigned __int8)(v56[i] ^ i) != *(&v20 + i) )
.18            {
.19                result = -2;
.20                goto LABEL_13;
.21            }
.22        }
.23        sub_410CC0((const __m128i *)"continue!");
.24        memset(&v59, 0, 0x40uLL);
.25        v61 = 0;
.26        v3 = &v59;
.27        sub_4406E0(0LL, (__int64)&v59);
.28        v60 = 0;
.29        v4 = &v59;
.30        LODWORD(v8) = sub_424BA0((const __m128i *)&v59);
.31        if ( v8 == 39 )
```

00000AC1 main:97 (400AC1) |

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```

LODWORD(v8) = sub_424BA0((const __m128i *)&v59);
if ( v8 == 39 )
{
    v9 = (const __m128i *)sub_400E44((const __m128i *)&v59);
    v10 = (const __m128i *)sub_400E44(v9);
    v11 = (const __m128i *)sub_400E44(v10);
    v12 = (const __m128i *)sub_400E44(v11);
    v13 = (const __m128i *)sub_400E44(v12);
    v14 = (const __m128i *)sub_400E44(v13);
    v15 = (const __m128i *)sub_400E44(v14);
    v16 = (const __m128i *)sub_400E44(v15);
    v17 = (const __m128i *)sub_400E44(v16);
    v18 = sub_400E44(v17);
    v3 = off_6CC090;
    v4 = (char *)v18;
    if ( !(unsigned int)sub_400360(v18, (_int64)off_6CC090) )
    {
        sub_410CC0((const __m128i *)"You found me!!!");
        v4 = "bye bye~";
        sub_410CC0((const __m128i *)"bye bye~");
    }
    result = 0;
}
else
{
    result = -3;
}

```

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第一组数据解出

Info:The first four chars are flag

然后看下面，打开函数发现是base64加密，加密十次，然后比较，把字符串解密十次，发现给了一个网址(不好意思，网址给关了，不想再解密10次了，就口头描述一下)，打开网址，发现是...

然后看大佬博客，才知道是个迷惑我们的函数，真正的核心函数在上图所示函数的下一个(wdnmd)口吐芬芳

```

__int64 v2; // rdi
__int64 result; // rax
unsigned __int64 v4; // rt1
unsigned int v5; // [rsp+Ch] [rbp-24h]
signed int i; // [rsp+10h] [rbp-20h]
signed int j; // [rsp+14h] [rbp-1Ch]
unsigned int v8; // [rsp+24h] [rbp-Ch]
unsigned __int64 v9; // [rsp+28h] [rbp-8h]

v9 = __readfsqword(0x28u);
v2 = 0LL;
v5 = sub_43FD20(0LL) - qword_6CEE38;
for ( i = 0; i <= 1233; ++i )
{
    v2 = v5;
    sub_40F790(v5);
    sub_40FE60();
    sub_40FE60();
    v5 = (unsigned __int64)sub_40FE60() ^ 0x98765432;
}
v8 = v5;
if ( ((unsigned __int8)v5 ^ byte_6CC0A0[0]) == 102 && (HIBYTE(v8) ^ (unsigned __int8)byte_6CC0A3) == 103 )
{
    for ( j = 0; j <= 24; ++j )
    {
        v2 = (unsigned __int8)(byte_6CC0A0[j] ^ *((_BYTE *)&v8 + j % 4));
        sub_410E90(v2);
    }
}
v4 = __readfsqword(0x28u);
result = v4 ^ v9;
if ( v4 != v9 )
    sub_444020(v2, a2);
return result;
}

```

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然后利用"flag"字符串把v5解出来，然后把给定字符串依次异或就行了

最后贴上脚本

```
s='flag'
a=[
    0x40, 0x35, 0x20, 0x56, 0x5D, 0x18, 0x22,
    0x45, 0x17, 0x2F, 0x24, 0x6E, 0x62, 0x3C,
    0x27, 0x54, 0x48, 0x6C, 0x24, 0x6E, 0x72,
    0x3C, 0x32, 0x45, 0x5B, 0x00
]
key=''
flag=''
for i in range(4):
    key+=chr(ord(s[i])^a[i])

for i in range(len(a)):
    flag+=chr(ord(key[i%4])^a[i])
print(flag)
```

最后的结果各位自己来吧！我就不写了  
结束