

攻防世界crypto easy_ECC

原创

小元砸 于 2020-12-05 20:49:42 发布 866 收藏 5

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订阅专栏

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刚开始看见6分就知道不简单

```
7dcbc9f6f95c48dfaa24f47ca1e74da5.txt - 记事本
文件(E) 编辑(E) 格式(O) 查看(V) 帮助(H)
已知椭圆曲线加密Ep(a,b)参数为

p = 15424654874903

a = 16546484

b = 4548674875

G(6478678675,5636379357093)

私钥为

k = 546768

求公钥K(x,y)

https://blog.csdn.net/qq_49480008
```

看见文本之后联想到了py脚本(但可惜卑微的我不会)所以就照搬了大佬的脚本

```
import collections
import random

EllipticCurve = collections.namedtuple('EllipticCurve', 'name p a b g n h')

curve = EllipticCurve(
    'secp256k1',
    # Field characteristic.
    p=int(input('p=')),
    # Curve coefficients.
    a=int(input('a=')),
```

```

b=int(input('b=')),
# Base point.
g=(int(input('Gx=')),
   int(input('Gy='))),
# Subgroup order.
n=int(input('k=')),
# Subgroup cofactor.
h=1,
)

# Modular arithmetic #####

def inverse_mod(k, p):
    """Returns the inverse of k modulo p.
    This function returns the only integer x such that (x * k) % p == 1.
    k must be non-zero and p must be a prime.
    """
    if k == 0:
        raise ZeroDivisionError('division by zero')

    if k < 0:
        # k ** -1 = p - (-k) ** -1 (mod p)
        return p - inverse_mod(-k, p)

    # Extended Euclidean algorithm.
    s, old_s = 0, 1
    t, old_t = 1, 0
    r, old_r = p, k

    while r != 0:
        quotient = old_r // r
        old_r, r = r, old_r - quotient * r
        old_s, s = s, old_s - quotient * s
        old_t, t = t, old_t - quotient * t

    gcd, x, y = old_r, old_s, old_t

    assert gcd == 1
    assert (k * x) % p == 1

    return x % p

# Functions that work on curve points #####

def is_on_curve(point):
    """Returns True if the given point lies on the elliptic curve."""
    if point is None:
        # None represents the point at infinity.
        return True

    x, y = point

    return (y * y - x * x * x - curve.a * x - curve.b) % curve.p == 0

def point_neg(point):
    """Returns -point."""

```

```

assert is_on_curve(point)

if point is None:
    #  $-\theta = \theta$ 
    return None

x, y = point
result = (x, -y % curve.p)

assert is_on_curve(result)

return result

def point_add(point1, point2):
    """Returns the result of point1 + point2 according to the group law."""
    assert is_on_curve(point1)
    assert is_on_curve(point2)

    if point1 is None:
        #  $\theta + point2 = point2$ 
        return point2
    if point2 is None:
        #  $point1 + \theta = point1$ 
        return point1

    x1, y1 = point1
    x2, y2 = point2

    if x1 == x2 and y1 != y2:
        #  $point1 + (-point1) = \theta$ 
        return None

    if x1 == x2:
        # This is the case point1 == point2.
        m = (3 * x1 * x1 + curve.a) * inverse_mod(2 * y1, curve.p)
    else:
        # This is the case point1 != point2.
        m = (y1 - y2) * inverse_mod(x1 - x2, curve.p)

    x3 = m * m - x1 - x2
    y3 = y1 + m * (x3 - x1)
    result = (x3 % curve.p,
              -y3 % curve.p)

    assert is_on_curve(result)

    return result

def scalar_mult(k, point):
    """Returns k * point computed using the double and point_add algorithm."""
    assert is_on_curve(point)

    if k < 0:
        #  $k * point = -k * (-point)$ 
        return scalar_mult(-k, point_neg(point))

```

```

result = None
addend = point

while k:
    if k & 1:
        # Add.
        result = point_add(result, addend)

        # Double.
        addend = point_add(addend, addend)

    k >>= 1

assert is_on_curve(result)

return result

# Keypair generation and ECDHE #####

def make_keypair():
    """Generates a random private-public key pair."""
    private_key = curve.n
    public_key = scalar_mult(private_key, curve.g)

    return private_key, public_key

private_key, public_key = make_keypair()
print("private key:", hex(private_key))
print("public key: (0x{:x}, 0x{:x})".format(*public_key))

```

解出X, Y的值

```
private key: 0x837d0  
public key: (0xcb19fe553fa, 0x50545408eb4)
```

再16进制转10进制可得

- 2进制 4进制 8进制 10进制 16进制 32进制

转换数字

- 2进制 4进制 8进制 10进制 16进制 32进制

转换结果

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题目要求答案是x+y所以得出答案为

5520194834100+13957031351290
19 477 226 185 390

()	%	C
7	8	9	÷
4	5	6	x
1	2	3	-
0	.	=	+

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