

# 南邮ctf mysql\_南邮ctf知识点汇总 - Crypto篇

原创

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本文链接: [https://blog.csdn.net/weixin\\_35089715/article/details/113710882](https://blog.csdn.net/weixin_35089715/article/details/113710882)

版权

easy!:

```
bmN0Znt0aGlzX2lzX2Jhc2U2NF9lbnNvZGV9
```

base64解码获得flag

```
import base64
```

```
encoded = "bmN0Znt0aGlzX2lzX2Jhc2U2NF9lbnNvZGV9"
```

```
base64.b64decode(encoded)
```

Keyboard:

按顺序在键盘上比划一下获得flag。。。

异性相吸:

提示说xor, 直接xor明文密文获得flag

```
#coding=utf-8
```

```
encrypted=[]
```

```
with open("./biubiubiu/密文.txt").decode('utf-8') as f:
```

```
while True:
```

```
c = f.read(1)
```

```
if not c:
```

```
break
```

```
encrypted.append(c)
```

```
plain=[]
```

```
with open("./biubiubiu/明文.txt").decode('utf-8') as f:
```

```
while True:
```

```
c = f.read(1)
```

```
if not c:
```

```
break
```

```
plain.append(c)
```

```
flag=""
for i in range(len(encrypted)):
    flag+=chr(ord(encrypted[i])^ord(plain[i]))
print flag
```

Wiener Wiener Chicken Dinner:

RSA wiener attack:

```
import math
def continued_fractions_expansion(numerator,denominator):#(e,N)
    result=[]
    dividend=numerator%denominator
    quotient=numerator/denominator
    result.append(quotient)
    while dividend!=0:
        numerator=numerator-quotient*denominator
        tmp=denominator
        denominator=numerator
        numerator=tmp
        dividend=numerator%denominator
        quotient=numerator/denominator
        result.append(quotient)
    return result
def convergents(expansion):
    convergents=[[expansion[0],1]]
    for i in range(1,len(expansion)):
        numerator=1
        denominator=expansion[i]
        for j in range(i-1,-1,-1):
            numerator+=expansion[j]*denominator
        if j==0:
            break
        tmp=denominator
```

```

denominator=numerator

numerator=tmp

convergents.append((numerator,denominator))#(k,d)

return convergents

def newtonSqrt(n):

approx = n/2

better = (approx + n/approx)/2

while better != approx:

approx = better

better = (approx + n/approx)/2

return approx

def wiener_attack(cons,e,N):

for cs in cons:

k,d=cs

if k==0:

continue

phi_N=(e*d-1)/k

#x**2-((N-phi_N)+1)*x+N=0

a=1

b=-((N-phi_N)+1)

c=N

delta = b*b - 4*a*c

if delta<=0:

continue

x1= (newtonSqrt(delta)-b)/(2*a)

x2=-((newtonSqrt(delta)+b)/(2*a)

if x1*x2==N:

return [x1,x2,k,d]

N=10630453212838444683445311689927785206511921621009485339915390974470314400900681918356
e=83716502291837631897269158916049137522937219562594013712174068543253013286054101017472

```

```
expansion=continued_fractions_expansion(e,N)
```

```
cons=convergents(expansion)
```

```
p,q,k,d=wiener_attack(cons,e,N)
```

```
print p
```

```
print q
```

```
print k
```

```
print d
```

获得p,q,k,d:

```
4264602184014247587924993536335689217032123661655205384003117212319820376607574005589577
```

```
2492718606365308154128369833763705307352576059745368683640484366658631317809868149671669
```

```
45596980786762657225114362796367275264021019060430838376323202096126956909718
```

```
57899763801722261062891290503559835904571946557258761154422546104824094670843
```

把d, p, q带入rsa,

```
privatekey = RSA.construct((long(n), long(e), long(d), long(p), long(q)))
```

```
from Crypto.PublicKey import RSA
```

```
from Crypto.Cipher import PKCS1_v1_5 as Cipher_pkcs1_v1_5
```

```
import base64
```

```
from Crypto.Hash import SHA
```

```
from Crypto import Random
```

```
privatekey=RSA.construct((106304532128384446834453116899277852065119216210094853399153909744'  
8371650229183763189726915891604913752293721956259401371217406854325301328605410101747276
```

```
cipher = Cipher_pkcs1_v1_5.new(privatekey)
```

```
cipher_text =
```

```
'AGgt1h6dudnkeoCr7SFclKYsYa65KZ8V29bbgbf+BDyjnyx5stCYjcyktat73aHs2EOaMgwGUwj3HwPTvT+T5L
```

```
flag=cipher.decrypt(base64.b64decode(cipher_text),83716502291837631897269158916049137522937219562
```

```
print flag
```

获得flag

Baby RSA:

尝试分解n:

总结两种方式:

Search	Sequences	Report results	Factor tables	Status	Downloads	Login
<input type="text" value="16903705997349646195704375376941855414691523387719679999999999999"/> <input type="button" value="Factorize!"/> (2)						
Result:						
status (2)	digits	number				
FF	65 (show)	1690370599...99<65> = 1578173871764844869716052171<28> · 10710927547195113973175047066215146269<38>				

<https://blog.csdn.net/u013596119>

## 2. msieve分解:

msieve153.exe 0x291733BAB061EF9C599139CB3E40A5C762B6F448FFFFFFFFFFFFFFFF -v

```

Msieve v. 1.53 (SVN 1005)
Sun Dec 23 19:45:46 2018
random seeds: 4421a080 3902936c
factoring 16903705997349646195704375376941855414691523387719679999999999999 (65 digits)
searching for 15-digit factors
commencing quadratic sieve (65-digit input)
using multiplier of 11
using generic 32kb sieve core
sieve interval: 12 blocks of size 32768
processing polynomials in batches of 17
using a sieve bound of 122849 (5838 primes)
using large prime bound of 6756695 (22 bits)
using trial factoring cutoff of 23 bits
polynomial 'A' values have 8 factors

sieving in progress (press Ctrl-C to pause)
6154 relations (2706 full + 3448 combined from 29596 partial), need 5934
6154 relations (2706 full + 3448 combined from 29596 partial), need 5934
sieving complete, commencing postprocessing
begin with 32302 relations
reduce to 9148 relations in 2 passes
attempting to read 9148 relations
recovered 9148 relations
recovered 7435 polynomials
attempting to build 6154 cycles
found 6154 cycles in 1 passes
distribution of cycle lengths:
  length 1 : 2706
  length 2 : 3448
largest cycle: 2 relations
matrix is 5838 x 6154 (0.8 MB) with weight 167383 (27.20/col)
sparse part has weight 167383 (27.20/col)
filtering completed in 3 passes
matrix is 5437 x 5501 (0.7 MB) with weight 146293 (26.59/col)
sparse part has weight 146293 (26.59/col)
commencing Lanczos iteration
memory use: 0.8 MB
lanczos halted after 88 iterations (dim = 5435)
recovered 64 nontrivial dependencies
p28 factor: 1578173871764844869716052171
p38 factor: 10710927547195113973175047066215146269
elapsed time 00:00:14

```

<https://blog.csdn.net/u013596119>

获得

p1=1578173871764844869716052171

p2=10710927547195113973175047066215146269

已知p1,p2,n,e,求d, 并且解密获得flag:

```
import gmpy2
```

p1=1578173871764844869716052171

p2=10710927547195113973175047066215146269

n=0x291733BAB061EF9C599139CB3E40A5C762B6F448FFFFFFFFFFFFFFF

e=0x10001

phi\_n=(p1-1)\*(p2-1)

d=gmpy2.invert(e,phi\_n)

print hex(d)

c=0x237200C0F72B97DB55BA37C7AACBB61A26A0CB47D294726259C4DF

m=pow(c,d,n)

m\_hex=hex(m)[2:]

m\_str = str(bytearray.fromhex(m\_hex))

print m\_str

简单的方法:

rsa-tool 2 by te! 写入数据直接获得答案

