

【CTF WriteUp】2020第四届强网杯部分Crypto题解

原创

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写在前边

强网杯还是难。。去年正赛赛题一道都不会，只能靠临时补充的强网先锋题目拿分的情景历历在目。今年也没好哪去，只能写一点是一点吧。

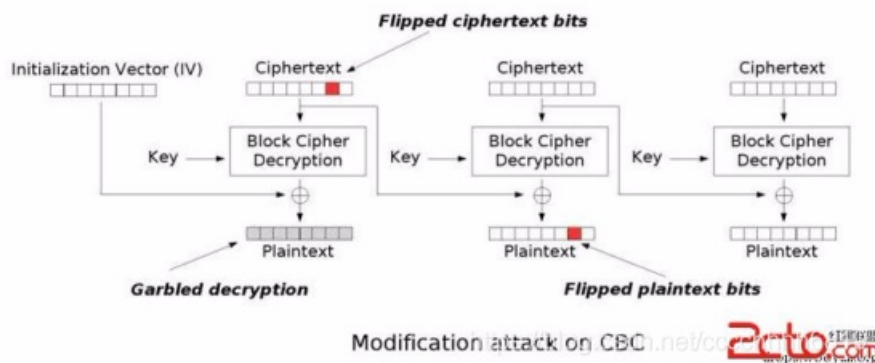
modestudy

这道题是一道六合一块密码大杂烩，考察基础知识与变换，六道小的题目全做完以后拿到flag。其中4、5、6三道题需要大量交互，因此答案不变，可以单次做完后保留答案一起提交。1、2、3三道题当场变换即可。

Stage 1

```
[/$] challenge 1
[+] cookie=session=6b1f33a78c5b9c17;admin=0;checksum=552ebbeb9276a8cd9f741b7d29f8c9e1eb455757207ac420659e7eab56dde25
[+] checksum=aes128cbc.encrypt(session=6b1f33a78c5b9c17;admin=0)
[+] only admin can see flag
[-] cookie: (待输入)
```

本题采用CBC模式加密。我们还是借用红黑联盟那个CBC解密的图说事



如图所示，需要将admin=0改为admin=1，只要把上一段密文的对应位异或'0'再异或'1'即可。这样虽然会破坏上一段明文，但是无所谓，admin=1构造出来了

```
session=6b1f33a7      552ebbeb9276a8cd9f741b7d29f8c9e1
8c5b9c17;admin=0    eb455757207ac420659e7eab56dde25
```

把第一行最后一位异或1即可过关

```
[ $\$$ ] challenge 1
[+] cookie:session=6043fb92858e4080;admin=0;checksum=e1b54fe73c59fc5d07542002d90b7ef53e730401cc8736d4312e6ca55f4758a6
[+] checksum=aes128cbc.encrypt(session=6043fb92858e4080;admin=0)
[+] only admin can see flag
[-] cookie:session=6043fb92858e4080;admin=0;checksum=e1b54fe73c59fc5d07542002d90b7ef43e730401cc8736d4312e6ca55f4758a6
[+] decrypt(checksum): 樽癩€萌窰賦?CH?58e4080;admin=1
[+] passed
```

Stage 2

```
[ $\$$ ] challenge 2
[+] sha256(iv)=11f595abc9d7b986d24fce986d1f9ddfb0d83f2f978db20b862ea730e570bff0
[+] 1. server's job: print aes_cbc_dec(key,iv,your_input_c).encode('hex')
[+] 2. your job: guess iv
[-] your choice: (待输入)
```

标准的CBC选择密文攻击。还用上边那张图。设我们输入为 c_0 c_1 ，得到输出为 m_0 m_1 ，则

```
dec( $c_0$ ) =  $m_0 \oplus iv$ 
dec( $c_1$ ) =  $m_1 \oplus c_0$ 
```

我们让 $c_1 = c_0$ ，可以看到

```
dec( $c_0$ ) =  $m_0 \oplus iv$ 
dec( $c_1$ ) = dec( $c_0$ ) =  $m_1 \oplus c_0 = m_0 \oplus iv$ 
 $iv = m_0 \oplus m_1 \oplus c_0$ 
```

所以输入两段一样的16字符内容，拿这段内容和两段解密出来的明文异或即为iv

```
[ $\$$ ] challenge 2
[+] sha256(iv)=11f595abc9d7b986d24fce986d1f9ddfb0d83f2f978db20b862ea730e570bff0
[+] 1. server's job: print aes_cbc_dec(key,iv,your_input_c).encode('hex')
[+] 2. your job: guess iv
[-] your choice:1
[-] c:80808080808080808080808080808080
[+] fee1313801108cd3dc7dfb547a3126d82893f3d8e1541cc11799a1546e4fef0e
[+] 1. server's job: print aes_cbc_dec(key,iv,your_input_c).encode('hex')
[+] 2. your job: guess iv
[-] your choice:2
[-] iv(encode hex):ee42fad0d874a822f3d462302c4ef1e6
[+] passed
```

Stage 3

```
[ $\$$ ] challenge 3
[+] cookie=session:0884ce7c;timed1=1;admin=0;guess_cookie_ma=1;guess_mp_ab=1;guess_cookie_mb=0;hell_pad=233
[+] 128bit_ecb_encrypt(cookie):0632b3e7adb2f6d5ae4a92f553f2f4a4c9f95b099cfd8e3408137d134eb51d147045f278246a831fabdbdccc0e99b4b694b07e699ffae6a82c6cfbc4454816b1c78e5d1be8b67e235fbbfd2a75e73f32bac806808bc2e102db8f2c159b250415
[+] only admin can see the flag
[-] input your encrypted cookie(encode hex): (待输入)
```

ECB模式的加密每块之间相互独立，加密的六个明文段依次为：

```
session:0884ce7c
;timedl=1;admin=
0;guess_cookie_m
a=1;guess_mp_ab=
1;guess_cookie_m
b=0;hell_pad=233
```

我们只需要把原始的第五段密文覆盖掉原始的第三段密文，即可做出admin=1，通过

```
[$] challenge 3
[+] cookie=session:31e820c8;timedl=1;admin=0;guess_cookie_ma=1;guess_mp_ab=1;guess_cookie_mb=0;hell_pad=233
[+] 128bit_ecb_encrypt(cookie):1ab7386c76b1c8749b05c6fc3b9ef740ef5a0fc8370638514069f522736af63aefa2c39ca7bc127d2
b02d5481a01562da88bf83f7efc639d8fcd953572936ef92397f151ae10e4e8180a0a576ba8327df98eb4a8787959889cb63476eca247a8
[+] only admin can see the flag
[-] input your encrypted cookie(encode hex):1ab7386c76b1c8749b05c6fc3b9ef740ef5a0fc8370638514069f522736af63a2397
f151ae10e4e8180a0a576ba8327da88bf83f7efc639d8fcd953572936ef92397f151ae10e4e8180a0a576ba8327df98eb4a8787959889cb6
3476eca247a8
[+] decrypted:session:31e820c8;timedl=1;admin=1;guess_cookie_ma=1;guess_mp_ab=1;guess_cookie_mb=0;hell_pad=233
[+] passed
```

Stage 4

```
[$] challenge 4
[+] sha256(secret)=d014cbddd2cbb0fa2404c519c166bc85c03ee3445d643f451a5f5d6244e7e34d
[+] assert len(secret)==16
[+] 1. server's job: print aes_ecb(key,input+secret+'\x00'*((16-(len(input+secret) % 16)) % 16))
[+] 2. your job: guess secret
[-] your choice: (待输入)
```

从这里开始就不是拿到题直接能够解决的了，需要计算出结果保存，

由于ECB模式加密每个块互不相关，所以这里可以按位爆破，如：

- (1)首先记录'a' * 15 + 任一字符加密后的的256个值
- (2)输入'a' * 15，看'a' * 15+secret首位是上述哪个
- (3)用'a' * 14 + secret首位代替'a' * 15，重复以上步骤，直到secret全部爆出

(代码最后一块给)

Stage 5

```
[$] challenge 5
[+] sha256(secret)=4c766c8749526dd1a14fdf37619d2fbbebbc2478e25e27e63e591ae0aafb305d
[+] assert len(secret)==16
[+] myblockencrypt_ecb(secret).encode("hex")=4c21cc2bc7941b224ed45bd02ee11b60
[+] In this challenge, you need to try something.
[+] 1. server's job: print myblockencrypt_ecb(your_input)
[+] 2. your job: guess secret
[-] your choice: (待输入)
```

这里果然是 need try something。我们尝试加密几个看看


```

def strxor(a, b):
    return ''.join(chr(ord(a[i])^ord(b[i%len(b)])) for i in range(len(a)))

def mysha256(text):
    mysha = hashlib.sha256()
    mysha.update(text)
    hashresult = mysha.digest()
    bits=''.join(bin(ord(j))[2:].zfill(8) for j in hashresult)
    return bits

def passPow(text):
    print text
    count = 0
    while True:
        bits = mysha256(text + "a" + str(count))
        if(bits.startswith("00000")):
            break
        count += 1
    return "a" + str(count)

def solve_step_4(level, known):
    print "Level %s" % str(level)
    if(level == 16):
        return known
    records = []
    for i in range(256):
        p.recvuntil("[-] your choice:")
        p.sendline("1")
        p.recvuntil("[-] input(encode hex):")
        tmpstr = ('a'*(15-level) + known + chr(i)).encode('hex')
        p.sendline(tmpstr)
        text = p.recvuntil("[+] 2. your job: guess secret")
        s = text.split("[+] ")[2].strip()
        s = s[15:47]
        print "Process: %s/256 %s %s" % (str(i), tmpstr, s)
        records.append(s)
    p.recvuntil("[-] your choice:")
    p.sendline("1")
    p.recvuntil("[-] input(encode hex):")
    tmpstr = ('a'*(15-level)).encode('hex')
    p.sendline(tmpstr)
    text = p.recvuntil("[+] 2. your job: guess secret")
    s = text.split("[+] ")[2].strip()
    s = s[15:47]
    newknown = known + chr(records.index(s))
    print "====="
    print newknown.encode('hex')
    print "====="
    return solve_step_4(level+1, newknown)

def pad4(text):
    return '0'*(4-len(text))+text

def solve_step_5(cipher):
    f = open("step5.txt", 'w+')
    for i in range(0, 65536, 8):
        print "Processing: %s/65536" % str(i)
        tmpstr = ""
        for j in range(8):

```

```

for j in range(8):
    tmpstr += pad4(long_to_bytes(i+j).encode('hex'))
p.recvuntil("[-] your choice:")
p.sendline("1")
p.recvuntil("[-] input(encode hex):")
p.sendline(tmpstr)
text = p.recvuntil("[+] 2. your job: guess secret")
cipher = text.split("[+] ")[1]
cipher = cipher[45:]
for j in range(8):
    f.write("'" + tmpstr[4*j:4*j+4] + "' ")
f.write("=== ")
for j in range(8):
    f.write("\'" + cipher[4*j:4*j+4] + "\' ")
f.write("\r\n")
return ""

def solve_step_6(c0, level, known):
    tmpknown = strxor(known, chr(level))
    newknown = ""
    if(level >= 17):
        return known
    for i in range(256):
        print "Process: %s/256" % str(i)
        tmpcipher = '1'*16 + '\x00'*(16-level) + chr(i) + tmpknown + c0
        p.recvuntil("[-] your choice:")
        p.sendline("1")
        p.recvuntil("[-] input your iv+c (encode hex):")
        p.sendline(tmpcipher.encode('hex'))
        text = p.recvuntil("[+] 2. your job: guess secret")
        if text.find("[+] unpadding success") >= 0:
            newknown = chr(i^level) + known
            print "======"
            print newknown.encode('hex')
            print "======"
            break
    return solve_step_6(c0, level+1, newknown)

print p.recvline().strip()
# pass the PoW
tmpline = p.recvline()
print tmpline.strip()
powtext = passPow(tmpline[11:19])
print p.recvuntil("[-] ?=")
print "SEND: %s" % powtext
p.sendline(powtext)
print p.recvuntil("[+] teamtoken=")
print "SEND: %s" % team_token
p.sendline(team_token)
print p.recvuntil("[-] your choice:")

# Step 1
print "SEND: 1"
p.sendline("1")
text = p.recvuntil("[-] cookie:")
print text
s = text.split("[+] ")[1].strip()
result = s[:79] + chr(ord(s[79:81].decode('hex'))^0x1).encode('hex') + s[81:]
print "SEND: %s" % result[7:]

```



```
print "SEND: %s" % result
p.sendline(result)
print p.recvuntil("[-] your choice:")

# # Step 5
# print "SEND: 5"
# p.sendline("5")
# text = p.recvuntil("[+] 2. your job: guess secret")
# print text
# ciphertext = text.split("[+] ")[3].strip()
# ciphertext = ciphertext[41:]
# result = solve_step_5(ciphertext)
# print "SEND: %s" % result
# p.sendline(result)
# print p.recvuntil("[-] your choice:")

# Step 5(finished)
result = "6122db344a73a14e8247fb2856fbbf94"
print "SEND: 5"
p.sendline("5")
print p.recvuntil("[-] your choice:")
print "SEND: 2"
p.sendline("2")
print p.recvuntil("[-] secret(encode hex):")
print "SEND: %s" % result
p.sendline(result)
print p.recvuntil("[-] your choice:")

# # Step 6
# print "SEND: 6"
# p.sendline("6")
# text = p.recvuntil("[+] 2. your job: guess secret")
# s = text.split("[+] ")[4].strip()
# s = s[70:]
# result = solve_step_6(s[:32].decode('hex'), 1, "")
# print result.encode('hex')
# result = strxor(result, '1')
# print result.encode('hex')

# Step 6(finished)
result = "0a1c9e1464925d23d4a3068313b407ee".decode('hex')
result = strxor(result, '1')
print "SEND: 6"
p.sendline("6")
print p.recvuntil("[-] your choice:")
print "SEND: 2"
p.sendline("2")
print p.recvuntil("[-] secret(encode hex):")
print "SEND: %s" % result.encode('hex')
p.sendline(result.encode('hex'))
print p.recvuntil("[-] your choice:")

p.interactive()
```



```
[+] menu (0 - unsolve, 1 - solved):
[1] challenge-1 status-1
[2] challenge-2 status-1
[3] challenge-3 status-1
[4] challenge-4 status-1
[5] challenge-5 status-1
[6] challenge-6 status-1
[7] get your flag
[-] your choice:
[*] Switching to interactive mode
$ 7
```

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强网先锋-baby_crt

(好久没见过这类纯数学变换的题目了，里边部分变换步骤可能有些跳步，熟悉数论的应该能看懂，不熟悉的请先去补充一下基础知识，或尝试自行证明)

首先我们知道

$$S = (Cp * Sp + Cq * Sq) \% (n * t1 * t2)$$

所以

$$\begin{aligned} S \% t1 &= (Cp * Sp + Cq * Sq) \% t1 \\ S \% t1 &= (q * t2 * \text{inv}(q * t2, p * t1)) * \text{pow}(m + k, dp, p * t1) + \text{一个}t1\text{的倍数} \% t1 \\ S \% t1 &= (a * p * t1 + 1) * \text{pow}(m + k, dp, p * t1) \% t1 \\ S \% t1 &= \text{pow}(m + k, dp, p * t1) \% t1 \\ S \& t1 &= \text{pow}(m + k, dp, t1) \end{aligned}$$

于是

$$\begin{aligned} c1 &= (m - \text{pow}(S, et1, t1) + 1) \% t1 \\ c1 &= (m - \text{pow}(m + k, dp * et1, t1) + 1) \% t1 \end{aligned}$$

我们知道

$$\begin{aligned} et1 &= \text{inv}(d, t1-1) \\ d * et1 &= x * (t1-1) + 1 \\ dp &= d \% ((p-1)*(t1-1)) \\ d &= dp + y * (p-1) * (t1-1) \end{aligned}$$

所以 $(dp + y * (p-1) * (t1-1)) * et1 = x * (t1-1) + 1$
 $dp * et1 = 1 \pmod{t1-1}$

t1是质数，m+k大概率与t1互质，根据费马小定理

$$\begin{aligned} c1 &= (m - \text{pow}(m + k, dp * et1, t1) + 1) \% t1 \\ c1 &= (m - (m + k) + 1) \% t1 \\ c1 &= (1-k) \% t1 \end{aligned}$$

同理可以推出

$$c2 = 1 \% t2 = 1$$

于是

```
sig = pow(S, c1 * c2, n)
```

变成了

```
sig = pow(S, c1, n)
```

这个c1是 $(1-k) \% t1$ ，而k是65536以内随机的质数，因此不能直接求。接下来我们想办法处理这个sig。处理sig的关键点在S。回到S的式子

```
S = (Cp * Sp + Cq * Sq) % (p * q * t1 * t2)
```

```
S % p = (Cp * Sp + Cq * Sq) % p
```

```
S % p = (q * t2 * inv(q * t2, p * t1)) * pow(m + k, dp, p * t1) + 一个p的倍数
```

```
S % p = (a * p * t1 + 1) * pow(m + k, dp, p * t1) % p
```

```
S % p = pow(m + k, dp, p)
```

同理推出

```
S % q = pow(m, dq, q)
```

根据中国剩余定理，可以求出 $S \% n$ 的通解：

```
S % n = pow(m+k,dp,p)*q*inv(q,p)+pow(m,dq,q)*p*inv(p,q)
```

于是有

```
sig = pow(S, c1, n)
```

```
sig = (pow(m+k,dp,p)*q*inv(q,p) + pow(m,dq,q)*p*inv(p,q)) ^ c1 % n
```

这里是形如 $(a+b)^n$ 的形式，用二项式定理展开，注意到前者能被q整除，后者能被p整除，所以展开后只剩下一头一尾不能被n整除，其余全部消掉。即：

```
sig = (pow(m+k,dp,p)*q*inv(q,p))^c1 + (pow(m,dq,q)*p*inv(p,q))^c1 % n
```

```
sig % q = (pow(m,dq,q)*p*inv(p,q))^c1 % q
```

```
sig % q = pow(m,dq*c1,q)
```

```
sig^e % q = pow(m,dq*c1*e,q)
```

```
sig^e % q = pow(m,c1,q)
```

此处最后一步是因为

```
dq = d % ((q-1)*(t2-1))
```

```
d = dq + a*(q-1)*(t2-1)
```

```
e*d = e*dq + e***(q-1)*(t2-1)
```

```
b*(p-1)*(q-1)+1 = e*dq + e***(q-1)*(t2-1)
```

```
1 = e*dq (mod q-1)
```

```
pow(m, e*dq, q) = m
```

于是我们可以通过爆破c1，计算 $sig^e - m^c1$ 与n的最大公约数得到q。完整代码如下：

```
#!/usr/bin/env python
# -*- coding: utf-8 -*-
import gmpy2
n = 263183583822582157708277707633846033595244445661461340392720652066571355134968973219839206522421821124794841
3534343620681572260575655709824188723383724851903187944474092278935135613832294710834683395640564757883887342565
8405513192437479359531790697924285889505666769580176431360506227506064132034621123828090480606055877425480739950
809109048177976884825589023444901953529913585288143291544181183810227538919739159609515261544693445870832956400
3487687431861099115305846281136961555547057146951747286546950202503054845129690985766766996372036629008406247058
3318590585472209798523021029182199921435625983186101089395997
m = 262754933207060261441969663988861968338151704138077058052877634130131009628317037746403327655038380874349048
3565798827606466030442780296160918599796466544086741690071112851785926750465762716059870024868973804524314211148
9179673375819308779535247214660694211698799461044354352200950309392321861021920968200334344131893259850468214901
2662080904692658097295142491439380435215796782347546700970562815568618055680966574159748055782991964403627919074
0888895891706366886720825737009932408484074243578596068180162518061132494895365766674219505149261061383062973163
3827861546693629268844700581558851830936504144170791124745540
sig = 2015294136912288841413007500284576404691272747171683985467128025584579892873810382459533988534540541994335
4215456598381228519131902698373225795339649300359363119754605698321052334731477127433796964107633109608706030111
1971567016073790867669440960666493233679767863830151066818964794468354191432258323209785305543998510741807623083
2209233972183956664214490886453046601761473167952539225979651178962408022858708062145408495716919334372451586746
8178242402356741884890739873250658960438450287159439457730127074563991513030091456771906853781028159857466498315
359846665211412644316716082898396009119848634426989676119219246
e = 65537
tmpsig = pow(sig, e, n)
for c1 in range(65536):
    tmpq = (pow(m, c1, n)-tmpsig)%n
    if(gmpy2.gcd(tmpq, n)>1):
        q = gmpy2.gcd(tmpq, n)
        p = n // q
        print "p = " + str(p)
# p = 1495804442330860257901795734148567115562926352190284922506763092333069266983476721148819388583646630096045
8429325108750521166559508456764501367184374293096528397369172598278322608545027685073789281983047910702936043730
3876681309102929919702367620122904659518881851295355796994990602745486313662851309278333637
```

强网先锋-红方辅助


```
#!/usr/bin/env python
# -*- coding: utf-8 -*-
import struct

lines = open("tcp.txt", 'r').readlines()

for i in range(0, len(lines), 5):
    btime = lines[i+1].strip().decode('hex')
    boffset = struct.unpack("<i", lines[i+2].strip().decode('hex'))[0]
    fn = chr(int(lines[i+3][16:18], 16))
    salt = int(lines[i+3][18:20], 16)

    t = struct.unpack("<i", btime)[0]
    t -= boffset
    t = struct.pack("<i", t)

    m = ""
    count = 0
    c = lines[i+3][20:].strip().decode('hex')
    for j in range(len(c)):
        if(fn == '0'):
            m += chr(((ord(c[j]) + salt) & 0xff) ^ ord(t[count]))
        elif(fn == '1'):
            m += chr(((ord(c[j]) - salt) & 0xff) ^ ord(t[count]))
        elif(fn == '2'):
            m += chr(((ord(c[j]) ^ salt) & 0xff) ^ ord(t[count]))
        count = (count + 1) % 4
    print m.strip()
```

强网先锋-bank

一个模拟区块链交易的题目，题目核心内容是要伪造交易记录。根据提示，每一条交易记录都是

```
enc(发送方) + enc(接收方) + enc(金额)
```

虽然我们不知道key，但是明显我们不需要知道，从view records的交易记录里直接搬运就可以了。解题所需要的三个要素如下获取：

1. enc(我)：请给Alice转1块钱，得到的交易记录前32位就是；
2. enc(别人)：同上，交易记录第32~64位就是；
3. enc(金额)：把交易记录里边的大额都读一遍就够了。

由于本题不允许使用重复记录，所以或者构造Alice给我转不同金额的记录，或者构造不同的人给我转相同金额的记录，不影响解题。解题代码如下：

```

#!/usr/bin/env python
# -*- coding: utf-8 -*-
from pwn import *
import hashlib
import string

def mysha256(text):
    mysha = hashlib.sha256()
    mysha.update(text)
    return mysha.hexdigest()

def passPoW(suffix, target):
    for a in string.printable:
        for b in string.printable:
            for c in string.printable:
                tmp = a+b+c+suffix
                if mysha256(tmp)==target:
                    return a+b+c

teamtoken = "队伍token"
p = remote("39.101.134.52", 8005)
text = p.recvuntil("Give me XXX:")
print text
suffix = text.split("+")[1][:17]
target = text.split(" == ")[1][:64]
result = passPoW(suffix, target)
p.sendline(result)
print p.recvuntil("teamtoken:").strip()
p.sendline(teamtoken)
print p.recvuntil("give me your name:").strip()
p.sendline("chainer")
print p.recvuntil("you can choose: transact, view records, provide a record, get flag, hint").strip()
p.sendline("transact")
print p.recvuntil("please give me the trader and the amount(for example:Alice 1)").strip()
p.sendline("Alice 1")
text = p.recvuntil("you can choose: transact, view records, provide a record, get flag, hint")
print text.strip()
s = text.split('\n')[1]
nameme = s[2:34]
nameAlice = s[34:66]

p.sendline("view records")
text = p.recvuntil("you can choose: transact, view records, provide a record, get flag, hint")
print text.strip()
newtrans = []
for i in range(2,12):
    s = text.split('\n')[i]
    newtrans.append(nameAlice + nameme + s[64:96])
print newtrans

for i in range(10):
    p.sendline("provide a record")
    print p.recvuntil("My system is secure if you can give me other records, the receiver can also get the money
.").strip()
    p.sendline(newtrans[i])
    print p.recvuntil("you can choose: transact, view records, provide a record, get flag, hint").strip()

p.sendline("get flag")
p.interactive()

```

```
your cash:850
you can choose: transact, view records, provide a record, get flag, hint
> My system is secure if you can give me other records, the receiver can also get the money.
>
your cash:956
you can choose: transact, view records, provide a record, get flag, hint
> My system is secure if you can give me other records, the receiver can also get the money.
>
your cash:1148
you can choose: transact, view records, provide a record, get flag, hint
> My system is secure if you can give me other records, the receiver can also get the money.
>
your cash:1330
you can choose: transact, view records, provide a record, get flag, hint
> My system is secure if you can give me other records, the receiver can also get the money.
>
your cash:1477
you can choose: transact, view records, provide a record, get flag, hint
[*] Switching to interactive mode

> you need pay 1000 for the flag!
```

<https://blog.csdn.net/ccchhhh6819>

最后说两句

国密那个题看着就不想做怎么破。。一点动力都没有