

[buuctf] crypto全解——85-120（不建议直接抄flag）

原创

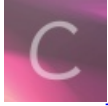
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订阅专栏

85.[ACTF新生赛2020]crypto-rsa3

查看题目

```
//output.txt
1776065048364992469709590302268716088859693217782110510805246340845169733314416449938980295736122900958530692640
3653045925365287558626794687783105514754691022710056649665814838183468303736613455384801190325125272647404766127
4223137727688689535823533046778793131902143444408735610821167838717488859902242863683
1457390378511382354771000540945361168984775052693073641682375071407490851289703070905749525830483035988737117653
9714284246123320209259266173955588681603806019124982999228259142295101669579104518417300289198838076344898341288
30801407228447221775264711349928156290102782374379406719292116047581560530382210049
```

```
//rsa3.py
from flag import FLAG
from Cryptodome.Util.number import *
import gmpy2
import random

e=65537
p = getPrime(512)
q = int(gmpy2.next_prime(p))
n = p*q
m = bytes_to_long(FLAG)
c = pow(m,e,n)
print(n)
print(c)
```

很明显文本里面的一堆数字就是n和c的值

那就很简单了, 我们用yafu来分解N, 也可以在线分解

```

factor(177606504836499246970959030226871608885969321778211051080524634084516973331441644993898029573612290095853
0692640365304592536528755862679468778310551475469102271005664966581483818346830373661345538480119032512527264740
47661274223137727688689535823533046778793131902143444408735610821167838717488859902242863683)

fac: factoring 1776065048364992469709590302268716088859693217782110510805246340845169733314416449938980295736122
9009585306926403653045925365287558626794687783105514754691022710056649665814838183468303736613455384801190325125
2726474047661274223137727688689535823533046778793131902143444408735610821167838717488859902242863683
fac: using pretesting plan: normal
fac: no tune info: using qs/gnfs crossover of 95 digits
div: primes less than 10000
fmt: 100000 iterations
Total factoring time = 0.6794 seconds

***factors found***

P155 = 133269090503574476435265858368339693780781470577230547014328421929887176493857314300950556223035495772334
95793715580004801634268505725255565021519817179293
P155 = 133269090503574476435265858368339693780781470577230547014328421929887176493857314300950556223035495772334
95793715580004801634268505725255565021519817179231

ans = 1

```

pq的值youlepy文件里面还有e

那么我们直接带入脚本就可以了（前面写了那么多脚本，我觉得你可以创一个文件夹，到时候把这些带入就可以了，而不是每次都重新敲）

```

import gmpy2

p = 133269090503574476435265858368339693780781470577230547014328421929887176493857314300950556223035495772334957
93715580004801634268505725255565021519817179293
q = 133269090503574476435265858368339693780781470577230547014328421929887176493857314300950556223035495772334957
93715580004801634268505725255565021519817179231
c=14573903785113823547710005409453611689847750526930736416823750714074908512897030709057495258304830359887371176
5397142842461233202092592661739555886816038060191249829992282591422951016695791045184173002891988380763448983412
8830801407228447221775264711349928156290102782374379406719292116047581560530382210049
n=17760650483649924697095903022687160888596932177821105108052463408451697333144164499389802957361229009585306926
4036530459253652875586267946877831055147546910227100566496658148381834683037366134553848011903251252726474047661
274223137727688689535823533046778793131902143444408735610821167838717488859902242863683

e=65537

phin = (p-1)*(q-1)

d=gmpy2.invert(e,phin)

print(d)

m = pow(c,d,n)

print(m)
print(hex(m))
print(bytes.fromhex(hex(m)[2:]))

```

运行得到

```
1142110340697533342134076532322394538183055540690133289094348258061535212792815766150851970390819896211244297276
8902360002419925826186240372948581699098548699684783004877571490662028879663621214255323576437411112094114014989
5595289559560740951667825123772519353303069022288432998660657137564563315041240411393
893441512863695667867454629314548913750576288726628646106246575194731343607717680126106810553506844477056381
0x616374667b705f616e645f715f73686f756c645f6e6f745f62655f736f5f636c6f73655f696e5f76616c75657d
b'actf{p_and_q_should_not_be_so_close_in_value}'
```

86.[RoarCTF2019]babyRSA

查看题目

```
import sympy
import random

def myGetPrime():
    A= getPrime(513)
    print(A)
    B=A-random.randint(1e3,1e5)
    print(B)
    return sympy.nextPrime((B!)%A)
p=myGetPrime()
#A1=21856963452461630437348278434191434000660767504190274938524635134698652620643408366138310666023009597726323
97773487317560339056658299954464169264467234407
#B1=21856963452461630437348278434191434000660767504190274938524635134698652620643408366138310666023009597726323
97773487317560339056658299954464169264467140596

q=myGetPrime()
#A2=164661131158392281197678878993088200257492609338634468882241671698576121786641395457263408674067907545602275
16013796269941438076818194617030304851858418927
#B2=164661131158392281197678878993088200257492609338634468882241671698576121786641395457263408674067907545602275
16013796269941438076818194617030304851858351026

r=myGetPrime()

n=p*q*r
#n=8549266378627529215983160339108387617514935430932767300871662765071816058563972310079334753464962833041663125
5660901307533909900431413447524262332232659153047067908693481947121069070451562822417357656432171870951184673132
5542136901233080426973619699863603750609547029206563641441541458128385583653341729359314414240962702061406918146
6231856269692576799193736978262790840823908735803316541002069015206771571111273225203858843289675840589870901034
2467882264362733
c=pow(flag,e,n)
#e=0x1001
#c=7570088302166957773932931679545070620450263580231073147715699883471082077024521946870324530200999893206708038
3977560299708060476222089630209972629755965140317526034680452483360917378812244365884527186056341888615564335560
765053550155758362271622330017433403027261127561225585912484778295885012139611106904519876255027013314851416396
8435642731690512299575982524113387273436271604181981994864566280329241880220443087452134210841362363515047596312
1220095236776428
#so,what is the flag?
```

做这题看到阶乘一下想到了 gxy2020 的一题，也是考到了威尔逊定理（Wilson's theorem）：当且仅当 p 为素数时： $(p-1)! \equiv -1 \pmod{p}$ 。

阶乘只乘到 B ，所以把 $(B+1)$ 乘到 $(A-1)$ 这一段也补上就得到了威尔逊公式，反之我们可以由用 -1 乘这一段的模反数，就得到了题目中的 $(B!)\%A$ 。

```

from Crypto.Util.number import *
from sympy import nextprime
A1=2185696345246163043734827843419143400006607675041902749385246351346986526206434083661383106660230095977263239
7773487317560339056658299954464169264467234407
B1=2185696345246163043734827843419143400006607675041902749385246351346986526206434083661383106660230095977263239
7773487317560339056658299954464169264467140596

A2=1646611311583922811976788789930882002574926093386344688822416716985761217866413954572634086740679075456022751
6013796269941438076818194617030304851858418927
B2=1646611311583922811976788789930882002574926093386344688822416716985761217866413954572634086740679075456022751
6013796269941438076818194617030304851858351026

def f(a,b,P):# a*(a+1)*...*b (mod P)
    ans=1
    for i in range(a,b+1):
        ans*=i
        ans%=P
    return ans%P

inv1=inverse(f(B1+1,A1-1,A1),A1)
ans1=((A1-1)*inv1)%A1
p=nextprime(ans1)

inv2=inverse(f(B2+1,A2-1,A2),A2)
ans2=((A2-1)*inv2)%A2
q=nextprime(ans2)

e=0x1001
c=75700883021669577739329316795450706204502635802310731477156998834710820770245219468703245302009998932067080383
9775602997080604762220896302099726297559651403175260346804524833609173788122443658845271860563418886155643355607
6505355015575836227162233001743340302726112756122558591248477782958850121396111069045198762550270133148514163968
4356427316905122995759825241133872734362716041819819948645662803292418802204430874521342108413623635150475963121
220095236776428
n=85492663786275292159831603391083876175149354309327673008716627650718160585639723100793347534649628330416631255
66090130753390990043141434475242623322326591530470679086934819471210690704515628224173576564321718709511846731325
5421369012330804269736196998636037506095470292065636414415414581283855836533417293593144142409627020614069181466
2318562696925767991937369782627908408239087358033165410020690152067715711112732252038588432896758405898709010342
467882264362733

r=(n//p)//q
assert isPrime(r)
d=inverse(e,(p-1)*(q-1)*(r-1))
m=pow(c,d,n)
print(long_to_bytes(m))

```

运行得到

```
b'RoarCTF{wm-CongrAtulation4-1t4-ju4t-A-bAby-R4A}'
```

87.坏蛋是雷宾

查看题目

老牌刺客之王混进了女王的住所。一天，女王得到了一个匿名举报，说她的侍卫里有一个刺客，叫做Rabin，而他的信息就在一份文件里，文件中有附带一个Pk，是523798549，密文是162853095，校验码二进制值是110001，根据说明是放在明文后一起加密的，明文与密文长度相同。加密算法和这位老牌刺客同名。快拯救女王，答案是求得的明文，进行32位md5小写哈希字符串，提交即可。注意：得到的 flag 请包上 flag{} 提交

n分解然后写脚本

```

'''// python2'''
from gmpy2 import *
import hashlib
n=523798549
p=10663
q=49123
e=2
c=162853095
inv_p = invert(p, q)
inv_q = invert(q, p)

mp = pow(c, (p + 1) / 4, p)
mq = pow(c, (q + 1) / 4, q)

a = (inv_p * p * mq + inv_q * q * mp) % n
b = n - int(a)
c = (inv_p * p * mq - inv_q * q * mp) % n
d = n - int(c)

for i in (a, b, c, d):
    print(bin(i)[2:])
#在得出的四个解中，找到二进制符合题目的解
#10010011100100100101010110001
m='10010011100100100101010'
mc=str(int(m,2))
md=hashlib.md5()
md.update(mc.encode("utf8"))
flag = md.hexdigest()
print("flag{"+str(flag)+'}')
#flag{ca5cec442b2734735406d78c88e90f35}

```

88.[AFCTF2018]Single

[查看题目](#)

Jmqrida rva Lfmz (JRL) eu m uqajemf seny xl enlxdomrexn uajiderc jxoqarerexnu. Rvada mda rvdaa jxooxn rcqau xl JRLu: Paxqmdyc, Mrrmjs-Yalanja mny oekay.

Paxqmdyc-urcfa JRLu vmu m jxiqfa xl giaurexnu (rmusu) en dmnza xl jmrazxdeau. Lxd akmoqfa, Wab, Lxdanuej, Jdcqrux, Benmdc xd uxoarvenz afua. Ramo jmn zmen uxoa qxenru lxd atadc uxftay rmus. Oxda qxenru lxd oxda jxoqfejmrax rmusu iuimffc. Rva nakr rmus en jvmen jmn ba xqanay xnfc mlrad uxoa ramo uxfta qdatexiu rmus. Rvan rva zmoa reoa eu xtad uio xl qxenru uvwxu cxi m JRL wenad. Lmoxiu akmoqfa xl uijv JRL eu Yaljxn JRL gimfu.

Waff, mrrmjs-yananja eu mnxrvad enradaurenz seny xl jxoqarerexnu. Vada atadc ramo vmu xwn narwxds(xd xnfc xna vxur) werv tifnmdmbfa uadtejau. Cxid ramo vmu reoa lxd qmrjvenz cxid uadtejau mny yatafxqenz akqfxeru iuimffc. Ux, rvan xdzmnehadu jxnajru qmdrejeqmru xl jxoqarerexn mny rva wmdzmoa urmdru! Cxi uvxify qdxrajr xwn uadtejau lxd yalanja qxenru mny vmjs xqqxanru lxd mrrmjs qxenru. Veurxdejmffc rveu eu m ledur rcqa xl JRLu, atadcbxyc snxu mbxir YAL JXN JRL - uxoarvenz fesa m Wxdfy Jiq xl mff xrvad jxoqarerexnu.

Oekay jxoqarerexnu omc tmde qxuuebfalxdomru. Er omc ba uxoarvenz fesa wmdzmoa werv uqajemf reoa lxd rmus-bmuay afaoanru (a.z. IJUB eJRL).

JRL zmoau xlraxrijv xn omnc xrvad muqajru xl enlxdomrexn uajiderc: jdcqruxzdmqvc, uraxx, benmdc mnmfcueu, datada anzanaadenz, oxbefa uajiderc mny xrvadu. Zxxy ramou zanadmffc vmta urdxnz useffu mny akqadeanja en mff rvaua euuiiau.

Iuimffc, lfmz eu uxoa urdenz xl dmnyxo ymrm xd rakr en uxoa lxdomr. Akmoqfa mljrl{Xv_I_lxiny_er_neja_rDc}

我看到这我就直接去在线爆破了

The screenshot shows the quipqiup website interface. At the top, there's a navigation bar with the site name 'quipqiup BETA'. Below that, a description states it's a fast and automated cryptogram solver. The main content area features a 'Puzzle:' section with a text box containing a cryptogram: 'Dekaz xkoqaterexnu om tndc quxebfa lxodnnu. Er onc ba uxoaarvna fesa vmdznoa werv uxajent roca lxd rmas-buuay afoanru (a.z. IJUB eJRL). JRL zmoax xlan xxiiv xn onoc xrvad maaixu xl enlkdonexm uxajere: jdcqxzdqvc, uraxx, benmdc anmofnuo, datada anzanadenz, oabefa ualiderc any xrvadu. Exxy rakou zanadaffc vata urdomz useffu nny xkqadeanja en aff rvaasa muuau. Duiffc. Ifaz eu uxoa urdomz xl dmozxo yman ad raxk en uxoa lxodnn. Almoefa aljil [Xv I lxavv er neja xlc] ClUES: For example G=R QVW=THE'. Below the puzzle is an 'auto' button and a 'Solve' button. A large advertisement for '爆款【云峰A型】主机' (Hot-selling Cloud Peak A-type Host) is overlaid on the page, offering 1G virtual hosts for 78 yuan/year. At the bottom, a list of search results is visible, with the first result being a definition of Capture the Flag (CTF) competitions.

爆破结果

```
0 -1.448 Capture the Flag (CTF) is a special kind of information security competitions. There are three common types of CTFs: Jeopardy, Attack-Defence and mixed. Jeopardy-style CTFs has a couple of questions (tasks) in range of categories. For example, Web, Forensic, Crypto, Binary or something else. Team can gain some points for every solved task. More points for more complicated tasks usually. The next task in chain can be opened only after some team solve previous task. Then the game time is over sum of points shows you a CTF winner. Famous example of such CTF is Defcon CTF quals. Well, attack-defence is another interesting kind of competitions. Here every team has own network(or only one host) with vulnarable services. Your team has time for patching your services and developing exploits usually. So, then organizers connects participants of competition and the wargame starts! You should protect own services for defence points and hack opponents for attack points. Historically this is a first type of CTFs, everybody knows about DEF CON CTF - something like a World Cup of all other competitions. Mixed competitions may vary possible formats. It may be something like wargame with special time for task-based elements (e.g. UCSB iCTF). CTF games often touch on many other aspects of information security: cryptography, stego, binary analysis, reverse engineering, mobile security and others. Good teams generally have strong skills and experience in all these issues. Usually, flag is some string of random data or text in some format. Example afctf{0h_U_found_it_nice_tTry}
```

```
1 -3.331 Capture the Flag (CTF) is a special kind of information security competitions. There are three common types of CTFs: Jeopardy, Attack-Defence and mixed. Jeopardy-style CTFs has a couple of questions (tasks) in range of categories. For example, Web, Forensic, Crypto, Binary or something else. Team can gain some points for every solved task. More points for more complicated tasks usually. The next task in chain can be opened only after some team solve previous task. Then the game time is over sum of points shows you a CTF winner. Famous example of such CTF is Defcon CTF quals. Well, attack-defence is another interesting kind of competitions. Here every team has own network(or only one host) with vulnarable services. Your team has time for patching your services and developing exploits usually. So, then organizers connects participants of competition and the wargame starts! You should protect own services for defence points and hack opponents for attack points. Historically this is a first type of CTFs, everybody knows about DEF CON CTF - something like a World Cup of all other competitions. Mixed competitions may vary possible formats. It may be something like wargame with special time for task-based elements (e.g. UCSB iCTF). CTF games often touch on many other aspects of information security: cryptography, stego, binary analysis, reverse engineering, mobile security and others. Good teams generally have strong skills and experience in all these issues. Usually, flag is some string of random data or text in some format. Example afctf{0h_U_found_it_nice_tTry}
```

```
2 -3.341 Capture the Flag (CTF) is a special kind of information security competitions. There are three common types of CTFs: Jeopardy, Attack-Defence and mixed. Jeopardy-style CTFs has a couple of questions (tasks) in range of categories. For example, Web, Forensic, Crypto, Binary or something else. Team can gain some points for every solved task. More points for more complicated tasks usually. The next task in chain can be opened only after some team solve previous task. Then the game time is over sum of points shows you a CTF winner. Famous example of such CTF is Defcon CTF quals. Well, attack-defence is another interesting kind of competitions. Here every team has own network(or only one host) with vulnarable services. Your team has time for patching your services and developing exploits usually. So, then organizers connects participants of competition and the wargame starts! You should protect own services for defence points and hack opponents for attack points. Historically this is a first type of CTFs, everybody knows about DEF CON CTF - something like a World Cup of all other competitions. Mixed competitions may vary possible formats. It may be something like wargame with special time for task-based elements (e.g. UCSB iCTF). CTF games often touch on many other aspects of information security: cryptography, stego, binary analysis, reverse engineering, mobile security and others. Good teams generally have strong skills and experience in all these issues. Usually, flag is some string of random data or text in some format. Example afctf{0h_U_found_it_nice_tTry}
```

you solve tasks. More points for more complicated tasks usually. The next task in chain can be opened only after some team solve previous tasks. Then the game time is over sum of points shows you a CTF timer. Famous example of such CTF is Defcon CTF Quals. Well, attack-defence is another interesting kind of competitions. Here every team has own network (or only one host) with vulnerable services. Your team has time for patching your services and developing exploits usually. So, then organizers connects participants of competition and the game starts! You should protect own services for defence points and hack opponents for attack points. Historically this is a first type of CTFs, everybody knows about DEF CON CTF - something like a World Cup of all other competitions. Mixed competitions may vary possible formats. It may be something like game with special time for task-based elements (e.g. UCSB iCTF). CTF games often touch on many other aspects of information security: cryptography, stego, binary analysis, reverse engineering, mobile security and others. Good teams generally have strong skills and experience in all these issues. Usually, flag is some string of random data or text in some format. Example afctf{Oh_U_found_it_nice_tRy}

3 -3.344 Capture the Flag (CTF) is a special kind of information security competitions. There are three common types of CTFs: Jeopardy, Attack-Defence and mixed. Jeopardy-style CTFs has a couple of questions (tasks) in range of categories. For example, Web, Forensic, Crypto, Binary or something else. Team can gain some points for every solved task. More points for more complicated tasks usually. The next task in chain can be opened only after some team solve previous tasks. When the game time is over sum of points shows you a CTF timer. Famous example of such CTF is Defcon CTF Quals. Well, attack-defence is another interesting kind of competitions. Where every team has own network (or only one host) with vulnerable services. Your team has time for patching your services and developing exploits usually. So, then organizers connects participants of competition and the game starts! You should protect own services for defence points and hack opponents for attack points. Historically this is a first type of CTFs, everybody knows about DEF CON CTF - something like a World Cup of all other competitions. Mixed competitions may vary possible formats. It may be something like game with special time for task-based elements (e.g. UCSB iCTF). CTF games often touch on many other aspects of information security: cryptography, stego, binary analysis, reverse engineering, mobile security and others. Good teams generally have strong skills and experience in all these issues. Usually, flag is some string of random data or text in some format. Example afctf{Oh_U_found_it_nice_tRy}

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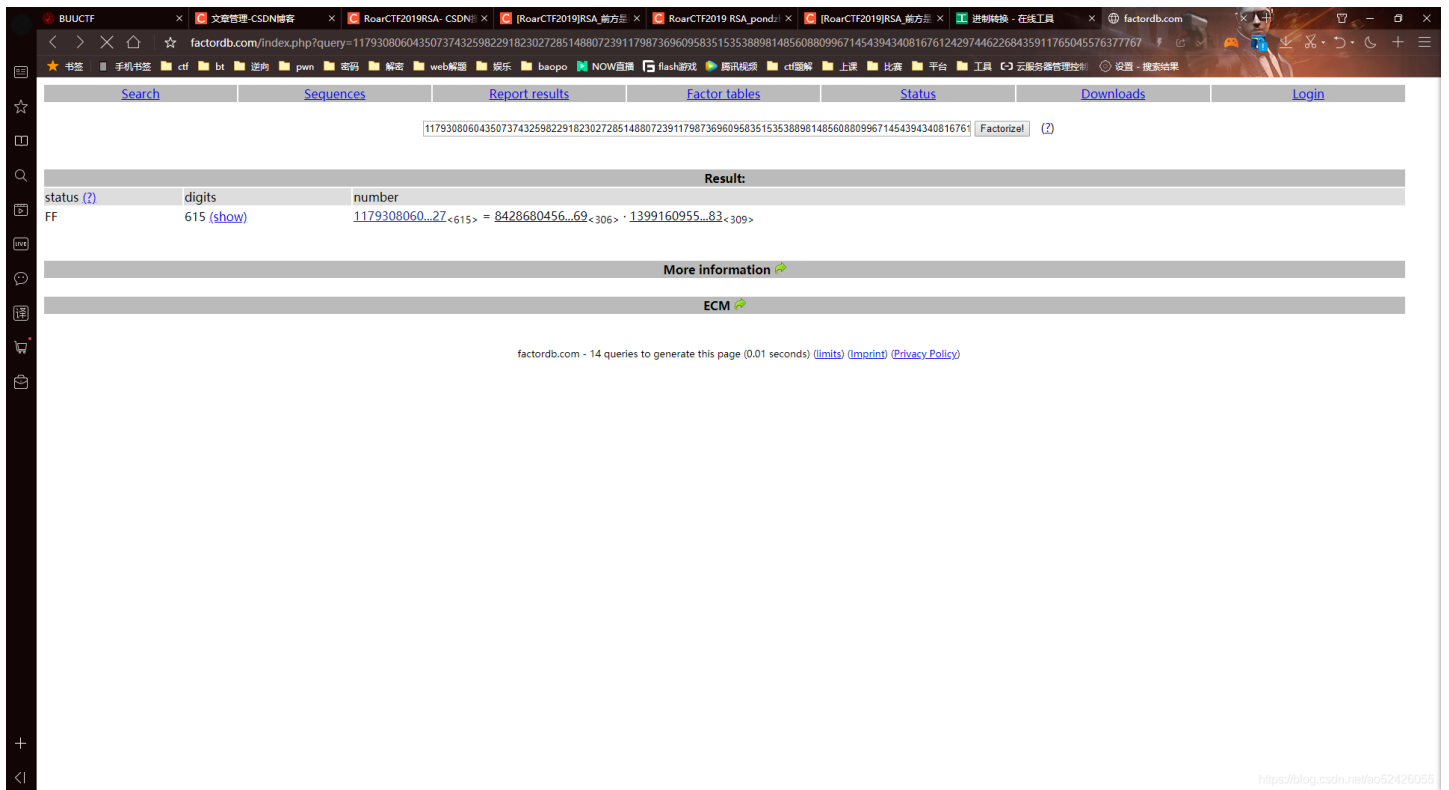
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拿第一个提交即可afctf{Oh_U_found_it_nice_tRy}

89.[RoarCTF2019]RSA

```
A=(((y%x)**5)%(x%y))**2019+y**316+(y+1)/x
p=next_prime(z*x*y)
q=next_prime(z)
A =
268334918267871452424746951279347600986101478100492490548412748030816
137776819286806156188657704864643238212896088148746342741417611448688
583069395940498974322910351692443251272419565442570345361271031058716
441703587830839067661259284875028738731812942419520862344029464781736
774087821194914752628709129830748050289746227910257255682223166943827
931747482847908971904638641197110544872391059471041809397704417994980
037322435472917983339321982778938907886929021756951123086896764796308
943059425881514636218725085516689755305607374458294614847206833416744
5499314471518357535261186318756327890016183228412253724
n =
117930806043507374325982291823027285148807239117987369609583515353889
814856088099671454394340816761242974462268435911765045576377767711593
100416932019831889059333166946263184861287975722954992219766493089630
810876984781113645362450398009234556085330943125568377741065242183073
882558834603430862598066786475299918395341014877416901185392905676043
795425126968745185649565106322336954427505104906770493155723995382318
346714944184577894150229037758434597242564815299174950147754426950251
419204917376517360505024549691723683358170823416757973059354784142601
436519500811159036795034676360028928301979780528294114933347127
p=8428680456813909345397399592018475522849801799588796679330784539509
685661516621472670062935717654631372705941511386957789861651113804288
065455935880783653313130842300146187144129595848434215866741626883219
428893699123920318826209949442419871530781563894703701955142858507365
41078623854327959382156753458569
```

这道题看着有A什么的但是给了NC
咱还是试一下在线分解N



结果就是一分解就分解出来了
然后有npqc直接写脚本即可（e一般为65537）

```

import gmpy2

q = 842868045681390934539739959201847552284980179958879667933078453950968566151662147267006293571765463137270594
1511386957789861651113804288065455935880783653313130842300146187144129595848434215866741626883219428893699123920
31882620994944241987153078156389470370195514285850736541078623854327959382156753458569

p = 139916095583110895133596833227506693679306709873174024876891023355860781981175916446323044732913066880786918
6290890234993117034084891511818865685356210086449979719821824267065925512910840079833879110062614425196354054570
77292515085160744169867410973960652081452455371451222265819051559818441257438021073941183

c=41971850275428383625653350824107291609587853887037624239544762751558838294718672159979929266922528917912189124
7132736739480514642265196058037451713407243437058321985546801967986232638066179980724960260199404763249716969285
5115937197020736574151706429595637680929727254180064774788517090573786856800010102914392379200348679327819705132
6716680212726111099439262589341050943913401067673851885114314709706016622157285023272496793595281054074260451116
2138159348433178948988832153622895993661010180815132151207282971313524390669304522818294465865620622425273296725
75620261776042653626411730955819001674118193293313612128

n=11793080604350737432598229182302728514880723911798736960958351535388981485608809967145439434081676124297446226
8435911765045576377767711593100416932019831889059333166946263184861287975722954992219766493089630810876984781113
6453624503980092345560853309431255683777410652421830738825588346034308625980667864752999183953410148774169011853
9290567604379542512696874518564956510632233695442750510490677049315572399538231834671494418457789415022903775843
4597242564815299174950147754426950251419204917376517360505024549691723683358170823416757973059354784142601436519
500811159036795034676360028928301979780528294114933347127

e=65537

phin = (p-1)*(q-1)

d=gmpy2.invert(e,phin)

print(d)

m = pow(c,d,n)

print(m)
print(hex(m))
print(bytes.fromhex(hex(m)[2:]))

```

运行得到

```

8599589881775512182490339390302384847126810744233969198532121090013876515514061191844004921719994842305490870513
6826880258908633192226330687534143784850786245106307099225133962814171533637778326481845442321992947664719004853
9278805029351560101212744826887241218280590799690114110729314081810416033936818232121737223480952384234472254960
4286239338414176997138752498663184064331483582259621245748238876057665171100280468834141443144340932719393320666
9179048022566244019931295809893893457165624563454551217020906061061854657248221799501001805487219916158911768825
67105125169912160252167465495939533501038099782250065
2222445378180646636797706006131586301300018090566246580838853918878845
0x526f61724354467b776d2d6c316c316c316c316c316c3131316c6c7d
b'RoarCTF{wm-1111111111111111}'

```

flag包裹提交即可

90.[WUSTCTF2020]B@se

查看题目

密文: MyLkTaP3FaA7KOWjTmKkVjWjVzKjdeNvTnAjoH9iZOlvTeHbvD==
JASGBWcQPRXEFLbCDIlmnHUVKTYZdMowwipatNOefghq56rs****kxyz012789+,

oh holy shit, something is missing...

哦，天哪，有东西不见了。

先看一下base编码少了那四位

```
import string
s = "JASGBWcQPRXEFLbCDIlmnHUVKTYZdMowwipatNOefghq56rs****kxyz012789+/"
for i in string.ascii_letters + string.digits:
    if(i not in s):
        print(i)
```

得到

```
j
u
3
4
```

补全跑一下

```
from Crypto.Util.number import *
from gmpy2 import *
from functools import reduce
import sympy
import itertools

# JASGBWcQPRXEFLbCDIlmnHUVKTYZdMowwipatNOefghq56rs****kxyz012789+/
# coding:utf-8
def My_base64_encode(inputs, s):
    bin_str = []
    for i in inputs:
        x = str(bin(ord(i))).replace('0b', '')
        bin_str.append('{:0>8}'.format(x))
    # print(bin_str)
    outputs = ""
    nums = 0
    while bin_str:
        temp_list = bin_str[:3]
        if (len(temp_list) != 3):
            nums = 3 - len(temp_list)
            while len(temp_list) < 3:
                temp_list += ['0' * 8]
        temp_str = "".join(temp_list)
        # print(temp_str)
        temp_str_list = []
        for i in range(0, 4):
            temp_str_list.append(int(temp_str[i * 6:(i + 1) * 6], 2))
        # print(temp_str_list)
        if nums:
            temp_str_list = temp_str_list[0:4 - nums]

        for i in temp_str_list:
            outputs += s[i]
        bin_str = bin_str[3:]
```

```

outputs += nums * '='
print("Encrypted String:\n%s " % outputs)

def My_base64_decode(inputs, s):
    bin_str = []
    for i in inputs:
        if i != '=':
            x = str(bin(s.index(i))).replace('0b', '')
            bin_str.append('{:0>6}'.format(x))
    # print(bin_str)
    outputs = ""
    nums = inputs.count('=')
    while bin_str:
        temp_list = bin_str[:4]
        temp_str = "".join(temp_list)
        # print(temp_str)
        if (len(temp_str) % 8 != 0):
            temp_str = temp_str[: -1 * nums * 2]
        for i in range(0, int(len(temp_str) / 8)):
            outputs += chr(int(temp_str[i * 8:(i + 1) * 8], 2))
        bin_str = bin_str[4:]
    print("Decrypted String:\n%s " % outputs)

# s = "ABCDEFGHGIJKLMNOPQRSTUVWXYZabcdefghijklmnopqrstuvwxyz0123456789+/"
h = ['j', 'u', '3', '4']
h1 = list(itertools.permutations(h, 4))
for i in h1:
    m = "".join(i)
    s = "JASGBWcQPRXEFLbCDIImnHUVKTYZdMovwipatNOefghq56rs" + m + "kxyz012789+/"
    input_str = "MyLkTaP3FaA7K0WjTmKkVjWjVzKjdeNvTnAjoH9iZ0IvTeHbvD=="
    # My_base64_encode(input_str)
    My_base64_decode(input_str, s)

```

运行得到

```
Decrypted String:
wctf2220{bape64_p_v0ry_e@py_and_fuN}
Decrypted String:
wctf2320{bape64_p_v0ry_e@py_and_fuN}
Decrypted String:
wctf2120{bape64_p_v0ry_e@py_and_fuN}
Decrypted String:
wctf2120{bape64_p_v0ry_e@py_and_fuN}
Decrypted String:
wctf2320{bape64_p_v0ry_e@py_and_fuN}
Decrypted String:
wctf2220{bape64_p_v0ry_e@py_and_fuN}
Decrypted String:
wctf2220{baqe64_q_v1ry_e@qy_and_fuN}
Decrypted String:
wctf2320{baqe64_q_v1ry_e@qy_and_fuN}
Decrypted String:
wctf2120{bare64_!r_v2ry_e@ry_and_fuN}
Decrypted String:
wctf2120{base64_1s_v3ry_e@sy_and_fuN}
Decrypted String:
wctf2320{bare64_!r_v2ry_e@ry_and_fuN}
Decrypted String:
wctf2220{base64_1s_v3ry_e@sy_and_fuN}
Decrypted String:
wctf2020{baqe64_q_v1ry_e@qy_and_fuN}
Decrypted String:
wctf2020{baqe64_q_v1ry_e@qy_and_fuN}
Decrypted String:
wctf2020{bare64_!r_v2ry_e@ry_and_fuN}
Decrypted String:
wctf2020{base64_1s_v3ry_e@sy_and_fuN}
Decrypted String:
wctf2020{bare64_!r_v2ry_e@ry_and_fuN}
Decrypted String:
wctf2020{base64_1s_v3ry_e@sy_and_fuN}
Decrypted String:
wctf2320{baqe64_q_v1ry_e@qy_and_fuN}
Decrypted String:
wctf2220{baqe64_q_v1ry_e@qy_and_fuN}
Decrypted String:
wctf2320{bare64_!r_v2ry_e@ry_and_fuN}
Decrypted String:
wctf2220{base64_1s_v3ry_e@sy_and_fuN}
Decrypted String:
wctf2120{bare64_!r_v2ry_e@ry_and_fuN}
Decrypted String:
wctf2120{base64_1s_v3ry_e@sy_and_fuN}
```

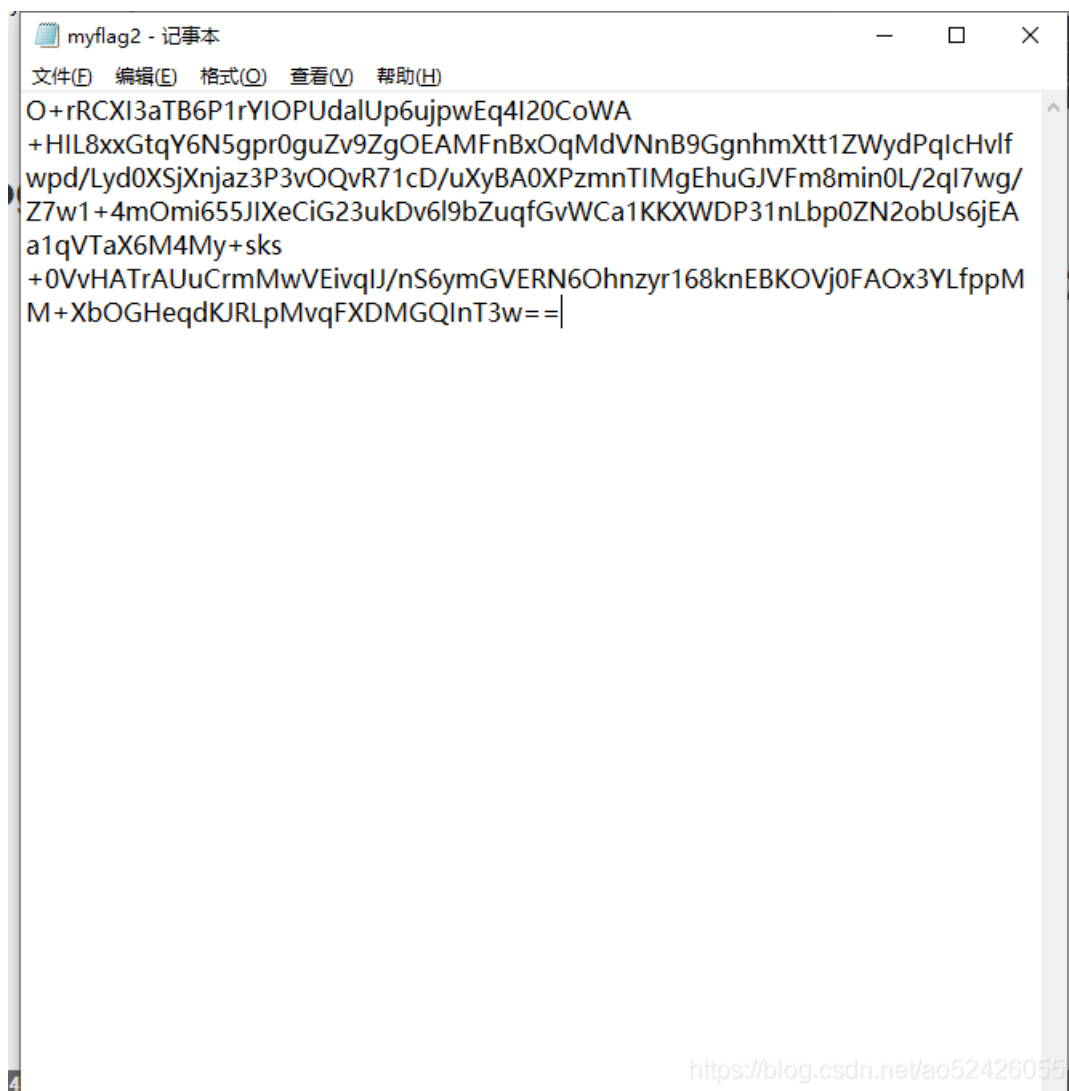
很明显这里的

```
Decrypted String:
wctf2220{base64_1s_v3ry_e@sy_and_fuN}
```

是最通顺的
flag包裹提交

91.[HDCTF2019]together

查看题目



```
myflag2 - 记事本
文件(F) 编辑(E) 格式(O) 查看(V) 帮助(H)
O+rRCXI3aTB6P1rYIOPUdalUp6ujpwEq4I20CoWA
+HIL8xxGtqY6N5gpr0guZv9ZgOEAMFnBxOqMdVnN9GgnhmXtt1ZWydPqlcHvlf
wPd/Lyd0XSjXnjaz3P3vOQvR71cD/uXyBA0XPzmnTIMgEhuGJVfM8min0L/2ql7wg/
Z7w1+4mOmi655JIXeCiG23ukDv6l9bZuqfGvWCa1KKXWDP31nLbp0ZN2obUs6jEA
a1qVTaX6M4My+sks
+0VvHATrAUuCrMmWVEivqIJ/nS6ymGVERN6OhnzYr168knEBKOVj0FAOx3YLfppM
M+XbOGHeqdKJRLpMvqFXDMGQInT3w==|
```

<https://blog.csdn.net/ao52426055>



```
myflag1 - 记事本
文件(F) 编辑(E) 格式(O) 查看(V) 帮助(H)
R3Noy6r3WLItytAmb4FmHEygoilucEEZbO9ZYXx5JN03HNpBLDx7fXd2fl
+UL5+11RCs/y0qITGURWWDtG66eNLzGwNpAKiVj6l7RtUJl2Pcm3NvFeAFwI9UsVR
Eyh7zIV6sI9ZP8l/2GVDorLAz5ULW
+f0OINGhJmZm8FL/aDnlfTElhQ87LPicWpXYoMtyr6WrxjK6Ontn8BqCt0EjQ7TeXZh
xIH9VTPWjDmFdmOqaqdvIT+LZemTgLNESwM5nn4g5S3aFDFwj1YiDYI0/
+8etvKfOrfoKOWR0CxsRHagwdUUTES8EcHlMGCxCkDZn3SzmmA6Nb3lgLeSgG
8P1A==|
```

```
pubkey1.pem - 记事本
文件(F) 编辑(E) 格式(O) 查看(V) 帮助(H)
-----BEGIN PUBLIC KEY-----
MIIBITANBgkqhkiG9w0BAQEFAAOCAQ4AMIIBCQKCAQB1qLiqKtKVDprtS
+NGGN++
q7jLqDJoXMIPRRczMBAGJIRsz5Dzwt1ulr0s5yu8RdaufiYeU6sYIKk92b3yygL
FvaYcZjdqBF2EyTWGVE7PL5lh3rPUfxwQFqDR8EhIH5x+Ob8rjlkftljHTBt1ThJ
JXvDBumXpQKGcBIknRaR9dwR1q8GU58/glk5ND3eCTAadhrhLByWkHbFARxalx4Q
q8s2ZUe8lDc/N6V93EOFjbKbqqqtDmhniF6jdXQDAIwWTpx6+jmzxlCJoVHd2MBs
ZCcQhvkIWtuKz4IYL4+iUpMKGHlhY1vCqFx2EzD4XiljFLP9rk7+9+CoyTulVL/D
AgMACR0=
-----END PUBLIC KEY-----
```

```
pubkey2.pem - 记事本
文件(F) 编辑(E) 格式(O) 查看(V) 帮助(H)
-----BEGIN PUBLIC KEY-----
MIIBIDANBgkqhkiG9w0BAQEFAAOCAQ0AMIIBCQKCAQB1qLiqKtKVDprtS
+NGGN++
q7jLqDJoXMIPRRczMBAGJIRsz5Dzwt1ulr0s5yu8RdaufiYeU6sYIKk92b3yygL
FvaYcZjdqBF2EyTWGVE7PL5lh3rPUfxwQFqDR8EhIH5x+Ob8rjlkftljHTBt1ThJ
JXvDBumXpQKGcBIknRaR9dwR1q8GU58/glk5ND3eCTAadhrhLByWkHbFARxalx4Q
q8s2ZUe8lDc/N6V93EOFjbKbqqqtDmhniF6jdXQDAIwWTpx6+jmzxlCJoVHd2MBs
ZCcQhvkIWtuKz4IYL4+iUpMKGHlhY1vCqFx2EzD4XiljFLP9rk7+9+CoyTulVL/D
AgJbJQ==
-----END PUBLIC KEY-----
```

看到第一个结尾两个等号就很兴奋

结果base解密不出来

看到后面两个就直接KEY就直接公钥解析

```
-----BEGIN PUBLIC KEY-----
MIIBIDANBgkqhkiG9w0BAQEFAAOCAQ0AMIIBCAKCAQElqLiQKtKVDprtS+NGGN++
q7jLqDJoxMlPFRczMBAGIIRsz5DzvtiLulr0s5yu8RdaufiYeU6sYIKk92b3vryL
FvaYCzjdqBF2EYTWGVE7FL5lh3rPUfxwQFqDR8EhIH5x+Ob8xjlkftIjHTBtLThJ
IXvDBumXpQKGCBIknRaR9dwr1q8GU58/gIk5ND3eCTAadhrhLByWkHbFArxaLx4Q
q8s2ZUe8Ldc/W6Y93E0FjbKbqaatDmhniF6jdXQDAIwWTPx6+jmzxLCJovHd2MBs
ZCcQhvk1WtuKz4IYL4+iUpMKGHlhYlvCqFx2EzD4XILjFLP9rk7+9+CoyTuIvL/D
AgIbIQ=
-----END PUBLIC KEY-----
```

解析

详细信息

密钥类型	RSA
密钥强度	2047
PN(e)	23333
PN(n)	1485308127790241124099171958226543729894160685098943265592807574 7449227799832389574251190347654658701773951599098366248661597113 0152215660413055019964516386243894170559569262385959478857400849 9480938293273355698610765349914458861410569451815059410571143898 3069306254763078820574239989253573144558449346681620784979079971 5599761023665272708675274230010831691274021575981834429233644803 8374265311728564302631991424407297555720035354606035274426363786 7557162046429886176035616570590229646013789737629785488326501654 202429466891022723268768841320111523816192606370230314305451686 18446134188815113100443559425057634959299
DER格式	30820120300d06092a864886f70d01010105000382010d00308201080282010075a8b8aa2ad2950e9aed4be34618dfbeabb8cba832685cc94f45173330100624846 cef90f3c2db75ba5af4b39caef1175ab9f898794eac6082a4f766f7cb280b16ff6980b38dda811761324d619513b3cbe65877acf51fc70405a8347c121207e71f8e6 fcae39647ed2231d306dd53849257bc306e997a502867012249d1691f5de11d0af06539f3f808939343dde09301a761ae12c1c969076c502bc5a971e10abcb36654 7bc94373f37a57ddc43858db29baaaad0e6867885ea3757403008c164e9c7afa39b3c65089a151ddd8c06c4271086f9255adb8acf82182f8fa252930a18796163 5bc2a85c761330f85c896314b3fdae4efef7e0a8c93b8854bfc302025b25

```
-----BEGIN PUBLIC KEY-----
MIIBITANBgkqhkiG9w0BAQEFAAOCAQ4AMIIBCAKCAQElqLiQKtKVDprtS+NGGN++
q7jLqDJoxMlPFRczMBAGIIRsz5DzvtiLulr0s5yu8RdaufiYeU6sYIKk92b3vryL
FvaYCzjdqBF2EYTWGVE7FL5lh3rPUfxwQFqDR8EhIH5x+Ob8xjlkftIjHTBtLThJ
IXvDBumXpQKGCBIknRaR9dwr1q8GU58/gIk5ND3eCTAadhrhLByWkHbFArxaLx4Q
q8s2ZUe8Ldc/W6Y93E0FjbKbqaatDmhniF6jdXQDAIwWTPx6+jmzxLCJovHd2MBs
ZCcQhvk1WtuKz4IYL4+iUpMKGHlhYlvCqFx2EzD4XILjFLP9rk7+9+CoyTuIvL/D
AgMACRO=
-----END PUBLIC KEY-----
```

详细信息

密钥类型	RSA
密钥强度	2047
PN(e)	2333
PN(n)	1485308127790241124099171958226543729894160685098943265592807574 7449227799832389574251190347654658701773951599098366248661597113 0152215660413055019964516386243894170559569262385959478857400849 9480938293273355698610765349914458861410569451815059410571143898 3069306254763078820574239989253573144558449346681620784979079971 5599761023665272708675274230010831691274021575981834429233644803 8374265311728564302631991424407297555720035354606035274426363786 7557162046429886176035616570590229646013789737629785488326501654 2024294668910227232687688413201111523816192606370230314305451686 18446134188815113100443559425057634959299
DER格式	30820120300d06092a864886f70d01010105000382010d00308201080282010075a8b8aa2ad2950e9aed4be34618dfbeabb8cba832685cc94f45173330100624846 ccf90f3c2db75ba5af4b39caef1175ab9f898794eac6082a4f766f7cb280b16f6980b38dda811761324d619513b3cbe65877acf51fc70405a8347c121207e71f8e6 fcae39647ed2231d306dd53849257bc306e997a502867012249d1691f5dc11d6af06539f3f808939343dde09301a761ae12c1c969076c502bc5a971e10abc36654 7bc94373f37a57ddc43858db29baaaad0e6867885ea3757403008c164e9c7afa39b3c65089a151ddd8c06c64271086f9255adb8acf82182f8fa252930a18796163 5bc2a85c761330f85c896314b3fdae4efef7e0a8c93b8854bfc30202091d

<https://blog.csdn.net/qq52426055>

两个得到的N是一样的，
base64是一种编码方式而不是一种加密算法
所以将他转换成unicode编码,然后再转换成数
结合共模攻击,flag就出来了

```

import base64
f1="R3Noy6r3WLlitytAmb4FmHEygoilucEEZb09ZYXx5JN03HNpBLDx7fXd2f1+UL5+11RCs/y0q1TGURWWDtG66eNLzGwNpAKiVj6I7RtUJl2Pc
m3NvFeAFwI9UsVREyh7zIV6sI9ZP81/2GVDorLaz5ULW+f00INGHJmZm8FL/aDn1fTElhQ87LPicWpXyOMtyr6WrxjK60ntn8BqCt0EjQ7TeXZhx
IH9VTPWjDmFdm0qaqdvIT+LZemTgLNESwM5nn4g5S3aFDFWj1YiDYl0/+8etvKf0rfoKOWr0CxsRHagwdUUTES8EcHLMGcxCkDzn3SzmA6Nb3l
gLeSg8P1A=="
f2="O+rRCXI3aTB6P1rYIOPUdalUp6ujpwEq4I20CoWa+HIL8xxGtqY6N5gpr0guZv9ZgOEAMFnBx0QMdVNnB9GgnhmXtt1ZWydpqIcHv1fwpd/L
yd0XSjXnjaz3P3v0QvR71cD/uXyBA0XPzmnTIMgEhuGJVfM8min0L/2qI7wg/Z7w1+4mOmi655JIXeCiG23ukDv6l9bZuqfGvWCa1KKXWDP31nLb
p0ZN2obUs6jEAa1qVTaX6M4My+sks+0VvHATrAUuCrMmWVeivqIJ/nS6ymGVERN6OhnzYr168knEBKOVj0FA0x3YLfppMM+Xb0GHeqdKJRLpMvqF
XDMGQInT3w=="

import Crypto.Util.number
c1=Crypto.Util.number.bytes_to_long(base64.b64decode(f1))
c2=Crypto.Util.number.bytes_to_long(base64.b64decode(f2))
n=14853081277902411240991719582265437298941606850989432655928075747449227799832389574251190347654658701773951599
0983662486615971130152215660413055019964516386243894170559569262385959478857400849948093829327335569861076534991
4458861410569451815059410571143898306930625476307882057423998925357314455844934668162078497907997155997610236652
7270867527423001083169127402157598183442923364480383742653117285643026319914244072975557200353546060352744263637
8675571620464298861760356165705902296460137897376297854883265016542024294668910227232687688413201111523816192606
37023031430545168618446134188815113100443559425057634959299
e1=2333
e2=23333
import gmpy2
import binascii
def exgcd(m, n, x, y):
    if n == 0:
        x = 1
        y = 0
        return (m, x, y)
    a1 = b = 1
    a = b1 = 0
    c = m
    d = n
    q = int(c / d)
    r = c % d
    while r:
        c = d
        d = r
        t = a1
        a1 = a
        a = t - q * a
        t = b1
        b1 = b
        b = t - q * b
        q = int(c / d)
        r = c % d
    x = a
    y = b
    return (d, x, y)
ans=exgcd(e1,e2,0,0)
s1=ans[1]
s2=ans[2]
m=(gmpy2.powmod(c1,s1,n)*gmpy2.powmod(c2,s2,n))%n
print(binascii.unhexlify(hex(m)[2:]))

```

运行得到

```
b'flag{23re_SDxF_y78hu_5rFgS}'
```

92.[AFCTF2018]可怜的RSA

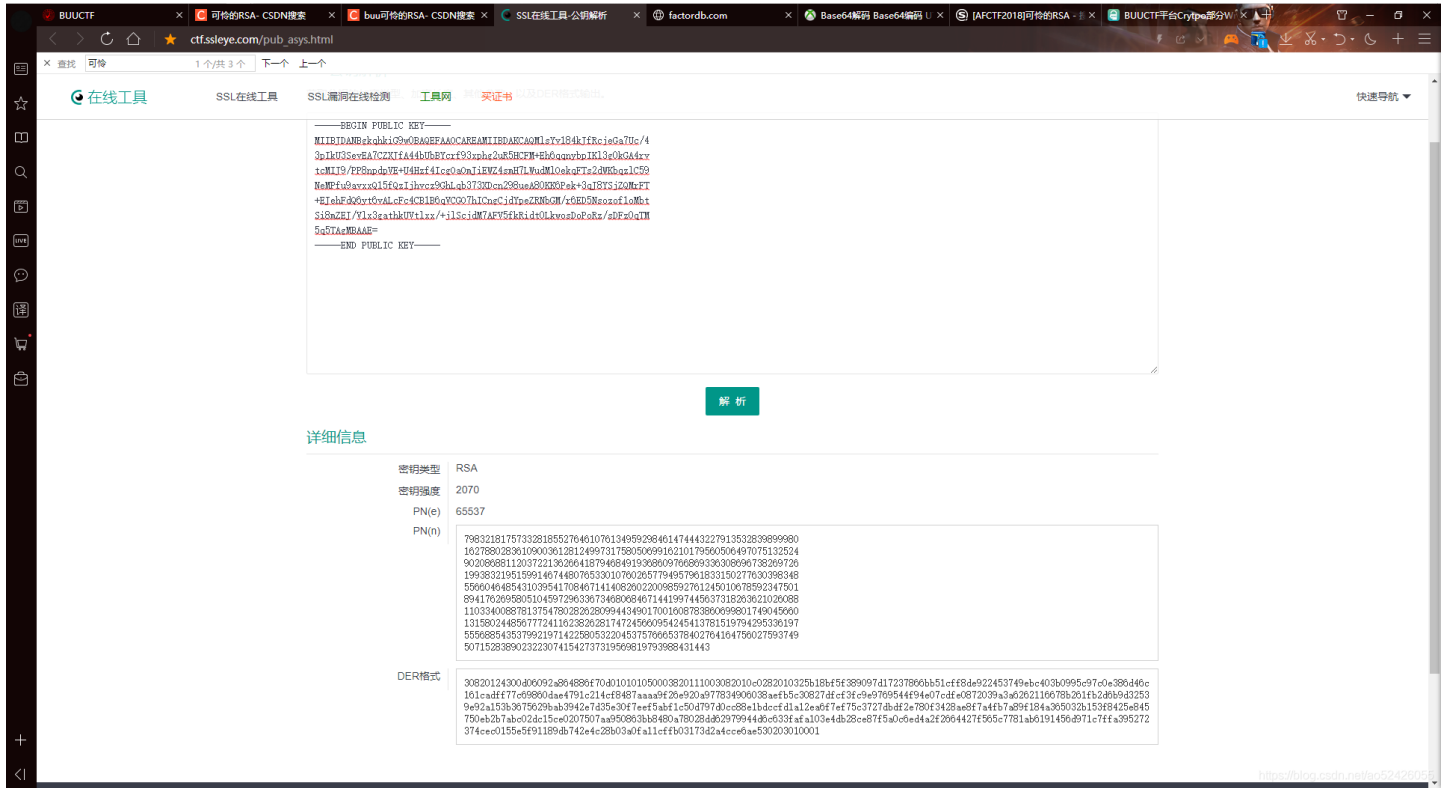
查看题目两个附件

```
flag.enc - 记事本
文件(F) 编辑(E) 格式(O) 查看(V) 帮助(H)
GVd1d3vilXFfcHapEYuo5fAvliUS83adrtMW/MgPwxVBSI46joFCQ1plcnIDGfL19K/3P
vChV6n5QGohzfVyz2Z5GdTIaknxvHDUGf5HCukokyPwK/1EYU7NzrhGE7J5jPdi0Aj7
xi/Odxy0hGMgpaBLd/nL3N8O6i9pc4Gg3O8soOlciBG/6/xdFN3SzSStMYIN8nfZZMS
q3xDDvz4YB7TcTBh4ik4wYhuC77gmT
+HWOv5gLTNQ3EkZs5N3EAopy11zHNYU80yv1jtFGcluNPYXYttU5qU33jcp0Wuzna
c+t+AZHeSQy5vk8DyWorSGMiS+J4KNqSVIDs12EqXEqqJ0uA==
```

<https://blog.csdn.net/ao52426055>

```
public.key - 记事本
文件(F) 编辑(E) 格式(O) 查看(V) 帮助(H)
-----BEGIN PUBLIC KEY-----
MIIBJDANBgkqhkiG9w0BAQEFAAOCAQMIIBDAKCAQMIsYv184kJfRcjeGa7Uc/4
3plkU3SevEA7CZXJfA44bUbBYcrf93xphg2uR5HCFM+Eh6qqnybpIKl3g0kGA4rv
tcMIJ9/PP8npdpVE+U4Hzf4lcgOaOmJiEWZ4smH7LWudMIOekqFTs2dWkbqzIC59
NeMPfu9avxxQ15fQzljhvcz9GhLqb373XDcn298ueA80KK6Pek+3qJ8YSjZQMrFT
s+EJehFdQ6yt6vALcFc4CB1B6qVCGO7hCngCjdYpeZRNbGM/r6ED5Nsozof1oMbt
Si8mZEJ/Vlx3gathkUVtlxx/+jIScjdM7AFV5fkRidt0LkwosDoPoRz/sDFz0qTM
5q5TAgMBAAE=
-----END PUBLIC KEY-----
```


很明显还是一个base一个公钥解密



而且这个N也是可以分解的

Sequences Report results Factor tables Status Downloads

7983218175733281855276461076134959298461474443227913532839899980 1627880283610900361281249973175805 Factorize! (2)

Result:

digits	number
623 (show)	7983218175...43<623> = 3133337 · 2547832606...39<617>

https://blog.csdn.net/ao52426055

分解网站我上面有就不设置超链接了

然后就是常规的求私钥

这边因为flag.enc是RSA的PKCS1_OAEP加密得来的。所以我们这边也是给生成一个私钥文件。

在做到这边的时候，如何导出一个私钥文件。找了一下百度上的方法。都是先generate后给参数分别赋值的。但是我发现我并不行，试了一下python3和python2下的Crypto库都得到一个报错

Exception has occurred: AttributeError
can't set attribute

也就是现在无法通过这么直接赋值了。

这种情况的话，可以去看下python调用的Crypto库里面的RSA模块的一个底层的实现。

发现有一个construct函数，传入一个rsa_components参数，是一个元组型的数据，也就是tuple类型的，分别是(n,e,d,p,q)

```
phi=(p-1)*(q-1)
d=gmpy2.invert(e,phi)
rsa_components=(n,e,int(d),p,q)
arsa=RSA.construct(rsa_components)
arsa.exportKey()
```

然后导出的私钥，对加密后的密文，使用PKCS1_OAEP模块进行解密即可。

脚本如下

```
`from Crypto.PublicKey import RSA
from Crypto.Cipher import PKCS1_OAEP
f=open("public.key","r")
key=RSA.import_key(f.read())
f.close()
e=key.e
n=key.n
import base64
from Crypto.Util.number import *
import gmpy2
p= 3133337
q=25478326064937419292200172136399497719081842914528228316455906211693118321971399936004729134841162974144246271
4864396957860365881174246118819559509962196468073788222782856382615820991083394389495730341012151411561564087428
4382004806683086381436237988572039508231846285000290160568976187631915114735273009095755694084214429988739467874
3607766937828094478336401159449035878306853716216548374273462386508307367713112073004011383418967894930554067582
4532489810220119228833744427368480459206763413618712317871634414675330768900817218821793691687872877247696426653
99992556052144845878600126283968890273067575342061776244939
print(p*q==n)
f=open("flag.enc","r")
c_base64=f.read().strip("\n")
c_bytes=base64.b64decode(c_base64)
c=bytes_to_long(c_bytes)
phi=(p-1)*(q-1)
d=gmpy2.invert(e,phi)
rsa_components=(n,e,int(d),p,q)
arsa=RSA.construct(rsa_components)
rsa_key = RSA.importKey(arsa.exportKey())
rsa_key = PKCS1_OAEP.new(rsa_key)
decrypted = rsa_key.decrypt(c_bytes)
print(decrypted)`
```

运行得到

```
b'afctf{R54_|5_$0_B0rin9}'
```

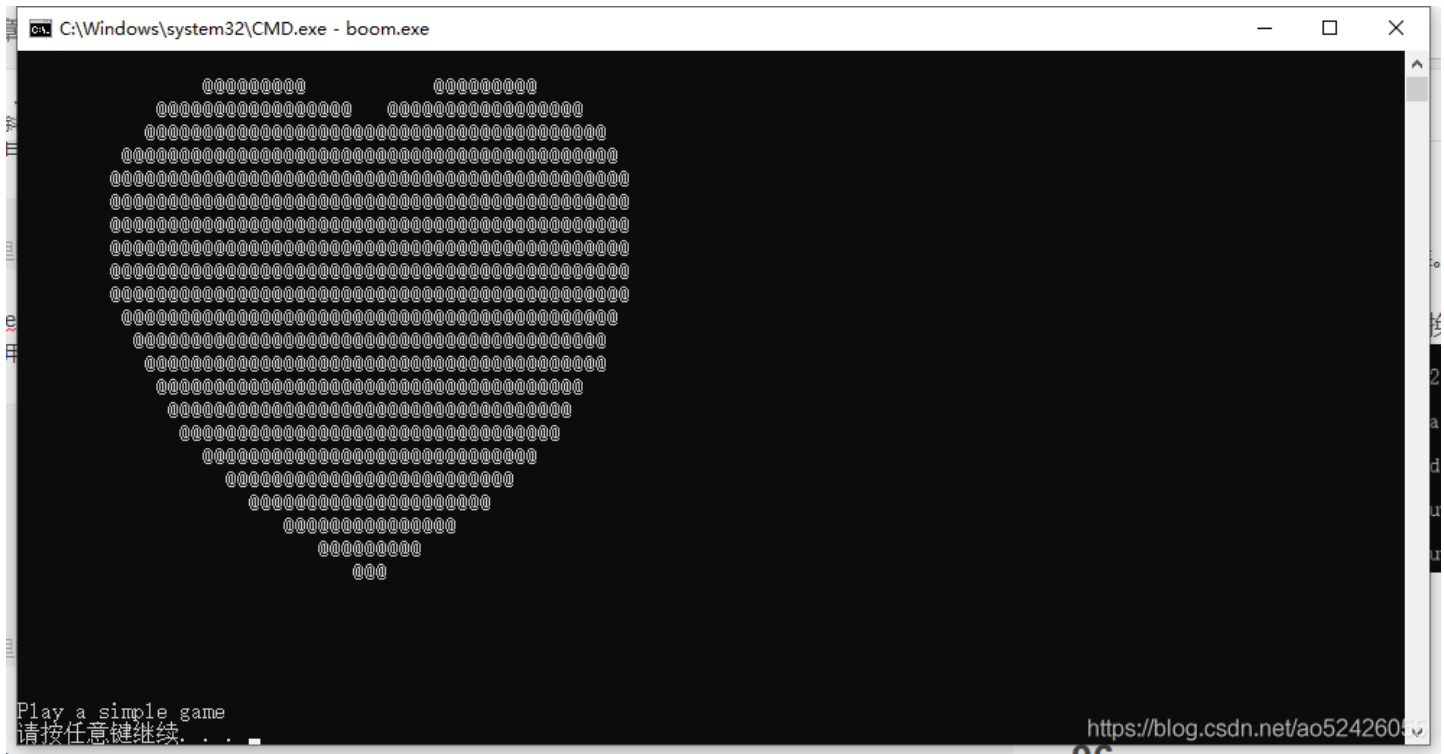
93.[网鼎杯 2020 青龙组]boom

查看题目

boom.exe

一个exe, 先说一下我翻车的时候。。我一开始是直接打开的这个exe但是你写到最后一步, 崩就没了
我这边用min+r打开cmd船靠在切换到exe的

```
C:\Users\25072>d:  
D:\>cd ctfbisai  
D:\ctfbisai>cd buumima  
D:\ctfbisai\buumima>cd "attachment(4)"  
D:\ctfbisai\buumima\attachment(4)>boom.exe
```



```
first:this string md5:46e5efe6165a5afb361217446a2dbd01
```

md5在线解密

```
en5oy
```

一个三元一次方程

```
This time:Here are have some formulas  
3x-y+z=185  
2x+3y-z=321  
x+y+z=173  
input: x =
```

脚本:

```
from sympy import *  
x = Symbol('x')  
y = Symbol('y')  
z = Symbol('z')  
print(solve([3*x-y+z-185,2*x+3*y-z-321,x+y+z-173],[x, y,z]))
```

运行得到

```
{x: 74, y: 68, z: 31}
```

一元二次方程

```
Last time: Kill it
x*x+x-7943722218936282=0
input x:
```

脚本

```
import math
print("ax*x+bx+c=0")
a = float(input("input a: "))
b = float(input("input b: "))
c = float(input("input c: "))
p = b*b-4*a*c
if p < 0:
    print("None")
    exit()
else:
    X1 = (-b+math.sqrt(p))/(2*a)
    X2 = (-b-math.sqrt(p))/(2*a)
print("X1: " + str(X1))
print("X2: " + str(X2))
```

运行得到

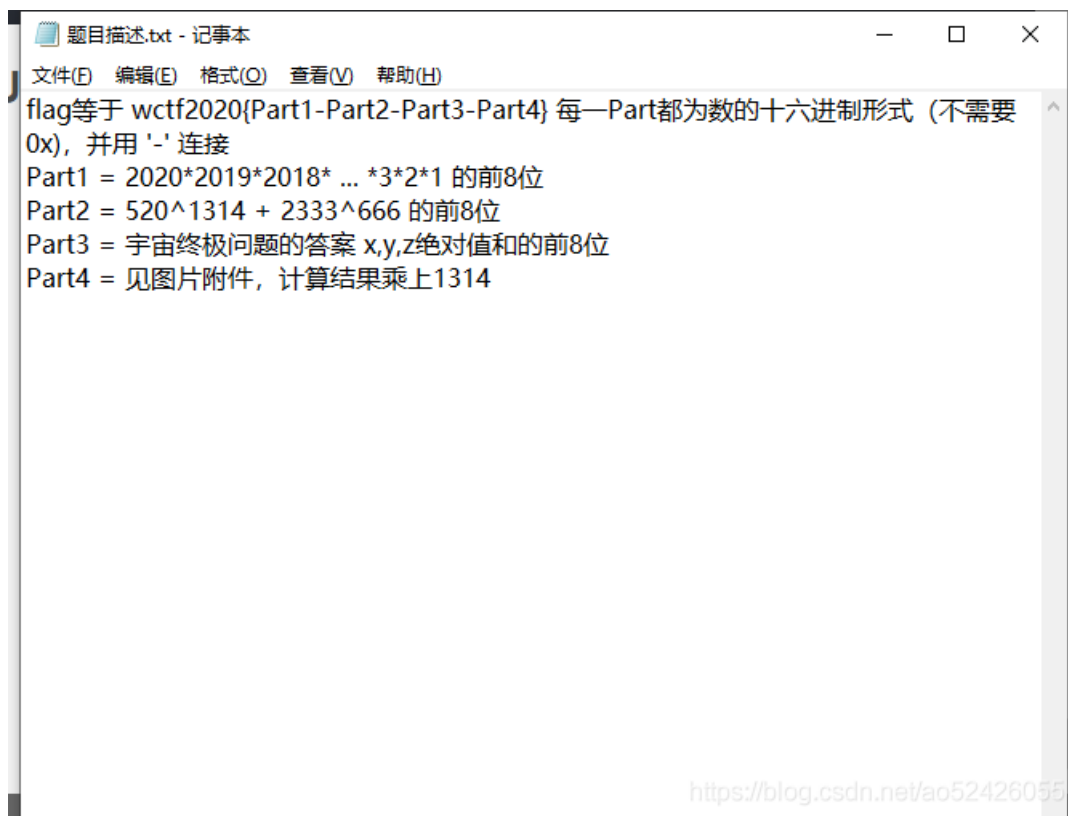
```
ax*x+bx+c=0
input a: 1
input b: 1
input c: -7943722218936282
X1: 89127561.0
X2: -89127562.0
```

在cmd框里最后得到

```
Last time: Kill it
x*x+x-7943722218936282=0
input x: 89127561
Great This is your FLAG
flag{en5oy_746831_89127561}
D:\ctfbisai\buumima\attachment(4)>
```

94.[WUSTCTF2020]大数计算

查看题目



```
题目描述.txt - 记事本
文件(F) 编辑(E) 格式(O) 查看(V) 帮助(H)
flag等于 wctf2020{Part1-Part2-Part3-Part4} 每一Part都为数的十六进制形式 (不需要
0x), 并用 '-' 连接
Part1 = 2020*2019*2018* ... *3*2*1 的前8位
Part2 = 520^1314 + 2333^666 的前8位
Part3 = 宇宙终极问题的答案 x,y,z绝对值和的前8位
Part4 = 见图片附件, 计算结果乘上1314

https://blog.csdn.net/ao52426055
```

$$\int_0^{22} 2x dx + 36 \leftarrow$$

我都不知道这个wp该怎么写。。。

脚本运行吧

```

#Part1 = 2020*2019*2018* ... *3*2*1 的前8位
p1=1
for i in range(1,2021):
    p1*=i
p1=str(p1)[0:8]
#print(p1)
part1=hex(int(p1))
print(part1)
#Part2 = 520^1314 + 2333^666 的前8位
p2=520**1314 + 2333**666
p2=str(p2)[0:8]
#print(p2)
part2=hex(int(p2))
print(part2)
#Part3 = 宇宙终极问题的答案 x,y,z绝对值和的前8位
p3=80538738812075974 + 80435758145817515 + 12602123297335631
p3=str(p3)[0:8]
#print(p3)
part3=hex(int(p3))
print(part3)
#Part4 = 见图片附件, 计算结果乘上1314
p4=(22**2+36)*1314
p4=str(p4)[0:8]
#print(p4)
part4=hex(int(p4))
print(part4)
print('wctf2020{' +str(part1)+'-' +str(part2)+'-' +str(part3)+'-' +str(part4)+'}')

```

运行得到

```

0x24d231f
0x403cfd3
0x108db5e
0xa6d10
wctf2020{0x24d231f-0x403cfd3-0x108db5e-0xa6d10}

```

去掉0x提交

95.RSA & what

查看题目

HUB1

```

7850954197182682868665082143048169854470772937668193987280464111669178108204847593142910289764982236612293950094
7406317370516262703761099353961775190544303927822758350460480825193108381890946761327758787454576107436442754996
6555519371913859875313577282243053150056274667798049694695703660313532933165449312949725581708965417273055582216
2959945876009759701248114962700808969770769460001027010302609905981814664472080547133915263137006813410939222403
1742817359903162412515518821648947682560619152118203496912034328769118130039968351541480926270045752587669180818
0257730351707673660380698973884642306898810000633684878715402823143549139850732982897459698089649561190746850698
1302994580802555823126968731492100282408981378228884925599576650679365733563675897845931190166240724338727445374
320059116684944557333068938514121465309188801778204904343486262030678343616985656417592987110066991343898089921
9579329897753233450934770193915434791427728636586218049874617231705308003720066269312729135764175698611068808404
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7

README.txt

素数生成算法太麻烦了，有没有取巧的方法呢？
诶，这里好像有个不错的想法哟。
看起来节约了不少时间呢，嘿嘿嘿.....
顺便问问，应该大家都知道base64吧，用来编码还是很方便的呢！

rsa.py

```

from Crypto.Util.number import bytes_to_long, getPrime
from random import randint
from gmpy2 import powmod

p = getPrime(2048)
q = getPrime(2048)
N = p*q
Phi = (p-1)*(q-1)
def get_enc_key(N,Phi):
    e = getPrime(N)
    if Phi % e == 0:
        return get_enc_key(N, Phi)
    else:
        return e
e1 = get_enc_key(randint(10, 12), Phi)
e2 = get_enc_key(randint(10, 12), Phi)

fr = open(r"./base64", "rb")#flag is in this file
f1 = open(r"./HUB1", "wb")
f2 = open(r"./HUB2", "wb")
base64 = fr.read(255)
f1.write("%d\n%d\n" % (N, e1))
f2.write("%d\n%d\n" % (N, e2))
while len(base64)>0:
    pt = bytes_to_long(base64)
    ct1 = powmod(pt, e1, N)
    ct2 = powmod(pt, e2, N)
    f1.write("\n%d" % ct1)
    f2.write("\n%d" % ct2)
    base64 = fr.read(255)
fr.close()
f1.close()
f2.close()

```

N一样 很容易想到共模攻击

```

from Crypto.Util.number import*
import base64

def egcd(a, b):
    if a == 0:
        return (b, 0, 1)
    else:
        g, y, x = egcd(b % a, a)
        return (g, x - (b // a) * y, y)

def CMA(n,e1,e2,c1,c2):
    s = egcd(e1, e2)
    s1 = s[1]
    s2 = s[2]
    if s1<0:
        s1 = - s1
        c1 = inverse(c1, n)
    elif s2<0:
        s2 = - s2
        c2 = inverse(c2, n)
    m = pow(c1,s1,n)*pow(c2,s2,n) % n
    return m

f1=open("HUB1")
f2=open("HUB2")
N=f1.readline()
N=f2.readline()
e1,e2=f1.readline(),f2.readline()
f1.readline()
f2.readline()
c1,c2=f1.readline(),f2.readline()
ans=b''
cnt=0
while len(c1)!=0:
    cnt+=1
    ans+=long_to_bytes(CMA(int(N),int(e1),int(e2),int(c1),int(c2)))
    #print(base64.b64decode(temp))
    c1,c2=f1.readline(),f2.readline()
temp=b''
M=b''
print(ans)
for i in ans:
    k=long_to_bytes(i)
    #print(i," ",end="")
    if k==b'\n':
        M+=base64.b64decode(temp)
        temp=b''
        continue
    temp+=k
print(M)

```

运行得到

```
b'VEhJUz==\nRkxBR3==\nSVN=\nSElEREVOLO==\nQ0FO\nWU9V\nRk1ORM==\nSVT=\nT1VUP4==\nRE8=\nwU9V\nS05PV9==\nQkFTRTY0P5
==\nwW91bmdD\nVEhJTKu=\nwU9V\nQVJF\nTk9U\nVEhBVE==\nRkFNSUXJQVI=\nv01US0==\nQkFTRTY0Lh==\nQmFzZTY0\naX0=\nYW==\n
Z3JvdXA=\nb2b=\nc21taWxhcn==\nYm1uYXJ5LXRvLXR1eHR=\nZW5jb2Rpbme=\nc2NoZW1lc0==\ndGhhdD==\ncmVwcmVzZW50\nYm1uYXJ5
\nZGF0YW==\naW5=\nYW6=\nQVNDU1=\nc3RyaW5n\nZm9ybWF0\nYnk=\ndHJhbnNsYXRpbmd=\naXS=\naW50b1==\nYT==\ncmFkaXgtNjQ=
\ncmVwcmVzZW50YXRpb24u\nVGh1\ndGVybc==\nQmFzZTY0\nb3JpZ21uYXRlc8==\nZnJvd==\nY==\nc3B1Y21maWN=\nTU1NRT==\nY29u
dGVudI==\ndHJhbnNmZSI=\nZW5jb2Rpbmcu\nVGh1\ncGFydG1jdWxhct==\nc2V0\nb2b=\nNjR=\nY2hhcmFjdGVyc5==\nY2hvc2Vu\ndG+
\ncmVwcmVzZW50\ndGh1\nNjQ=\ncGxhY2UtZmFsdWVz\nZm9y\ndGh1\nYmFzZd==\ndmFyaWVz\nYmV0d2V1bT==\naW1wbGVtZW50YXRpb25z
Lp==\nVGh1\nZ2VuZlZhbI==\nc3RyYXR1Z3n=\naX0=\ndG9=\nY2hvb3N1\nNjR=\nY2hhcmFjdGVyc5==\ndGhhdA==\nYXJ1\nYm90aN==\n
bWVtYmVyc5==\nb2a=\nYS==\nc3Vic2V0\nY29tbW9u\ndG8=\nbW9zdM==\nZW5jb2RpbmdzLA==\nYW5k\nYWxz8==\ncHJpbmRhYm1lg==
\nVGhpc9==\nY29tYm1uYXRpb25=\nbGVhdmVz\ndGh1\nZGF0YW==\ndW5saWt1bHk=\ndG/= \nYmV=\nbW9kaWZpZS=\naW5=\ndHJhbnNpdE
==\ndGhyb3VnaN==\naW5mb3JtYXRpb26=\nc3lzdGVtcyw=\nc3VjaN==\nYXM=\nRS1tYW1sLD==\ndGhhdA==\nd2VyZQ==\ndHJhZG10aw9u
YWxseQ==\nbm90\nOC1iaXQ=\nY2x1YW4uWzFd\nRm9y\nZXhhbXBsZSsw=\nTU1NRSdz\nQmFzZTY0\naW1wbGVtZW50YXRpb24=\ndXNlcw==\n
QahDwiw=\nYahDeiw=\nYW5k\nMKhDOQ==\nZm9y\ndGh1\nZmlyc3Q=\nNjI=\ndmFsdWVzLg==\nT3RoZXI=\ndmFyaWF0aW9ucw==\nc2hhcm
U=\ndGhpcw==\ncHJvcGVydHk=\nYnV0\nZGlMmZmVy\naW4=\ndGh1\nc3ltYm9scw==\nY2hvc2Vu\nZm9y\ndGh1\nbGFzdA==\ndHdv\ndmFs
dWVzOw==\nYW4=\nZXhhbXBsZQ==\naXM=\nVVRGLTcu'
```

```
b"THISFLAGISHIDDEN.CANYOUFINDITOUT?DOYOUKNOWBASE64?YoungCTHINKYOUARENOTTHATFAMILIARWITHBASE64.Base64isagroupofsi
milarbinary-to-textencodingschemes that represent binary data in an ASCII string format by translating it into a radix-64 repres
entation.The term Base64 originates from a specific MIME content transfer encoding. The particular set of 64 characters chosen to r
epresent the 64 place-values for the base varies between implementations. The general strategy is to choose 64 characters that are b
oth members of a subset common to most encodings, and also printable. This combination leaves the data unlikely to be modified in tra
nsit through information systems, such as E-mail, that were traditionally not 8-bit clean.[1] For example, MIME's Base64 implement
ation uses A\xa8CZ, a\xa8Cz, and \0\xa8C9 for the first 62 values. Other variations share this property but differ in the symbols cho
sen for the last two values; an example is"
```

再看第三个提示，就知道是base隐写
我白嫖了一个隐写脚本


```

from Crypto.Util.number import*
import base64
c = b'VEhJUz==\nRkxBR3==\nSVN=\nSEIEREVOLo==\nQ0F0\nWU9V\nRk1ORM==\nSVT=\nT1VUP4==\nRE8=\nWU9V\nS05PV9==\nQkFTRTY0P5==\nWw91bmdD\nVEhJTKu=\nWU9V\nQVJF\nTk9U\nVEHBE==\nRkFNSUxJQVI=\nV01US0==\nQkFTRTY0Lh==\nQmFzZTY0\naX0=\nYW==\nZ3JvdXA=\nb2b=\nc21taWxhcnc==\nYm1uYXJ5LXRvLXRleHR=\nZW5jb2Rpbme=\nc2NoZW1lc0==\ndGhhdD==\ncmVwcmVzZW50\nYm1uYXJ5\nZGF0YW==\naW5=\nYW6=\nQVNDUSU1=\nc3RyaW5n\nZm9ybWw0\nYnk=\ndHJhbnNsYXRpbmd=\naXS=\naW50b1==\nYT==\ncmFkaXgtNjQ=\ncmVwcmVzZW50YXRpb24u\nVGhl\ndGVybc==\nQmFzZTY0\nb3JpZ21uYXRlc8==\nZnJvbd==\nYY==\nc3B1Y2lmaWN=\nTU1NRT==\nY29udGVudI==\ndHJhbnNmZXI=\nZW5jb2Rpbmcu\nVGhl\ncGFydG1jdWxhct==\nc2V0\nb2b=\nNjR=\nY2hhcmFjdGVyc5==\nY2hvc2Vu\nnG+=\ncmVwcmVzZW50\ndGhl\nNjQ=\ncGxhY2UtdmFsdWVz\nZm9y\ndGhl\nYmFzZd==\ndmFyaWVz\nYmV0d2V1bt==\naW1wbGVtZW50YXRpb25zLp==\nVGhl\nZ2VuZUxhbnI==\nc3RyYXR1Z3n=\naX0=\ndG9=\nY2hvbnN1\nNjR=\nY2hhcmFjdGVyc5==\ndGhhdA==\nYXJ1\nYm90aA==\nbwVtYmVyc5==\nb2a=\nYS==\nc3Vic2V0\nY29tbW9u\ndG8=\nbw9zdM==\nZW5jb2RpbmdzLA==\nYW5k\nYmVzbn8=\ncHJpbmRhYmX1Lg==\nVGhpc9==\nY29tYm1uYXRpb25=\nbGVhdmVz\ndGhl\nZGF0YW==\ndW5sawt1bHk=\ndG/= \nYmV=\nbW9kawZpZWS=\naW5=\ndHJhbnNpdE==\ndGhyb3VnaN==\naW5mb3JtYXRpb26=\nc3lzdGVtcyw=\nc3VjaN==\nYXM=\nRS1tYW1sLD==\ndGhhdA==\nd2VyZQ==\ndHJhZGI0aW9uYXwseQ==\nbm90\nOC1iaXQ=\nY2x1YW4uWzFd\nRm9y\nZXhhbXBsZSsw=\nTU1NRSdz\nQmFzZTY0\naW1wbGVtZW50YXRpb24=\ndXN1cw==\nQahDWiw=\nYahDeiw=\nYW5k\nMKhD0Q==\nZm9y\ndGhl\nZmlyc3Q=\nNjI=\ndmFsdWVzLg==\nT3RoZXI=\ndmFyaWV0aW9ucw==\nc2hhcmU=\ndGhpcw==\ncHJvcGVydHk=\nYnV0\nZGlmZmVy\naW4=\ndGhl\nc3ltYm9scw==\nY2hvc2Vu\nZm9y\ndGhl\nbGFzdA==\ndHdv\nZmFsdWVzOw==\nYW4=\nZXhhbXBsZQ==\naXM=\nVVRGLTcu'

def get_base64_diff_value(s1, s2):
    base64chars = b'ABCDEFGHIJKLMNOPQRSTUVWXYZabcdefghijklmnopqrstuvwxyz0123456789+/'
    res = 0
    for i in range(len(s2)):
        if s1[i] != s2[i]:
            return abs(base64chars.index(s1[i]) - base64chars.index(s2[i]))
    return res

def solve_stego():
    line=b''
    bin_str=''
    for i in c:
        k=long_to_bytes(i)
        if k==b'\n':
            steg_line = line
            norm_line = base64.b64encode(base64.b64decode(line))
            diff = get_base64_diff_value(steg_line, norm_line)
            #print(diff)
            pads_num = steg_line.count(b'=')
            if diff:
                bin_str += bin(diff)[2:].zfill(pads_num * 2)
            else:
                bin_str += '0' * pads_num * 2
            print(goflag(bin_str))
            line=b''
            continue
        line+=k

def goflag(bin_str):
    res_str = ''
    for i in range(0, len(bin_str), 8):
        res_str += chr(int(bin_str[i:i + 8], 2))
    return res_str

if __name__ == '__main__':
    solve_stego()

```

运行得到

flag包裹提交即可 注意（共模攻击py要和H1，2放在一个文件夹）

96.[MRCTF2020]babyRSA

[查看题目一个加密脚本](#)

```

'''
import sympy
import random
from gmpy2 import gcd, invert
from Crypto.Util.number import getPrime, isPrime, getRandomNBitInteger, bytes_to_long, long_to_bytes
from z3 import *
flag = b"MRCTF{xxxx}"
base = 65537

def GCD(A):
    B = 1
    for i in range(1, len(A)):
        B = gcd(A[i-1], A[i])
    return B

def gen_p():
    P = [0 for i in range(17)]
    P[0] = getPrime(128)
    for i in range(1, 17):
        P[i] = sympy.nextprime(P[i-1])
    print("P_p :", P[9])
    n = 1
    for i in range(17):
        n *= P[i]
    p = getPrime(1024)
    factor = pow(p, base, n)
    print("P_factor :", factor)
    return sympy.nextprime(p)

def gen_q():
    sub_Q = getPrime(1024)
    Q_1 = getPrime(1024)
    Q_2 = getPrime(1024)
    Q = sub_Q ** Q_2 % Q_1
    print("Q_1: ", Q_1)
    print("Q_2: ", Q_2)
    print("sub_Q: ", sub_Q)
    return sympy.nextprime(Q)

if __name__ == "__main__":
    _E = base
    _P = gen_p()
    _Q = gen_q()
    assert (gcd(_E, (_P - 1) * (_Q - 1)) == 1)
    _M = bytes_to_long(flag)
    _C = pow(_M, _E, _P * _Q)
    print("Ciphertext = ", _C)
'''

```

拿到加密脚本以后,发现p,q的生成方式不同,其中,q的生成方式比较简单.三数相乘以后取结果的下一个素数.p的求解则需要先求出n的欧拉函数值,然后求出base对应的逆元然后求解p,p求出来以后就可以解密密文了.

脚本如下

```

P_p=206027926847308612719677572554991143421
P_factor=2136717427659089807871165799762896005958647045741344691731117909652336299095138847041584469464099104757
2758434264184859785894220915111462730628639339025970023969886948746908088126718280306248804346913825278638182264
6126962323295676431679988602406971858136496624861228526070581338082202663895710929460596143281673761666804565161
4359639576550120110519361805365814884990595179463086501353004286724868196452799696935190394078929416727843628686
5324363272792827969858817769417179725464486455416284869621076368119727975813081172370015461828076412339631233003
2986093579531909363210692564988076206283296967165522152288770019720928264542910922693728918198338839
Q_1=103766439849465588084625049495793857634556517064563488433148224524638105971161051763127718438062862548184814
7476012994940528136628514597401274995577853987144819094616319960200483157901679676999329679744844812098796641730
09585231469785141628982021847883945871201430155071257803163523612863113967495969578605521
Q_2=151010734276916939790591461278981486442548035032350797306496105136358723586953123484087860176438629843688462
6716817775136529475553256074148585145660535132430836278106860848902611206411619876144351148875654918661205078445
66210561620503961205851409386041194326728437073995372322433035153519757017396063066469743
sub_Q=1689925297935933157578959951014302419949536383309193148001305368098018249711120395725623894495843506439243
9198480097819370779590995647299263100429047927352511695946185622726223260008917695081072947505826033217762696128
6009876630340945093629959302803189668904123890991069113826241497783666995751391361028949651
Ciphertext=17091872405163671414608621877494510476440948857917616735746743308408427921897950499683941222168544917
5792264765643090858705999707048867422033084787181183672454190766698304237621641156182664006073430701345879492502
5684062804589439843027290282034999617915124231838524593607080377300985152179828199569474241678651559771763395596
6971402060725376881297901264720539873915382800070822030063480291257296502076613623719361967895626584587783125335
0593885895964454123357865434092590196395798004763911417003393657006025043890613059137790418211162223656750702271
1176457301476543461600524993045300728432815672077399879668276471832
import gmpy2
import Crypto
import sympy
base = 65537
q=sympy.nextprime(gmpy2.powmod(sub_Q,Q_2,Q_1))
print(q)
P=[]
for i in range(9):
    P_p=sympy.prevprime(P_p)
P.append(P_p)
for i in range(1,17):
    P.append(sympy.nextprime(P[i-1]))
n=1
phi=1
for i in range(17):
    n*=P[i]
    phi*=(P[i]-1)
based=gmpy2.invert(base, phi)
_p=gmpy2.powmod(P_factor, based, n)
p=sympy.nextprime(_p)
print(p)
d=gmpy2.invert(base, (p-1)*(q-1))
print(d)
plaintext=gmpy2.powmod(Ciphertext, d, p*q)
import binascii
print(binascii.unhexlify(hex(plaintext)[2:]))

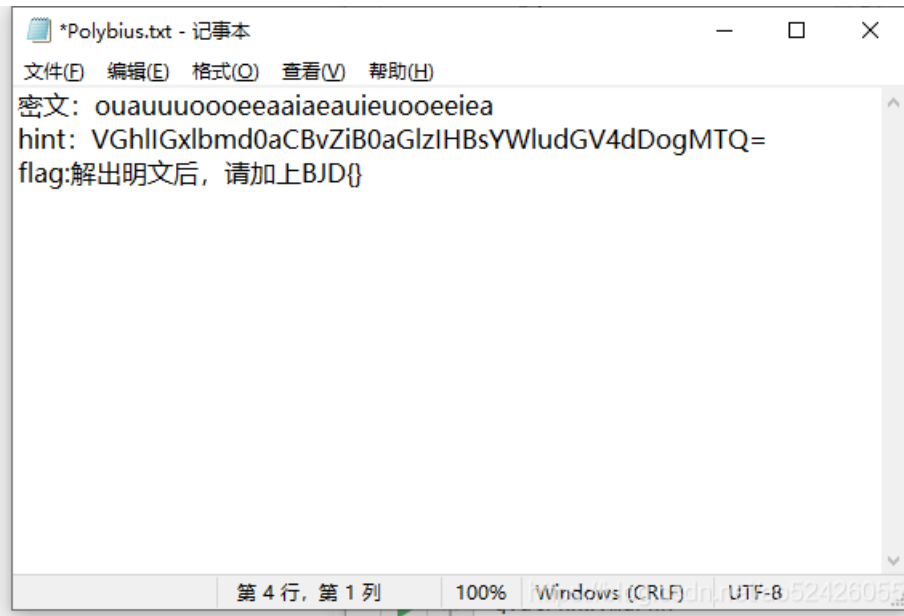
```

我喜欢输出pq
运行得到

```
951706537140816870887605854409067687004194597677433375733684286450760708180919337087074776999321825692511110026
0761958233280546585624501259121060195932474781731613458132842656517609786144352755126076860272047457230913808406
105832246663969943550533958139118721153456230616182820319799156494938586844573835221
1607353802641185641618355367827829241600056206316799298554452902073519458632582820882652022328622021806688449472
0580626132371394581887285230324859035563266588690092852053342177472159093548577323461955818151303338564271170620
5607543347313747616062185115981201425568780146693758544521883683953378438266703113683
3511955744337428151031979941361296216306473937469271350521070131816991178710055704862733703600712476342357982869
3369248039257489113761845745849187589419531504626751267311173511163470030839119103143532372091977653267149268524
3496013965439642965417782657024324810813460943009803523746219386004686158627347562391960350225884309848872664977
7621733372184366568775297877185960241115496391276974016445360524169770129867843553985769813282863097708315136140
0659334142547950320704518925659970267166574056694840174902568549981607232865248331879281680147343835445705493729
47304031812767519366827995325795367403195819253459303793
b'MRCTF{sti11_@_b@by_qu3st10n}'
```

97.[BJDCTF2020]Polybius

查看题目



hint用base64解密出来The length of this plaintext: 14

,结合题目polybius 猜测这是波利比奥斯方阵密码。

但是a,e,o,i,u这五个字符的代表顺序却不知道,因此可能有54321种情况,在结合刚才所说的i,j同时占一个位置,所以情况数要再乘上2,将这些情况全部都打印出来,然后去找有真实语义的句子就可以了。

```

import itertools
s="aeoiu"
sumresult=[]
numsumresult=[]
ciper="ouuuuoooeaaiaeauieuooeeiea"

for i in itertools.permutations(s,5):#找出所有全排列
    sumresult.append("".join(i))
for i in sumresult:
    temp=""
    for j in ciper:
        temp+=str(i.index(j)+1)
    numsumresult.append(temp)
for i in numsumresult:
    ans_=""
    for j in range(0, len(i),2):
        xx=(int(i[j])-1)*5+int(i[j+1])+96
        if xx>ord('i'):
            xx+=1
        ans_+=chr(xx)
    print(ans_)

```

运行得到

```

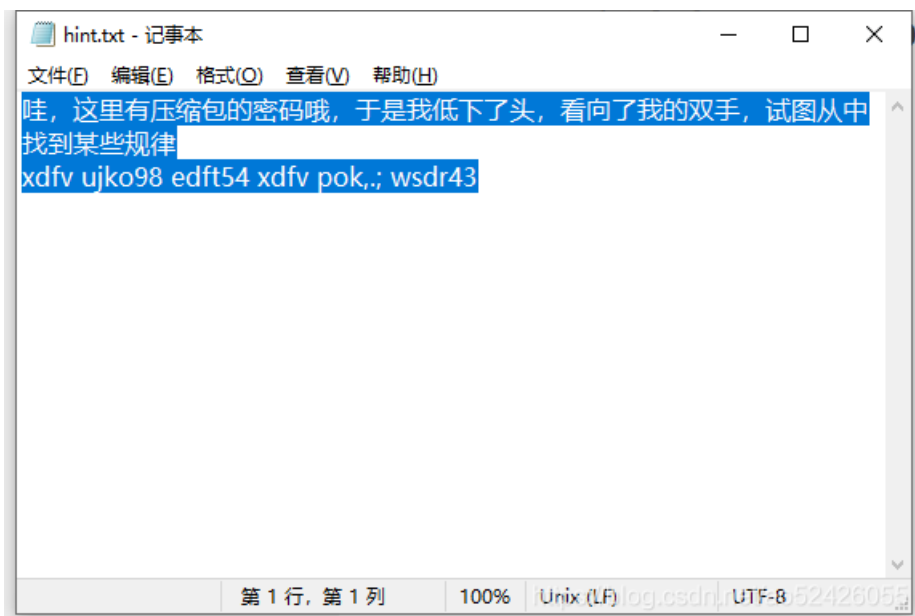
flagispolybius
flagkxoplubkyx
fqaghousqxbhpo
fvaghyxvsbhop

```

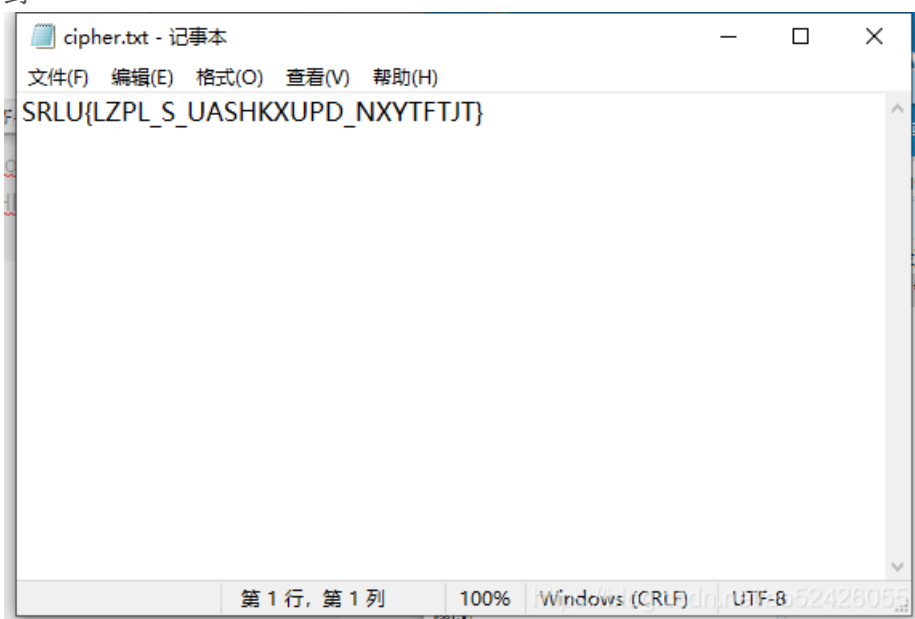
很多的结果里面找到flagispolybius

98.[ACTF新生赛2020]crypto-classic1

查看题目



低下头看，就是看键盘，比如xdfv包围的就是c
按照这个规律把键盘密码解出来就是circle
这就是压缩包密码，得到



脚本如下

```

s = 'ABCDEFGHIJKLMNOPQRSTUVWXYZ'
s1 = 'ACTF'
s2 = 'SRLU'
key = ''
for i in range(len(s1)):
    key += s[(s.find(s2[i]) - s.find(s1[i])) % 26]
print(key)

# 解密
cipher = 'SRLU{LZPL_S_UASHKXUPD_NXYTFTJT}'
key = 'SPSP'
# decode
flag = ''
for i in range(0, len(cipher)):
    flag += s[(s.find(cipher[i]) + 26 - s.find(key[i % len(key)])) % 26]
print(flag)

```

运行得到

```

SPSP
ACTFHZEDQKAKZOXVPLZDIKVIGENEREH

```

按照原题顺序把{}这些换回来，ACTF{WHAT_A_CLASSICAL_VIGENERE}提交发现不对
然后就越走越远了，我看别的大佬的wp，还有说题出错的直接改题目的等等

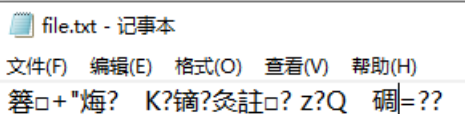
我发现他们的我提交也不行啊

最后想起一个大小写转换

flag{what_a_classical_vigenerE}提交这个成功了，顿时觉得自己好沙雕

99.EasyProgram

查看题目



file.txt - 记事本

文件(F) 编辑(E) 格式(O) 查看(V) 帮助(H)

箬口+"悔? K?摘?灸註口? z?Q 碯|=??

附件.txt


```

get buf unsign s[256]
get buf t[256]
we have key:whoami
we have flag:????????????????????????????????????????????????????????????

for i:0 to 256
  set s[i]:i

for i:0 to 256
  set t[i]:key[(i)mod(key.lenth)]

for i:0 to 256
  set j:(j+s[i]+t[i])mod(256)
  swap:s[i],s[j]

for m:0 to 38
  set i:(i + 1)mod(256)
  set j:(j + S[i])mod(256)
  swap:s[i],s[j]
  set x:(s[i] + (s[j]mod(256))mod(256))
  set flag[m]:flag[m]^s[x]

fprintf flagx to file

```

代码是一系列运算后，最后进行异或。解题思路就是，还原源代码，再进行一次异或就可以了

```

#读文件方法一:
def filehex(file):
    fhex=[]
    f = open(file,'rb')
    ff = f.read().hex()

    for i in range(0,len(str(ff)),2):
        fhex.append(int(str(ff)[i:i+2],16))
    return fhex

flagx2=filehex('file.txt')
print(flagx2)

#读文件方法二: 使用010edit 读取'file.txt'文件16进制
flagx=[0x00,0xBA,0x8F,0x11,0x2B,0x22,0x9F,0x51,0xA1,0x2F,0xAB,0xB7,0x4B,0xD7,0x3F,0xEF,0xE1,0xB5,0x13,0xBE,0xC4,
0xD4,0x5D,0x03,0xD9,0x00,0x7A,0xCA,0x1D,0x51,0xA4,0x73,0xB5,0xEF,0x3D,0x9B,0x31,0xB3]

s=[]
t=[]
key='whoami'
j=0
for i in range(0,256):
    s.append(i)

for i in range(0,256):
    t.append(key[i % len(key)])

for i in range(0,256):
    j=(j+int(s[i])+int(ord(t[i])))%256
    s[i], s[j] = s[j], s[i]

i=0;j=0;x=0
for m in range(0,38):
    i=(i + 1)%256
    j=(j + s[i])%256
    s[i],s[j] = s[j],s[i]
    x=(s[i] + (s[j]%256))%256
    flagx[m] = flagx[m] ^ s[x]
    flagx2[m]=flagx2[m]^s[x]

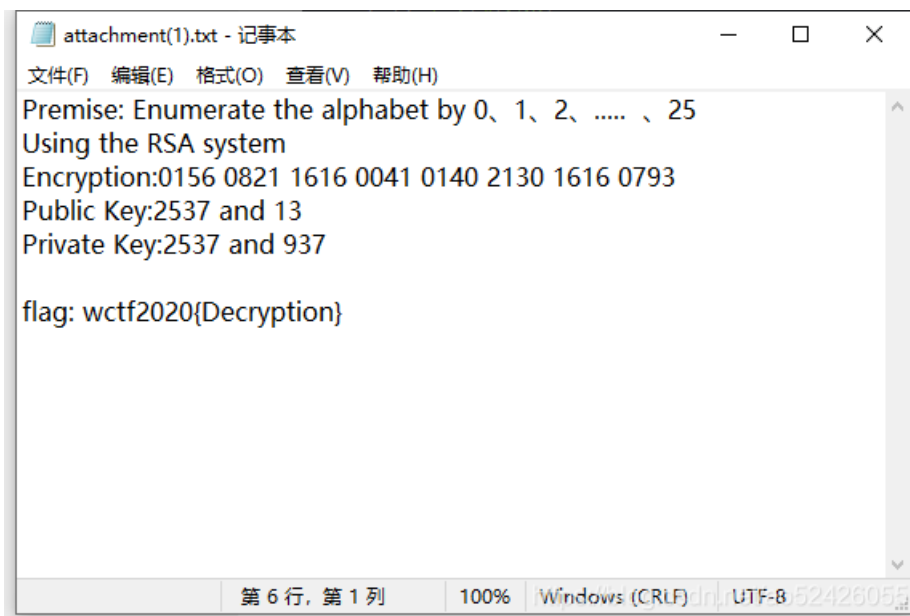
print(''.join(chr(flagx[i]) for i in range(0,38)))
print(''.join(chr(flagx2[i]) for i in range(0,38)))

```

得到flag{f238yu28323uf28u2yef2ud8uf289euf}

100.[WUSTCTF2020]情书

查看题目



可以看到RSA加密系统
2537分解得到43和59

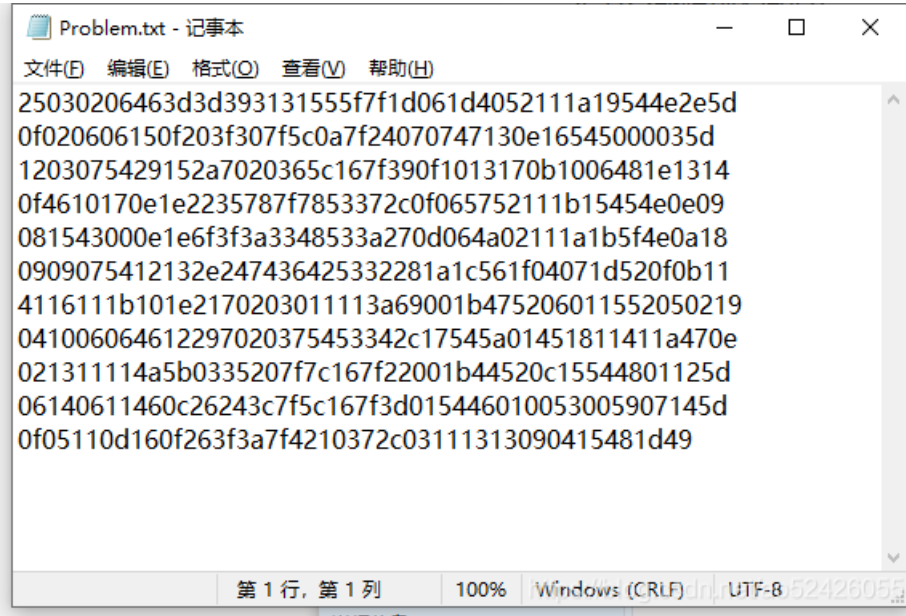
```
a = "abcdefghijklmnopqrstuvwxyz"
c = "0156 0821 1616 0041 0140 2130 1616 0793".split(" ")
N = 2537
e = 13
d = 937
p = 43
q = 59
phi_N = (p-1)*(q-1)

m = "".join(a[pow(int(i),d,N)] for i in c)
print (m)
```

运行得到iloveyou

101.[AFCTF2018]你听过一次一密么？

查看题目



加密方法

一次一密的加密算法就是异或加密。将明文分组，与密钥按位异或即可。

我这里出的题目是用同一个密钥去加密多条明文，当密文条数较多时就很容易被攻击，例如Many Time Pad。

Many Time Pad攻击

这个攻击的原理是 $c1 \oplus c2 = m1 \oplus m2$ ，而通过 $m1 \oplus m2$ 可以分析出 $m1$ 和 $m2$ ，因此 $m1$ 与 $m2$ 不再安全。

代码搬运自GitHub。

```
#!/usr/bin/python
## OTP - Recovering the private key from a set of messages that were encrypted w/ the same private key (Many time
pad attack) - crypto100-many_time_secret @ alexctf 2017
# Original code by jwomers: https://github.com/Jwomers/many-time-pad-attack/blob/master/attack.py

import string
import collections
import sets, sys

# 11 unknown ciphertexts (in hex format), all encrypted with the same key

c1='25030206463d3d393131555f7f1d061d4052111a19544e2e5d'
c2='0f020606150f203f307f5c0a7f24070747130e16545000035d'
c3='1203075429152a7020365c167f390f1013170b1006481e1314'
c4='0f4610170e1e2235787f7853372c0f065752111b15454e0e09'
c5='081543000e1e6f3f3a3348533a270d064a02111a1b5f4e0a18'
c6='0909075412132e247436425332281a1c561f04071d520f0b11'
c7='4116111b101e2170203011113a69001b475206011552050219'
c8='041006064612297020375453342c17545a01451811411a470e'
c9='021311114a5b0335207f7c167f22001b44520c15544801125d'
c10='06140611460c26243c7f5c167f3d015446010053005907145d'
c11='0f05110d160f263f3a7f4210372c03111313090415481d49'
ciphers = [c1, c2, c3, c4, c5, c6, c7, c8, c9, c10, c11]
# The target ciphertext we want to crack
#target_cipher = "0529242a631234122d2b36697f13272c207f2021283a6b0c7908"

# XORs two string
def strxor(a, b): # xor two strings (trims the longer input)
    return "".join([chr(ord(x) ^ ord(y)) for (x, y) in zip(a, b)])

def target_cipher(target_ciphers):
```

```

def target_fix(target_cipher):
    # To store the final key
    final_key = [None]*150
    # To store the positions we know are broken
    known_key_positions = set()

    # For each ciphertext
    for current_index, ciphertext in enumerate(ciphers):
        counter = collections.Counter()
        # for each other ciphertext
        for index, ciphertext2 in enumerate(ciphers):
            if current_index != index: # don't xor a ciphertext with itself
                for index_of_char, char in enumerate(strxor(ciphertext.decode('hex'), ciphertext2.decode('hex'))):
                    # Xor the two ciphertexts
                    # If a character in the xored result is an alphanumeric character, it means there was probably
                    # a space character in one of the plaintexts (we don't know which one)
                    if char in string.printable and char.isalpha(): counter[index_of_char] += 1 # Increment the counter at this index
                known_space_indexes = []

                # Loop through all positions where a space character was possible in the current_index cipher
                for ind, val in counter.items():
                    # If a space was found at least 7 times at this index out of the 9 possible XORS, then the space character was likely from the current_index cipher!
                    if val >= 7: known_space_indexes.append(ind)
                #print known_space_indexes # Shows all the positions where we now know the key!

                # Now Xor the current_index with spaces, and at the known_space_indexes positions we get the key back!
                xor_with_spaces = strxor(ciphertext.decode('hex'),' '*150)
                for index in known_space_indexes:
                    # Store the key's value at the correct position
                    final_key[index] = xor_with_spaces[index].encode('hex')
                    # Record that we know the key at this position
                    known_key_positions.add(index)

        # Construct a hex key from the currently known key, adding in '00' hex chars where we do not know (to make a complete hex string)
        final_key_hex = ''.join([val if val is not None else '00' for val in final_key])
        # Xor the currently known key with the target cipher
        output = strxor(target_cipher.decode('hex'),final_key_hex.decode('hex'))

        print "Fix this sentence:"
        print ''.join([char if index in known_key_positions else '*' for index, char in enumerate(output)]+"\n")

    # WAIT.. MANUAL STEP HERE
    # This output are printing a * if that character is not known yet
    # fix the missing characters like this: "Let*M*k*ow if *o{*a" = "cure, Let Me know if you a"
    # if is too hard, change the target_cipher to another one and try again
    # and we have our key to fix the entire text!

    #sys.exit(0) #comment and continue if u got a good key

    target_plaintext = "cure, Let Me know if you a"
    print "Fixed:"
    print target_plaintext+"\n"

    key = strxor(target_cipher.decode('hex'),target_plaintext)

    print "Decrypted msg:"
    for cipher in ciphers:

```

```

print strxor(cipher.decode('hex'),key)

print "\nPrivate key recovered: "+key+"\n"

for i in ciphers:
    target_fix(i)

```

运行得到

```
'afctf{OPT_1s_Int3rest1ng}'
```

102.[BJDCTF2020]编码与调制

查看题目 曼彻斯特



6进制024A447B4469664D616E63686573746572636F64657D

通过解码16进制即可得到flag{DifManchestercode}

103.[网鼎杯 2020 青龙组]you_raise_me_up

查看题目

```

#!/usr/bin/env python
# -*- coding: utf-8 -*-
from Crypto.Util.number import *
import random

n = 2 ** 512
m = random.randint(2, n-1) | 1
c = pow(m, bytes_to_long(flag), n)
print 'm = ' + str(m)
print 'c = ' + str(c)

# m = 3911907091245274289594896625652740393183059521729368594038550795814027709868903084690847354512078853863189
86881041563704825943945069343345307381099559075
# c = 6665851394203214245856789450723658632520816791621796775909766895233000234023642878786025644953797995373211
308485605397024123180085924117610802485972584499

```

mbytes_to_long(flag) %n=c的flag值.


```
m = 391190709124527428959489662565274039318305952172936859403855079581402770986890308469084735451207885386318986
881041563704825943945069343345307381099559075
c = 666585139420321424585678945072365863252081679162179677590976689523300023402364287878602564495379799537321130
8485605397024123180085924117610802485972584499
n = 2 ** 512
import sympy
flag=sympy.discrete_log(2**512,c,m)
import binascii
print(binascii.unhexlify(hex(flag)[2:]))
```

运行得到

```
b'flag{5f95ca93-1594-762d-ed0b-a9139692cb4a}'
```

104.[AFCTF2018]BASE

查看题目

 flag_encode.txt 21.73 MB 3.29 MB 文本文档




拿到附件是一个flag加密后的文本，其实是flag经过很多次base系列编码后的一个内容。文件有20+MB，文本编辑器发现是一行编码后的内容。

这题主要是三种编码，base64/base32/base16。base16也就是16进制。

写了一个很简单的脚本

```
import base64
file = open("flag_encode.txt", 'r')
file2 = open("flag.txt", 'w')
base = file.read()
while(1):
    try:
        base = base64.b32decode(base).decode()
    except:
        try:
            base = base64.b64decode(base).decode()
        except:
            try:
                base = base64.b16decode(base).decode()
            except:
                print("完成!")
                file2.write(base)
                break
```

运行得到

 Basebaoli.py	2020/12/1 8:56	JetBrains PyCha
 flag.txt	2020/12/1 8:55	文本文档
 flag_encode.txt	2018/4/9 12:14	文本文档

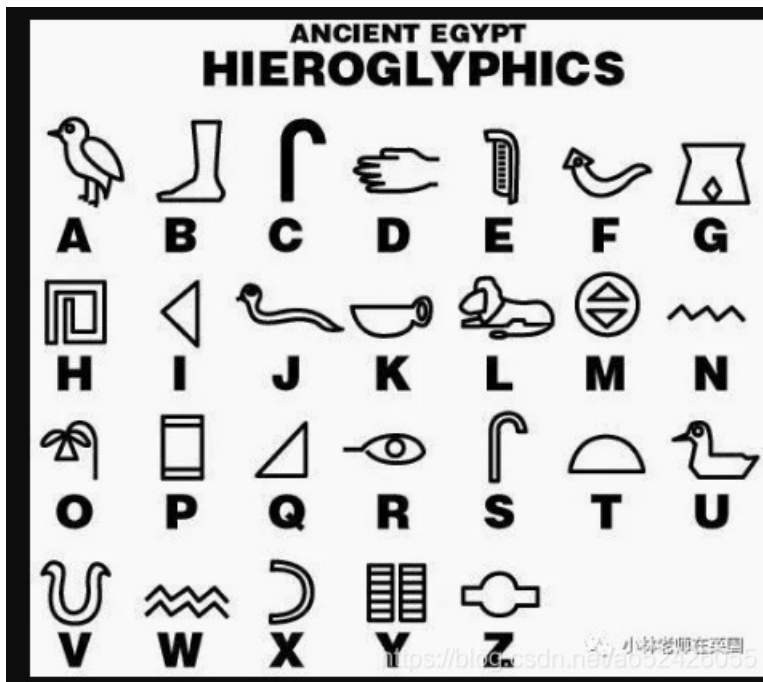
```
flag.txt - 记事本
文件(F) 编辑(E) 格式(O) 查看(V) 帮助(H)
afctf{U_5h0u1d_Us3_T0015}
```

105.[NPUCTF2020]Classical Cipe

放在<https://quipqiup.com/>上解，试出来压缩包密码：the_key_is_atdash
解出密码得到图片



是象形文字和猪圈密码的结合



flag{classicalcode}

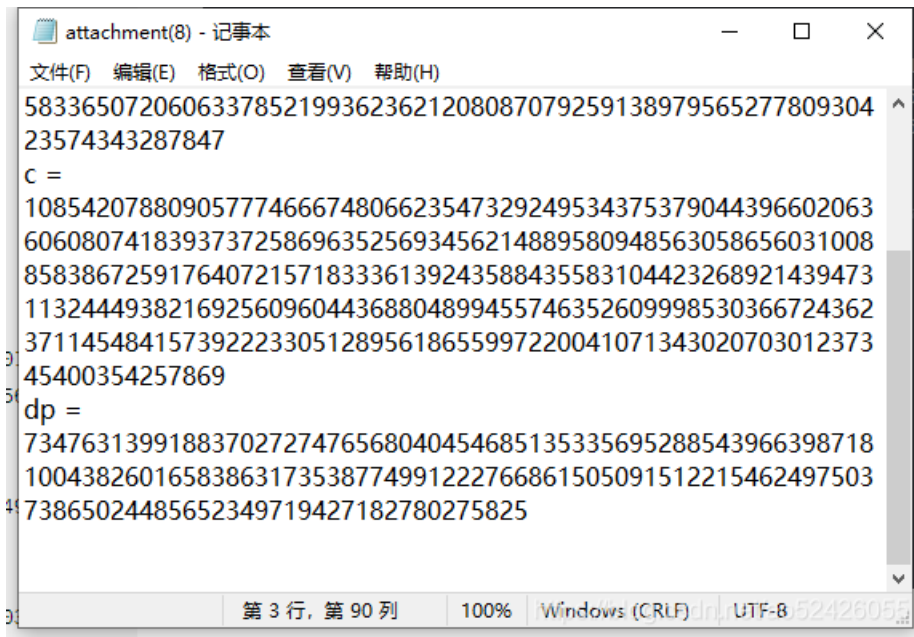
解压得到一张图，变形猪圈很容易看出来（那个鸟和狗是啥真没见过...），根据题目名称里的classical还是很容易顺出来的，答案是 classicalcode，flag{classicalcode}

106.[WUSTCTF2020]dp_leaking_1s_very_d@angerous

这个一看到题目我相信有的大佬就直接把题目flag包裹提交了哈哈

没错题目就是答案

查看题目



```
attachment(8) - 记事本
文件(F) 编辑(E) 格式(O) 查看(V) 帮助(H)
58336507206063378521993623621208087079259138979565277809304
23574343287847
c =
10854207880905777466674806623547329249534375379044396602063
60608074183937372586963525693456214889580948563058656031008
85838672591764072157183336139243588435583104423268921439473
11324449382169256096044368804899455746352609998530366724362
37114548415739222330512895618655997220041071343020703012373
45400354257869
dp =
73476313991883702727476568040454685135335695288543966398718
10043826016583863173538774991222766861505091512215462497503
73865024485652349719427182780275825
```

bp泄露，把以前的脚本带一下数值就可以了

```
import gmpy2 as gp

e = 65537
n = 156808343598578774957375696815188980682166740609302831099696492068246337198792510898818496239166339015207305
1021014316342831685444929845865667999964711502523821441482572367072472675061656708775063702531276953141639870840
76462560095456635833650720606337852199362362120808707925913897956527780930423574343287847
dp = 73476313991883702727476568040454685135335695288543966398718100438260165838631735387749912227668615050915122
1546249750373865024485652349719427182780275825
c = 108542078809057774666748066235473292495343753790443966020636060807418393737258696352569345621488958094856305
8656031008858386725917640721571833361392435884355831044232689214394731132444938216925609604436880489945574635260
99985303667243623711454841573922233051289561865599722004107134302070301237345400354257869

for i in range(1, e): # 在范围(1, e)之间进行遍历
    if (dp * e - 1) % i == 0:
        if n % (((dp * e - 1) // i) + 1) == 0: # 存在p, 使得n能被p整除
            p = ((dp * e - 1) // i) + 1
            q = n // (((dp * e - 1) // i) + 1)
            phi = (q - 1) * (p - 1) # 欧拉定理
            d = gp.invert(e, phi) # 求模逆
            m = pow(c, d, n) # 快速求幂取模运算

print(m) # 10进制明文
print('-----')
print(hex(m)[2:]) # 16进制明文
print('-----')
print(bytes.fromhex(hex(m)[2:])) # 16进制转文本
```

运行得到

```
3891178547072937865927641698099974931243767865417536099629795080648874795850439816682933678973
-----
77637466323032307b64705f6c65616b696e675f31735f766572795f6440616e6765726f75737d
-----
b'wctf2020{dp_leaking_1s_very_d@angerous}'
```

脚本这些东西该存就存一下

107.[UTCTF2020]basic-crypto

题目是一堆二进制

```
str='01010101 01101000 00101101 01101111 01101000 00101100 00100000 01101100 01101111 01101111 01101011 01110011
00100000 01101100 01101001 01101011 01100101 00100000 01110111 01100101 00100000 01101000 01100001 01110110 011
00101 00100000 01100001 01101110 01101111 01110100 01101000 01100101 01110010 00100000 01100010 01101100 0110111
1 01100011 01101011 00100000 01101111 01100110 00100000 01110100 01100101 01111000 01110100 00101100 00100000 01
110111 01101001 01110100 01101000 00100000 01110011 01101111 01101101 01100101 00100000 01110011 01101111 011100
10 01110100 00100000 01101111 01100110 00100000 01110011 01110000 01100101 01100011 01101001 01100001 01101100 0
0100000 01100101 01101110 01100011 01101111 01100100 01101001 01101110 01100111 00101110 00100000 01000011 01100
001 01101110 00100000 01111001 01101111 01110101 00100000 01100110 01101001 01100111 01110101 01110010 01100101
00100000 01101111 01110101 01110100 00100000 01110111 01101000 01100001 01110100 00100000 01110100 01101000 0110
1001 01110011 00100000 01100101 01101110 01100011 01101111 01100100 01101001 01101110 01100111 00100000 01101001
01110011 00111111 00100000 00101000 01101000 01101001 01101110 01110100 00111010 00100000 01101001 01100110 001
00000 01111001 01101111 01110101 00100000 01101100 01101111 01101111 01101011 00100000 01100011 01100001 01110001
0 01100101 01100110 01110101 01101100 01101100 01111001 00101100 00100000 01111001 01101111 01110101 00100111 01
101100 01101100 00100000 01101110 01101111 01110100 01101001 01100011 01100101 00100000 01110100 01101000 011000
01 01110100 00100000 01110100 01101000 01100101 01110010 01100101 00100000 01101111 01101110 01101100 01111001 0
0100000 01100011 01101000 01100001 01110010 01100001 01100011 01110100 01100101 01110010 01110011 00100000 01110
000 01110010 01100101 01110011 01100101 01101110 01110100 00100000 01100001 01110010 01100101 00100000 01000001
00101101 01011010 00101100 00100000 01100001 00101101 01111010 00101100 00100000 00110000 00101101 00111001 0010
1100 00100000 01100001 01101110 01100100 00100000 01110011 01101111 01101101 01100101 01110100 01101001 01101101
01100101 01110011 00100000 00101111 00100000 01100001 01101110 01100100 00100000 00101011 00101110 00100000 010
10011 01100101 01100101 00100000 01101001 01100110 00100000 01111001 01101111 01110101 00100000 01100011 0110000
1 01101110 00100000 01100110 01101001 01101110 01100100 00100000 01100001 01101110 00100000 01100101 01101110 01
100011 01101111 01100100 01101001 01101110 01100111 00100000 01110100 01101000 01100001 01110100 00100000 011011
00 01101111 01101111 01101011 01110011 00100000 01101100 01101001 01101011 01100101 00100000 01110100 01101000 0
1101001 01110011 00100000 01101111 01101110 01100101 00101110 00101001 00001010 01010100 01101101 01010110 00110
011 01001001 01000111 01001110 01101111 01011001 01010111 01111000 01110011 01011010 01010111 00110101 01101110
01011010 01010011 01000101 01100111 01010001 00110010 01000110 01110101 01001001 01001000 01101100 01110110 0110
0100 01010011 01000010 01101101 01100001 01010111 01100100 00110001 01100011 01101101 01010101 01100111 01100010
00110011 01010110 00110000 01001001 01001000 01100100 01101111 01011001 01011000 01010001 01101110 01100011 011
11001 01000010 01101110 01100010 00110010 01101100 01110101 01011010 01111001 01000010 01110110 01100010 0110100
1 01000010 01101111 01011010 01011000 01001010 01101100 01010000 01111001 01000010 01001010 01100100 01000011 01
000010 01110011 01100010 00110010 00111001 01110010 01100011 01111001 01000010 01110011 01100001 01010111 011101
00 01101100 01001001 01001000 01010010 01101111 01011010 01010011 01000010 01110011 01011010 01011000 01010010 0
0110000 01011010 01011000 01001010 01111010 01001001 01000111 01000110 01111001 01011010 01010011 01000010 01111
010 01100001 01000111 01101100 01101101 01100100 01000111 01010110 01101011 01001001 01000111 01001010 00110101
01001001 01001000 01001110 01110110 01100010 01010111 01010101 01100111 01011001 00110010 00111001 01110101 0110
0011 00110011 01010010 01101000 01100010 01101110 01010001 01110101 01001001 01000011 01101000 01101111 01100001
01010111 00110101 00110000 01001111 01101001 01000010 00110101 01100010 00110011 01010101 01100111 01100010 010
10111 01101100 01101110 01100001 01001000 01010001 01100111 01100100 00110010 01000110 01110101 01100100 0100001
1 01000010 00110000 01100010 01111001 01000010 01111010 01100100 01000111 01000110 01111001 01100100 01000011 01
000010 01110011 01100010 00110010 00111001 01110010 01100001 01010111 00110101 01101110 01001001 01001000 010101
10 01110111 01001001 01000110 01001010 01110110 01100010 01010111 01000110 01110101 01001001 01001000 01000010 0
1101100 01100010 00110011 01000010 01110011 01011010 01010011 01101011 01110101 01000011 01101101 01110100 00110
010 01011001 01101110 01001110 01111000 01100011 01101101 01010001 01110011 01001001 01000111 01101100 00110101
01011010 01010011 01100100 01101001 01100010 01111001 01000010 01110001 01100010 01110001 01101110 01100100 00110101 0101
1001 00110010 01010001 01100111 01011010 01001000 01001010 01110110 01011001 01101101 00111000 01101000 01001001
01000110 01101000 00110101 01011010 01111001 01000010 01110111 01100101 01010111 01001001 01100111 01011010 010
01000 01001010 01110110 01001001 01001000 01000010 01111010 01100101 01000111 01110100 00110010 01001001 0100001
1 01101000 01110010 01100101 01000111 00110100 01100111 01100100 00110010 01110100 01110000 01100010 01000111 00
111000 01100111 01011010 01001000 01001010 01110110 01001001 01001000 01001010 01110010 01011001 01101101 001101
01 01110110 01011001 00110010 01010001 01110101 01001100 01101001 00110100 01110000 01001001 01001000 01110000 0
1110010 01011001 01101101 01010001 00110110 01001001 01000111 01110011 01100111 01011001 00110010 01010110 01110
011 01011001 00110010 01010010 01111010 01011010 01000111 01010110 01101011 01100011 00110011 01101100 00110100
01001001 01000111 00110001 01111010 01100101 01100110 01001010 01110110 01001010 01110110 01011001 01101001 00110100 01100111 0101
```

01001001 01000111 00110001 01110100 01100101 01101110 01001010 01101110 01011001 01001000 01000010 00110101 01100100
0101 00110011 01100111 01100111 01011010 01001000 01001010 01110110 01001001 01001000 01000010 00110101 01100100
01101110 01011010 00110101 01011010 00110011 01001110 00110100 01100011 01010011 01000010 01101011 01100010 001
10010 01101000 01101011 01001100 01000011 01000010 01010100 01001010 00110010 01011010 01110110 01001001 0100011
1 01010010 01110010 01100100 01010111 00111001 00110100 01001001 01001000 01100100 01110000 01001001 01001000 01
100100 01110110 01011001 00110010 01001110 01110010 01100011 01010111 00111000 01100111 01100001 00110011 011010
00 01110101 01001001 01000111 01001010 01110110 01100101 01101110 01011010 01110010 01100010 01010111 00111001 0
1110101 01001001 01000111 00111001 01101101 01100010 00110010 01001010 01110000 01001001 01000111 01110100 00110
010 01100101 01101110 01001010 01110010 01100010 01000111 00111001 01101011 01100011 00110010 00110000 01100111
01100010 01011000 01001010 01110010 01011001 01101101 01110100 01110100 01011010 01000111 00111001 01101001 0100
1001 01000111 01100100 01111010 01011010 01001000 01001001 01100111 01100001 01111001 01000010 01110100 01100101
01010111 01001010 01101001 01100010 00110010 01001110 00110110 01100101 01011000 01101000 01110101 01100010 001
10011 01101000 01110100 01100010 01111001 01000010 01101011 01100101 01010011 01000010 01110010 01001001 0100011
1 00110101 01111010 01100011 01001000 01000010 01110110 01011001 01101101 00111001 00110100 01011010 01000011 01
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: fdcysn{h0v_di4du_vi4d_t_r4yy_rxqld0}. qgf vtyy cthe disd s ygd gc rxqldgnxslmq tu pfud zftyethn gcc ditu ugxd
gc zsutr bhgyykenk, she td xksyyq tu hgd ug zse scdkx syy. iglk qgf khpgqke dik risyykhnk!
```

词频分析得到

```
congratulations! you have finished the beginner cryptography challenge. here is a flag for all your hard efforts
: utflag{n0w_th4ts_wh4t_i_c4ll_crypt0}. you will find that a lot of cryptography is just building off this sort
of basic knowledge, and it really is not so bad after all. hope you enjoyed the challenge!
```

flag{n0w_th4ts_wh4t_i_c4ll_crypt0}这个就是flag了

108.[GKCTF2020]Backdoor

查看题目 是一个cve漏洞

前面具体过程就不写了 那个流量的用记事本打开，然后是一个base64解密，另一个是公钥解密即可

脚本如下

```
from Crypto.Util import number
from gmpy2 import *
'''#密文c
02142af7ce70fe0ddae116bb7e96260274ee9252a8cb528e7fdd29809c2a6032727c05526133ae4610ed944572ff1abfcd0b17aa22ef44a2
'''
'''e = 65537
n = 155189610416250748761824045853940987814871410592854559270243212767838311221687450763597803430780112164805875
75072479784829258678691739
'''
'''p:3386619977051114637303328519173627165817832179845212640767197001941
q:4582433561127855310805294456657993281782662645116543024537051682479
'''
vals=39
M=1
n = mpz(15518961041625074876182404585394098781487141059285455927024321276783831122168745076359780343078011216480
587575072479784829258678691739)
primes = [2, 3, 5, 7, 11, 13, 17, 19, 23, 29, 31, 37, 41, 43, 47, 53, 59, 61, 67, 71, 73, 79, 83, 89, 97, 101, 1
03, 107, 109, 113, 127, 131, 137, 139, 149, 151, 157, 163, 167, 173, 179, 181, 191, 193, 197, 199, 211, 223, 227
, 229, 233, 239, 241, 251, 257, 263, 269, 271, 277, 281, 283, 293, 307, 311, 313, 317, 331, 337, 347, 349, 353,
359, 367, 373, 379, 383, 389, 397, 401, 409, 419, 421, 431, 433, 439, 443, 449, 457, 461, 463, 467, 479, 487, 49
1, 499, 503, 509, 521, 523, 541, 547, 557, 563, 569, 571, 577, 587, 593, 599, 601, 607, 613, 617, 619, 631, 641,
643, 647, 653, 659, 661, 673, 677, 683, 691, 701, 709, 719, 727, 733, 739, 743, 751, 757, 761, 769, 773, 787, 7
97, 809, 811, 821, 823, 827, 829, 839, 853, 857, 859, 863, 877, 881, 883, 887, 907, 911, 919, 929, 937, 941, 947
, 953, 967, 971, 977, 983, 991, 997, 1009, 1013, 1019, 1021, 1031, 1033, 1039, 1049, 1051, 1061, 1063, 1069, 108
7, 1091, 1093, 1097, 1103, 1109, 1117, 1123, 1129, 1151, 1153, 1163, 1171, 1181, 1187, 1193, 1201, 1213, 1217, 1
223, 1229, 1231, 1237, 1249, 1259, 1277, 1279, 1283, 1289, 1291, 1297, 1301, 1303, 1307, 1319, 1321, 1327, 1361,
1367, 1373, 1381, 1399, 1409, 1423, 1427, 1429, 1433, 1439, 1447, 1451, 1453, 1459, 1471, 1481, 1483, 1487, 148
9, 1493, 1499, 1511, 1523, 1531, 1543, 1549, 1553, 1559, 1567, 1571, 1579, 1583, 1597, 1601, 1607, 1609, 1613, 1
619, 1621, 1627, 1637, 1657, 1663, 1667, 1669, 1693, 1697, 1699, 1709, 1721, 1723, 1733, 1741, 1747, 1753, 1759,
1777, 1783, 1787, 1789, 1801, 1811, 1823, 1831, 1847, 1861, 1867, 1871, 1873, 1877, 1879, 1889, 1901, 1907, 191
3, 1931, 1933, 1949, 1951, 1973, 1979, 1987, 1993, 1997, 1999, 2003, 2011, 2017, 2027, 2029, 2039, 2053, 2063, 2
069, 2081, 2083, 2087, 2089, 2099, 2111, 2113, 2129, 2131, 2137, 2141, 2143, 2153, 2161, 2179, 2203, 2207, 2213,
2221, 2237, 2239, 2243, 2251, 2267, 2269, 2273, 2281, 2287, 2293, 2297, 2309, 2311, 2333, 2339, 2341, 2347, 235
1, 2357, 2371, 2377, 2381, 2383, 2389, 2393, 2399, 2411, 2417, 2423, 2437, 2441, 2447, 2459, 2467, 2473, 2477, 2
503, 2521, 2531, 2539, 2543, 2549, 2551, 2557, 2579, 2591, 2593, 2609, 2617, 2621, 2633, 2647, 2657, 2659, 2663,
2671, 2677, 2683, 2687, 2689, 2693, 2699, 2707, 2711, 2713, 2719, 2729, 2731, 2741, 2749, 2753, 2767, 2777, 278
9, 2791, 2797, 2801, 2803, 2819, 2833, 2837, 2843, 2851, 2857, 2861, 2879, 2887, 2897, 2903, 2909, 2917, 2927, 2
939, 2953, 2957, 2963, 2969, 2971, 2999]
for x in range(0, vals):
```

```

M=M*primes[x]

for a in range(1,20):
    for k in range(50):
        p=mpz(k*M+(65537**a %M))
        if is_prime(p):
            q = mpz(n//p)
            if is_prime(q):
                print('p=%d\nq=%d'%(p,q))

import gmpy2
p =gmpy2.mpz(4582433561127855310805294456657993281782662645116543024537051682479)
q =gmpy2.mpz(3386619977051114637303328519173627165817832179845212640767197001941)
e =gmpy2.mpz(65537)
phi_n= (p - 1) * (q - 1)
d = gmpy2.invert(e, phi_n)
print("d is:%d"%(d))
# d is:114995697859901812901421504475409862997293136893980437948652229147514562710973371046228849923451202789592
13140333860537563347711742153
import gmpy2,binascii

from Crypto.Util.number import long_to_bytes
c = int("02142af7ce70fe0ddae116bb7e96260274ee9252a8cb528e7fdd29809c2a6032727c05526133ae4610ed944572ff1abfcd0b17a
a22ef44a2",16)

d = 114995697859901812901421504475409862997293136893980437948652229147514562710973371046228849923451202789592131
40333860537563347711742153
n = 155189610416250748761824045853940987814871410592854559270243212767838311221687450763597803430780112164805875
75072479784829258678691739
m = pow(c,d,n)
print('m = ',m)
print('m = ',hex(m))
print('long_to_bytes(m) =',long_to_bytes(m))

```

运行得到

```

p=4582433561127855310805294456657993281782662645116543024537051682479
q=3386619977051114637303328519173627165817832179845212640767197001941
d is:11499569785990181290142150447540986299729313689398043794865222914751456271097337104622884992345120278959213
140333860537563347711742153
m = 56006392793406552883106744981771255916153714828118097099130014407421330832850082353964262145657222269
m = 0x666c61677b37363039353863392d636361392d343538622d396362652d6561303761613136363865347d
long_to_bytes(m) = b'flag{760958c9-cca9-458b-9cbe-ea07aa1668e4}'

```

109.[V&N2020 公开赛]easy_RSA

[查看题目](#)

```

from random import randint
from gmpy2 import *
from Crypto.Util.number import *

def getprime(bits):
    while 1:
        n = 1
        while n.bit_length() < bits:
            n *= next_prime(randint(1,1000))
        if isPrime(n - 1):
            return n - 1

m = bytes_to_long(b'flag{*****}')

p = getprime(505)
q = getPrime(512)
r = getPrime(512)
assert m < q

n = p * q * r
e = 0x10001
d = invert(q ** 2, p ** 2)
c = pow(m, 2, r)
cipher = pow(c, e, n)

print(n)
print(d)
print(cipher)

...

7941371739956577280160664419383740967516918938781306610817149744988379280561359039016508679365806108722198157199
0588078927038375582806787114204112429140596580553663481231064733351865056174189566307806498949452333459852794711
0688863517725601146897908332060510325617844699323032044379024028515826023692651904241337820429851471489072532583
1769281505530787739922007367026883959544239568886349070557272869042275528961483412544495589811933856131557221673
534170105409
7515987842794170949444517202158067021118454558360145030399453487603693522695746732547224100845570119375977629070
7023089912213887219522589697523059043787244020025459471825298596045844000489830918615947202997917438875212284927
14135449584003054386457751933095902983841246048952155097668245322664318518861440
1618155233923718966393124032999431934705026408748451436388483012584983753140040289666712916510617403356206112730
6134852270841283140436659133571063017368170624129271357162815443486121503288672265151840789663971807716241487975
2803654824334331650150336478309255048043974940430112227705673285739941380529389924931304568466214633344866820956
7898831091274930053147799756622844119463942087160062353526056879436998061803187343431081504474584816590199768034
450005448200

...

```

还是用分解N的网站在线分解得到pqr

sequences	Report Results	Factor Tables	Status
<input type="text" value="7941371739956577280160664419383740967516918938781306610817149744988379280561359039016508679365"/>			<input type="button" value="Factorize!"/> (?)
Result:			
number			
$7941371739956577280160664419383740967516918938781306610817149744988379280561359039016508679365$			

脚本如下

```
from Crypto.Util.number import *
from gmpy2 import *
from sympy.ntheory.residue_ntheory import nthroot_mod

n = 794137173995657728016066441938374096751691893878130661081714974498837928056135903901650867936580610872219815
7199058807892703837558280678711420411242914059658055366348123106473335186505617418956630780649894945233345985279
4711068886351772560114689790833206051032561784469932303204437902402851582602369265190424133782042985147148907253
2583176928150553078773992200736702688395954423956888634907055727286904227552896148341254449558981193385613155722
1673534170105409
d = 751598784279417094944451720215806702111845455836014503039945348760369352269574673254722410084557011937597762
9070702308991221388721952258969752305904378724402002545947182529859604584400048983091861594720299791743887521228
492714135449584003054386457751933095902983841246048952155097668245322664318518861440
cipher = 1618155233923718966393124032999431934705026408748451436388483012584983753140040289666712916510617403356
2061127306134852270841283140436659133571063017368170624129271357162815443486121503288672265151840789663971807716
2414879752803654824334331650150336478309255048043974940430112227705673285739941380529389924931304568466214633344
8668209567898831091274930053147799756622844119463942087160062353526056879436998061803187343431081504474584816590
199768034450005448200
#分解得到的pqr
p = 102634610559478918970860957918259981057327949366949344137104804864768237961662136189827166317524151288799657
758536256924609797810164397005081733039415393
q = 753481019642093255216870893701969199468105266006827590697348061760453538130604158384110638368865442612905093
1519275383386503174076258645141589911492908993
r = 102690287677543062175637216649762619244079408837841938177866604137448661846459842388664637118733800728037470
92361041245422348883639933712733051005791543841
phn=(p-1)*(q-1)*(r-1)
e = 0x10001
d = invert(e,phn)
print(d)
#696507389127827123706661871654450815284301235334466795840166268431469309011195076922454367593796687564813693144
0284118553986550846849465375779258102181906351936906869410334002715845817435116465257017702306973898081396888890
5589773868014614548914146162321030615201748237234028805613294778275446387109635156263439686832386133548009179785
5608845695507892749062713128285934572259972389743839646950269761373893877746254763641876362977045864056456449244
814748352513

c=pow(cipher,d,n)
print(c)
#808109245511251639736110581690049008535531557408753834078830988533410679632559382367878788756992040481498664381
9898763828872716522338864714182757065213683

m=nthroot_mod(c,2,r)
print(m)
#56006392793430016468251971646527328995718100207125432393433900875091739391190683811783574991236326013

print(long_to_bytes(m))
```

运行得到

```

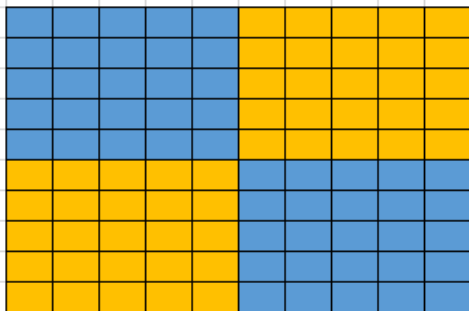
6965073891278271237066618716544508152843012353344667958401662684314693090111950769224543675937966875648136931440
2841185539865508468494653757792581021819063519369068694103340027158458174351164652570177023069738980813968888905
5897738680146145489141461623210306152017482372340288056132947782754463871096351562634396868323861335480091797855
6088456955078927490627131282859345722599723897438396469502697613738938777462547636418763629770458640564564492448
14748352513
8081092455112516397361105816900490085355315574087538340788309885334106796325593823678787887569920404814986643819
898763828872716522338864714182757065213683
56006392793430016468251971646527328995718100207125432393433900875091739391190683811783574991236326013
b'flag{fd462593-25e4-4631-a96a-0cd5c72b2d1b}'

```

四方密码简介

四方密码是用四个方块，然后配合字母来完成加密。四方密码需要密钥，要用任意两个单词作为密钥。比如说题目中的security和information来作为密钥，然后去处单词里面的重复字母作为真正的密钥，也就是security和informat。

然后就是介绍四方了，四方，也就是四个方块，为4个5x5的方块，方块的左上角和右下角，也就是图中的蓝色方块，是用来按顺序存放26个字母的。



然而这里只有25个小方格，所以一般是把字母q或z删去不要，或者是把i、j放在同一个格子里面。

然后右上角和左下角的方块存放的是我们的密钥，然后分别把密钥写在田字格的右上角和左下角，然后再按字母排列表的顺序写下来。注意，密钥中已经有的字母就去除不写，不要再次出现！

加密

首先第一步，填写好这个方格，就拿security和information来说，填写之后的内容是这样的，这里是去除了q。

a	b	c	d	e	s	e	c	u	r
f	g	h	i	j	i	t	y	a	b
k	l	m	n	o	d	f	g	h	j
p	r	s	t	u	k	l	m	n	o
v	w	x	y	z	p	v	w	x	z
i	n	f	o	r	a	b	c	d	e
m	a	t	b	c	f	g	h	i	j
d	r	g	h	j	k	l	m	n	o
k	l	p	s	u	p	r	s	t	u
v	w	k	y	z	v	w	x	y	z

<https://blog.csdn.net/ao52426055>

假设我们需要加密的内容是iamafool。首先把他拆分为两个两个一组，拆分完毕之后是ia ma fo ol。这里需要注意：四方密码加密只对偶数个字符有效，最后余下的一个字符将无法加密。然后按照顺序来，首先是ia，找到i在左上角方格中的位置和a在右下角方格中的位置。然后画出以他们连线为对角线的矩形，即为下图中的红色方块。

a	b	c	d	e	s	e	c	u	r
f	g	h	i	j	i	t	y	a	b
k	l	m	n	o	d	f	g	h	j
p	r	s	t	u	k	l	m	n	o
v	w	x	y	z	p	v	w	x	z
i	n	f	o	r	a	b	c	d	e
m	a	t	b	c	f	g	h	i	j
d	r	g	h	j	k	l	m	n	o
k	l	p	s	u	p	r	s	t	u
v	w	x	y	z	v	w	x	y	z

加密得到的字母与和明文在同一行的边角点，这里也就是i对应i，a对应o，ia加密得到的内容就是io，剩下的字符也是同样的操作方式。

解密

揭秘也很简单了，就是加密的逆过程，比如说对应上面加密的例子，密文是io，明文就是ia，也就是现在左上角的方格中找到i，左下角找到o，然后取同一行的边角点就好了。

flag{youngandsuccessful}

110[De1CTF2019]babyrsa

查看题目

```
import binascii
from data import e1,e2,p,q1p,q1q,hint,flag

n = [2012961535249176549934011294318831718054876159786130084730582714151046561967053684463455824643923037165883
6928103063432870245707180355907194284861510906071265352409579441048101084995923962148527097370705452070577098780
2462828200655737110156642919913720851570169012091141910685742086803977100428428359404284519495006076136346826841
1320876669402878927574852825428770575952849898630649426781719834065824187302480033601394629489168759101341493523
7821291805123285905335762719823771647853378892868896078424572232934360940672962436849523915563328779942134504499
568866135266628078485232098208237036724121481835035731201383423L, 3122165015562784996446641374941470061382384106
0149524451234901677160009099014018926581094879840097248543411980533066831976617023676225625067854003317018794041
7236125560084715790604288981177905879910556813804082633827618416257144158790874780727719681603849099199580109836
6936836078850528885594612415951311884774799865642252141498029521264667585069093788376400057166757438141914437282
4211798018586804674824564606122592483286575800685232128273820087791811663878057827386379787882962763290066072231
2488149204682647416540860110726382110754454478436910498472624857593932908531170728684068618407938958162159568695
23289231421L, 299445375153979533615209227741241926055247113067538353037034788904141635107746059798334313021216
3893562518749177920076382992258210188496485206738137867724528228095465711298163102072328832397713241228848049934
1895830946000940634287217318900844923795957746911415899120243347671058135624381571376280247845439027380837743068
5157110095496727966308001254107517967559384019734279861840997239176254236069001453544559786063915970071130087811
1239120443122195355138806639138313587903766504390836606118311562051138737931068802558821144220257469864033550669
96567909581710647746463994280444700922867397754748628425967488232530303L, 25703437855600135215185778453583925446
912731661604054184163883272265503323016295700357253105301146726667897497435325799749514783545704155542214017785
3610473729615431605631403944911638649432366848374983314780055740336848954227316948908022200936890399365849826390
5567516798684211462607069796613434661148186901892016282065916190920443378756167250809872483501712225782004396969
9969830574239426071743141325984212691697225182244782488368810764846398373430793246369971451998350348333677430799
3536127614999099787590531364277521448604638136861963855189229278778313762226143352891526933342676894735855291974
0901860982679180791L]

c = [1913143266121790847026233842129969199852615779058354415674198123882215856398852022598691523457003738388811
2724408392918113942721994125505014727545946133307329781747600302829588248042922635714391033431930411180545085316
4380843179273487052419275704327578929850913960449500854624295754400606529672538450413983996484423400429708144155
7190405766702815751297107938460172481630807863184448011020178734358307381518677179047771204005115718031880442212
047200763672206398931532086358063133064711699381977750684150950416298085261478841177681677867236865666207391847
```

```
046483954029213495373613490690687473081930148461830425717614569L, 1534189843322663823516007202987573382695679998
2958107910250055958334922460202554924743144122170018355117452459472017133614642242411479849369061482860570279863
6924256215260568628084251352676085448558333583140712006873404425128565752787129866415730124567294026605973396094
4377114534718126828505072892599351870489900541618725000330458123070144470515741279078702792681071099864619146713
0550713600765898234392350153965811595060656753711278308005193370936296124790772689433773414703645703910742193898
4718000813214690552117093398463925007065236701452590242678583682169021764898147896794722273433630354285419151183
78163012031L, 18715065071648040017967211297231106538139985087685358555650567057715550586464814763683688299037897
1828450075785714013590612137776451144146429030770035681555084658196285537471732442359365868124454400954507551543
5764673708707160581198416341659027835260543336232794904824372255626297990948820244253030750581937159474793622383
523358694542352256938701002370646382097846105014981763307729234675737702252155130837154876831885888669150418885
0880893245348925061997244867834462673367898727821378955525093535833058801449477141100098931341621853823099926044
35664777436197587312317224862723813510974493087450281755452428746194446L, 22822845612248582931384804474633192624
7491884763014877011247270312854903259218779728996559261519970985787900827176643346203232849858034096887126018966
9707518557157836592424973257334362931639831072584824103123486522582531666152363874396482744561758133655406410364
4421749832270055018609278208712607118610088301206170568835145257987096017440881359994655983386357942751231491654
9893358015994503236388061352492191302334120943965714596233221346857340286379692057181241820081481708623426228033
8221161622789516829363805084715652121739036183264026120868756523770196284142271849879003202190966150390061195469
351716819539183797L]
f=lambda m,e,n,c:pow(m,e,n)==c
assert(sum(map(f,[p]*4,[4]*4,n,c))==4)

ee1 = 42
ee2 = 3
ce1 = 457226517863401239469608150030593225288104818413782472806428685536076921495091269628725830371424613988066
8948914174149497483688234150523425532568321909216305284346163233844252901150237893114035611175693271282251681402
3166068902569458299933391973504078898958921809723346229893913662577294963528318424676803942288386430172430880307
6197481868638900501139345738205055709281090178426475982666343444471823478493677145646863418710075058867283937511
4703355688921760464735562855750220836441226994490801130506412294144651699016892470968409220018386065317385627238
4
ce2 = 13908468332335671584691364399323259923496968891291039354007602393194544095397253897470592138352383730478
9919821112868937404972957814687530923196293655440328788299996784034621669520842458273977703426107955039591804842
1086843927009452479936045850799096750074359160775182238980989229190157551197830879877097703347301072427149474991
8038683257699673323569508635185049654865654640597704514585577449497352821317279560562792928006942038661672702689
8843738994570311707060448899924775013956861493996588521127682198758688290815958586351456119190504024496765544421
9603287214405014887994238259270716355378069726760953320025828158
tmp = 864078778078609835167779565982540757684070450697854309005171742813414963447462554999012718960925081621571
487444725528982424037419052194840720949809891134854871222612682162490991065015935449289960707882463387
n = 1591158155579679861471162528850830970479183751623212241044095883072607882106905040401282089626007175138043
6992710638364294658173571101596931605797509712839622479368850251206419748090059752427303611760004621378226431226
983665746837790562715301818656481158629475272127878246295162048323130264563900477681747656870409506365304805490
1440127905434609803039510038700411157427881374963098672470626365516628958623045397595377379194540858948467937185
4113457758157492241225180907090235116325034822993748409011554673180494306003272836905082473475046277554085737627
846557240367696214081276345071055578169299060706794192776825039
assert(pow(e1,ee1,n)==ce1)
assert(pow(e2+tmp,ee2,n)==ce2)

e = 46531
n = 162785240342783648429643860624761135170679118916997899913559821210849739517383240633051906308655115548883302
1582772488796456597960780829416828299582586498260375938132304890781496127901237534649778104641720495410107645735
0988751188332353062731641153547102721113593787978587135707313755661153376485647168543680503160420091693269984008
7644442912894868058404399066203131623440579565948361975215017553783879446092461206623357901109016237409904515866
2184621204795008420725159516914101564544921784718068335762638356563131725391394288639649439618983743242907825157
322937891740084183219073751876329732390158686664595327850603
c = 149921321409961603309673075585031172556269257774266119785183390506710130414907246168926349110309183608679748
9437153916085382718059610089218073577068872327076538769760442671567044527081962670936456647878127367611592165796
7761494619448095207169386364541164659123273236874649888236433399127407801843412677293516986398190165291102109310
458304626261648346825196743539220198199366711858135271877662410355585767124059539212746916068251033553103486076
1123305272580523676322034324987384964621985095494534679101585826171596795246102165030730745443451085186986296423
6227932964442289459508441345652423088404453536608812799355469
hint=int(binascii.hexlify(hint).16)
```



```

hint=int(binascii.hexlify(hint),16)
assert(q1p*q1q==n)
assert(q1p<q1q)
assert(c==pow(hint,e,n))

flag=int(binascii.hexlify(flag),16)
q1=q1p
q2 = 1144011882274795846808840461512997046569205361687671329165891823575834610533363869961237832949325665677736
9542668944741031196945645857473118751297486829709263867751528358499441638287245016704641657347265884162769098722
8528798356894803559278308702635288537653192098514966089168123710854679638671424978221959513
c1 = 2627399757539302816909427843212523390359061968463407132375103823645576853795434987650744488257993421943326
8118112977004607501812203342198322788771961011202823060316652730302103638635078141444734715038378381686978400659
8225583375458609586450854602862569022571672049158809874763812834044257419199631217527367046624888837755311215081
1733865238060867832661983902890972311681726923266536573935225617419479518875771566666635842491088993270539518914
8635517993977015055099581247832773591700619457441251881929930378324388696245539978360122922771878708178539101042
4030509937403600351414176138124705168002288620664809270046124
c2 = 7395591129228876649030819616685821899204832684995757724924450812977470787822266387122334722132760470911599
1763626172252183454044682700145488172677276698728968381064515203928064974665769070632956037466600031884401709194
9015725082930817331071531892577164310506488262074617126649985904903801690216259926140905090714082335299075029823
9508355767238575709803167676810456559665476121149766947851911064706646506705397091626648713684511780456955453552
0204609096380161341245904384257388268286947739605142219101094739414514714316379031822057387381094297364250256213
08300895473186381826756650667842656050416299166317372707709596
assert(c1==pow(flag,e1,p*q1))
assert(c2==pow(flag,e2,p*q2))

```

这是一道综合的rsa

包括中国剩余定理特解为解

小指数攻击(枚举爆破)

在线质因数分解factordb

rsa中e关于 $(p-1)(q-1)$ 的逆元不唯一,同时兼有中国剩余定理方程合并.

得到两个方程

$$c1 = m e_1 \%(p q_1)$$

$$c2 = m e_2 \%(p q_2)$$

在通常情况下,我们只需要一个方程就可以求得一个满足条件的最小解作为解

这道题目最后给了两个方程,显然我们求得的解必须满足这两个方程.

首先看第一个方程.

通过计算得知 $\gcd(e_1, (p-1)(q_1-1))=14$.

可以知道该rsa式子并不能用常规方法.通过变形.

$c1 = m e_1 \%(p q_1) = (m/14) e_1 / 14 \%(p q_1)$,我们可以在此条件下求出 $m/14$ 的通解(显然最小特解很大可能不是答案,因为这个解还需要满足第二个方程)

记 $p q_1 = n_1, p q_2 = n_2$,记 $e_1/14$ 的逆元为 d_1 ,方程2中 $e_2/14$ 的逆元为 d_2

记该通解为方程

$$m/14 \equiv c_1 d_1 \pmod{n_1}$$

同理可得:

$$\gcd(e_2, (p-1)(q_2-1))=14$$

则:

$$m/14 \equiv c_2 d_2 \pmod{n_2}$$

此时尝试用中国剩余定理来求出一个特解,但是遗憾的是该特解不是最终解.

于是尝试爆破,但是同样发现过了很长时间还是爆破不了.

于是转变方法.将该同余方程组进行细化.

$$c_1 d_1 \equiv a_1, c_2 d_2 \equiv a_2$$

得到

$$m/14 \equiv a_1 \pmod{p}$$

$$m/14 \equiv a_1 \pmod{q_1}$$

$$m/14 \equiv a_2 \pmod{p}$$

$$m/14 \equiv a_2 \pmod{q_2}$$

由于 m 的指数过大,我们尝试通过构造一个新的rsa式子来降解 m 的指数.

理论上4个方程有6种合并方式.但是通过计算

$\gcd(p-1, 7) \neq 1$ 所以如果选择 p 的话显然是行不通的.

于是舍弃 p

选择 q_1, q_2 进行合并.

合并函数

脚本如下

```
import gmpy2
import math
def merge(a1, n1, a2, n2):
    d = math.gcd(n1, n2)
    c = a2 - a1
    if c % d != 0:
        return 0
    c = (c % n2 + n2) % n2
    c = c // d
    n1 = n1 // d
    n2 = n2 // d
    c *= gmpy2.invert(n1, n2)
    c %= n2
    c *= n1 * d
    c += a1
    global n3
    global a3
    n3 = n1 * n2 * d
    a3 = (c % n3 + n3) % n3
    return 1
```

```

def exCRT(a,n):
    a1=a[0]
    n1=n[0]
    le= len(a)
    for i in range(1,le):
        a2 = a[i]
        n2=n[i]
        if not merge(a1,n1,a2,n2):
            return -1
        a1 = a3
        n1 = n3
    global mod
    mod=n1
    return (a1%n1+n1)%n1
def exCRT_getequation(a,n):
    a1=a[0]
    n1=n[0]
    le= len(a)
    for i in range(1,le):
        a2 = a[i]
        n2=n[i]
        if not merge(a1,n1,a2,n2):
            return -1
        a1 = a3
        n1 = n3
    return (a1,n1)
#a为余数列表
#n为模数列表
n = [2012961535249176549934011294318831718054876159786130084730582714151046561967053684463455824643923037165883
6928103063432870245707180355907194284861510906071265352409579441048101084995923962148527097370705452070577098780
2462828200655737110156642919913720851570169012091141910685742086803977100428428359404284519495006076136346826841
1320876669402878927574852825428770575952849898630649426781719834065824187302480033601394629489168759101341493523
7821291805123285905335762719823771647853378892868896078424572232934360940672962436849523915563328779942134504499
568866135266628078485232098208237036724121481835035731201383423, 31221650155627849964466413749414700613823841060
1495244512349016771600090990140189265810948798400972485434119805330668319766170236762256250678540033170187940417
2361255600847157906042889811779058799105568138040826338276184162571441587908747807277196816038490991995801098366
9368360788505288855946124159513118847747998656422521414980295212646675850690937883764000571667574381419144372824
2117980185868046748245646061225924832865758006852321282738200877918116638780578273863797878829627632900660722312
4881492046826474165408601107263821107544544784369104984726248575939329085311707286840686184079389581621595686952
3289231421, 2994453751539795336152092277412419260552471130675383530370347889041416351077746055979833431302121638
9356251874917792007638299225821018849648520673813786772452822809546571129816310207232883239771324122884804993418
9583094600094063428721731890084492379595774691141589912024334767105813562438157137628024784543902738083774306851
5711009549672796630800125410751796755938401973427986184099723917625423606900145354455978606391597007113008781112
3912044312219535513880663913831358790376650439083660611831156205113873793106880255882114422025746986403355066996
567909581710647746463994280444700922867397754748628425967488232530303, 25703437855600135215185778453583925446912
7316616040541841638832722655033230162957003572531053011467266678974974355325799749514783545704155542214017785361
0473729615431605631403944911638649432366848374983314780055740336848954227316948908022200936890399365849826390556
7516798684211462607069796613434661148186901892016282065916190920443378756167250809872483501712225782004396969996
9830574239426071743141325984212691697225182244782488368810764846398373430793246369971451998350348333677430799353
6127614999099787590531364277521448604638136861963855189229278778313762226143352891526933342676894735855291974090
1860982679180791]
c = [1913143266121790847026233842129969199852615779058354415674198123882215856398852022598691523457003738388811
2724408392918113942721994125505014727545946133307329781747600302829588248042922635714391033431930411180545085316
4380843179273487052419275704327578929850913960449500854624295754400606529672538450413983996484423400429708144155
7190405766702815751297107938460172481630807863184448011020178734358307381518677179047771204005115718031880442212
047200763672206398931532086358063133064711699381977750684150950416298085261478841177681677867236865666207391847
046483954029213495373613490690687473081930148461830425717614569, 15341898433226638235160072029875733826956799982
9581079102500559583349224602025549247431441221700183551174524594720171336146422424114798493690614828605702798636
9242562152605686280842513526760854485583335831407120068734044251285657527871298664157301245672940266059733960944

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3771145347181268285050728925993518704899005416187250003304581230701444705157412790787027926810710998646191467130
5507136007658982343923501539658115950606567537112783080051933709362961247907726894337734147036457039107421938984
7180008132146905521170933984639250070652367014525902426785836821690217648981478967947222734336303542854191511837
8163012031, 1871506507164804001796721129723110653813998508768535855565056705771555058646481476368368829903789718
2845007578571401359061213777645114414642903077003568155508465819628553747173244235936586812445440095450755154357
6467370870716058119841634165902783526054333623279490482437225562629799094882024425303075058193715947479362238352
335869454235225693870100237064638209784610501498176330772923467573770225215513083715487683188588866915041888508
8089324534892506199724486783446267336789872782137895552509353583305880144947714110009893134162185382309992604435
66477436197587312317224862723813510974493087450281755452428746194446, 22822845612248582931384804474633192624749
1884763014877011247270312854903259218779728996559261519970985787900827176643346203232849858034096887126018966970
7518557157836592424973257334362931639831072584824103123486522582531666152363874396482744561758133655406410364442
1749832270055018609278208712607118610088301206170568835145257987096017440881359994655983386357942751231491654989
3358015994503236388061352492191302334120943965714596233221346857340286379692057181241820081481708623426228033822
1161622789516829363805084715652121739036183264026120868756523770196284142271849879003202190966150390061195469351
716819539183797]
p_4=exCRT(c,n)
p=gmpy2.iroot(p_4,4)[0]
print(p)
#中国剩余定理
ee1 = 42
ee2 = 3
ce1 = 457226517863401239469608150030593225288104818413782472806428685536076921495091269628725830371424613988066
8948914174149497483688234150523425532568321909216305284346163233844252901150237893114035611175693271282251681402
3166068902569458299933391973504078898958921809723346229893913662577294963528318424676803942288386430172430880307
6197481868638900501139345738205055709281090178426475982666343444471823478493677145646863418710075058867283937511
4703355688921760464735562855750220836441226994490801130506412294144651699016892470968409220018386065317385627238
4
ce2 = 13908468332335671584691364399323259923496968891291039354007602393194544095397253897470592138352383730478
9919821112868937404972957814687530923196293655440328788299996784034621669520842458273977703426107955039591804842
1086843927009452479936045850799096750074359160775182238980989229190157551197830879877097703347301072427149474991
8038683257699673323569508635185049654865654640597704514585577449497352821317279560562792928006942038661672702689
8843738994570311707060448899924775013956861493996588521127682198758688290815958586351456119190504024496765544421
9603287214405014887994238259270716355378069726760953320025828158
tmp = 86407877807860983516779565982540757684070450697854309005171742813414963447462554999012718960925081621571
487444725528982424037419052194840720949809891134854871222612682162490991065015935449289960707882463387
n = 1591158155579679861471162528850830970479183751623212241044095883072607882106905040401282089626007175138043699
2710638364294658173571101596931605797509712839622479368850251206419748090059752427303611760004621378226431226983
6657468377790562715301818656481158629475272127878246295162048323130264563900477681747656870409506365304805490144
0127905434609803039510038700411157427881374963098672470626365516628958623045397595377379194540858948467937185411
3457758157492241225180907090235116325034822993748409011554673180494306003272836905082473475046277554085737627846
557240367696214081276345071055578169299060706794192776825039
#assert(pow(e1, ee1, n)==ce1)
#assert(pow(e2+tmp, ee2, n)==ce2)
#n = 159115815557967986147116252885083097047918375162321224104409588307260788210690504040128208962600717513804
3699271063836429465817357110159693160579750971283962247936885025120641974809005975242730361176000462137822643122
6983665746837779056271530181865648115862947527212787824629516204832313026456390047768174765687040950636530480549
0144012790543460980303951003870041115742788137496309867247062636551662895862304539759537737919454085894846793718
5411345775815749224122518090709023511632503482299374840901155467318049430600327283690508247347504627755408573762
7846557240367696214081276345071055578169299060706794192776825039
def doit(ee,n,ce):
    k=0
    while True:
        x=ce+k*n
        if gmpy2.iroot(x,ee)[1]:
            return gmpy2.iroot(x,ee)[0]
        k=k+1
e1=doit(ee1,n,ce1)
e2=doit(ee2,n,ce2)-tmp
##小指数爆破
n1n=127587319253436643569312142058559706815497211661083866592534217079310497260365307426095661281103710042392775
```

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4538661746574049855390667416841960201378404729501023802320677864003226009029389849163556317144396683266713101609
16766472897536055371474076089779472372913037040153356437528808922911484049460342088834871
q1=q1p
##factordb爆破
q2 = 1144011882274795846808840461512997046569205361687671329165891823575834610533363869961237832949325665677736
9542668944741031196945645857473118751297486829709263867751528358499441638287245016704641657347265884162769098722
8528798356894803559278308702635288537653192098514966089168123710854679638671424978221959513
c1 = 2627399757539302816909427843212523390359061968463407132375103823645576853795434987650744488257993421943326
8118112977004607501812203342198322788771961011202823060316652730302103638635078141444734715038378381686978400659
8225583375458609586450854602862569022571672049158809874763812834044257419199631217527367046624888837755311215081
1733865238060867832661983902890972311681726923266536573935225617419479518875771566666635842491088993270539518914
8635517993977015055099581247832773591700619457441251881929930378324388696245539978360122922771878708178539101042
4030509937403600351414176138124705168002288620664809270046124
c2 = 7395591129228876649030819616685821899204832684995757724924450812977470787822266387122334722132760470911599
1763626172252183454044682700145488172677276698728968381064515203928064974665769070632956037466600031884401709194
9015725082930817331071531892577164310506488262074617126649985904903801690216259926140905090714082335299075029823
9508355767238575709803167676810456559665476121149766947851911064706646506705397091626648713684511780456955453552
0204609096380161341245904384257388268286947739605142219101094739414514714316379031822057387381094297364250256213
08300895473186381826756650667842656050416299166317372707709596
#assert(c1==pow(flag,e1,p*q1))
#assert(c2==pow(flag,e2,p*q2))
phi1=(p-1)*(q1-1)
phi2=(p-1)*(q2-1)
xx1=gmpy2.gcd(e1,phi1)
xx2=gmpy2.gcd(e2,phi2)
d1=gmpy2.invert(e1//xx1,phi1)
d2=gmpy2.invert(e2//xx2,phi2)
nn=[]
aa=[]
nn.append(q1)
nn.append(q2)
a1=gmpy2.powmod(c1,d1,p*q1)%q1
a2=gmpy2.powmod(c2,d2,p*q2)%q2
aa.append(a1)
aa.append(a2)
last=exCRT_getequation(aa,nn)#最终方程组 aa=n^14*q1*q2
new_e=7
new_phi=(q1-1)*(q2-1)
new_d=gmpy2.invert(new_e,new_phi)
m_2=gmpy2.powmod(last[0],new_d,last[1])#特解m_2
flag=gmpy2.iroot(m_2,2)[0]
import binascii
print(binascii.unhexlify(hex(flag)[2:]))

```

运行得到

```

1099358579338678297289853985632354554811203008593114217625408587627219550383101176094567633380822379070059373808
7315127935183160022527099534409653275027107080705198409752490095780942786144143679693401239370777001255660447906
5826879107677002380580866325868240270494148512743861326447181476633546419262340100453
b 'de1ctf{9b10a98b-71bb-4bdf-a6ff-f319943de21f}'

```

111.[UTCTF2020]hill

查看题目

```
wznqca{d4uqop0fk_q1nwofDbzg_eu}
```

这题就是个二阶希尔密码加密。知道密文前6位wznqca和明文前6位utflag，等于4个未知数6个一阶方程，居然还有解，那就是二阶跑不了了。求个逆乘回来就好。

```
from mpmath.libmp.backend import xrange

a00, a01, a10, a11 = 13, 6, 3, 21
dic = 'abcdefghijklmnopqrstuvwxyz'
c = 'wznqcaduqopfkqnwofdbzgeu'
m = ''
for i in xrange(0, len(c), 2):
    a, b = dic.index(c[i]), dic.index(c[i+1])
    a0, b0 = (a00*a+a01*b)%26, (a10*a+a11*b)%26
    m += dic[a0]
    m += dic[b0]
print(m)
```

记得把中间下划线、空格、数字之类的补回去，还有一个字母是大写别忘了
utflag{d4nger0us_c1pherText_qq}

112.[V&N2020 公开赛]CRT

[查看题目](#)

```
import hashlib
from functools import reduce
from Crypto.Util.number import *

ms = [getRandomNBitInteger(128) for i in range(8)]
p = reduce(lambda x,y: x*y, ms)
x = getRandomRange(1, p)
cs = [x % m for m in ms]

flag = "flag{" + hashlib.sha256(str(x).encode()).hexdigest() + "}"
# assert("4b93deeb" in flag)

# ms = [284461942441737992421992210219060544764, 218436209063777179204189567410606431578, 2886734381099336499112
76214358963643204, 239232622368515797881077917549177081575, 206264514127207567149705234795160750411, 33891554756
8169045185589241329271490503, 246545359356590592172327146579550739141, 219686182542160835171493232381209438048]
# cs = [273520784183505348818648859874365852523, 128223029008039086716133583343107528289, 5111091025406771271167
772696866083419, 33462335595116820423587878784664448439, 145377705960376589843356778052388633917, 12815842172585
6807614557926615949143594, 230664008267846531848877293149791626711, 94549019966480959688919233343793910003]
```

中国剩余定理，模不互质

根据assert里的说明这道题需要爆破

```

import hashlib
import gmpy2
import math
def merge(a1,n1,a2,n2):
    d = math.gcd(n1,n2)
    c = a2-a1
    if c%d!=0:
        return 0
    c = (c%n2+n2)%n2
    c = c//d
    n1 = n1//d
    n2 = n2//d
    c *= gmpy2.invert(n1,n2)
    c %= n2
    c *= n1*d
    c += a1
    global n3
    global a3
    n3 = n1*n2*d
    a3 = (c%n3+n3)%n3
    return 1
def exCRT(a,n):
    a1=a[0]
    n1=n[0]
    le= len(a)
    for i in range(1,le):
        a2 = a[i]
        n2=n[i]
        if not merge(a1,n1,a2,n2):
            return -1
        a1 = a3
        n1 = n3
    global mod
    mod=n1
    return (a1%n1+n1)%n1
ms = [284461942441737992421992210219060544764, 218436209063777179204189567410606431578, 288673438109933649911276
214358963643204, 239232622368515797881077917549177081575, 206264514127207567149705234795160750411, 3389155475681
69045185589241329271490503, 246545359356590592172327146579550739141, 219686182542160835171493232381209438048]
cs = [273520784183505348818648859874365852523, 128223029008039086716133583343107528289, 511109102540677127116777
2696866083419, 33462335595116820423587878784664448439, 145377705960376589843356778052388633917, 1281584217258568
07614557926615949143594, 230664008267846531848877293149791626711, 94549019966480959688919233343793910003]
x=exCRT(cs,ms)
flag=hashlib.sha256(str(x).encode()).hexdigest()
while "4b93deeb" not in flag:
    x=x+mod
    flag = hashlib.sha256(str(x).encode()).hexdigest()
print(flag)

```

python代码需要声明

运行得到

```
fa71921acdc2a756897d6b0c7ee41a7397386de2e7cde5b6adb525414b93deeb
```

113.[MRCTF2020]Easy_RSA

[查看题目](#)

```

import sympy
from gmpy2 import gcd, invert

```

```

from random import randint
from Crypto.Util.number import getPrime, isPrime, getRandomNBitInteger, bytes_to_long, long_to_bytes
import base64

from zlib import *
flag = b"MRCTF{XXXX}"
base = 65537

def gen_prime(N):
    A = 0
    while 1:
        A = getPrime(N)
        if A % 8 == 5:
            break
    return A

def gen_p():
    p = getPrime(1024)
    q = getPrime(1024)
    assert (p < q)
    n = p * q
    print("P_n = ", n)
    F_n = (p - 1) * (q - 1)
    print("P_F_n = ", F_n)
    factor2 = 2021 * p + 2020 * q
    if factor2 < 0:
        factor2 = (-1) * factor2
    return sympy.nextprime(factor2)

def gen_q():
    p = getPrime(1024)
    q = getPrime(1024)
    assert (p < q)
    n = p * q
    print("Q_n = ", n)
    e = getRandomNBitInteger(53)
    F_n = (p - 1) * (q - 1)
    while gcd(e, F_n) != 1:
        e = getRandomNBitInteger(53)
    d = invert(e, F_n)
    print("Q_E_D = ", e * d)
    factor2 = 2021 * p - 2020 * q
    if factor2 < 0:
        factor2 = (-1) * factor2
    return sympy.nextprime(factor2)

if __name__ == "__main__":
    _E = base
    _P = gen_p()
    _Q = gen_q()
    assert (gcd(_E, (_P - 1) * (_Q - 1)) == 1)
    _M = bytes_to_long(flag)
    _C = pow(_M, _E, _P * _Q)
    print("Ciphertext = ", _C)
'''
P_n = 140573321395373957012384636448279482040305765285585432834059669335099444446812575211087693039996799553714
7454621319605138680293634309296520251950411123857226982307219903981220810030193936508032851857870407676914748492
2508482686658959347725753762078590928561862163337382463252361958145933210306431342748775024336556028267742021320

```



```
8916817625436604684840186868658910731107573941540248335525588636715374910899570386483289737906923560147784203338
9670559525271151411747807282888019850618766792402026060012471724306742087636398053899410192943797866870912865258
7073901337310278665778299513763593234951137512120572797739181693
P_F_n = 1405733213953739570123846364482794820403057652855854328340596693350994444468125752110876930399967995537
1474546213196051386802936343092965202519504111238572269823072199039812208100301939365080328518578704076769147484
9225084826866589593477257537620785909285618621633373824632523619581459332103064313427487750240994273639673211101
2756203987901861608292693556795137818528088242690306459837666810661669462354007405721043279030957101877828172371
0994930151635857933293394780142192586806292968028305922173313521186946635709194350912242693822450297748434301924
950358561859804256788098033426537956252964976682327991427626735740
Q_n = 207142983381604497495453607436880188428772740545408520964594852839368023412713637661579761125250340043199
3805403493488086095696658505168448366253578062167331677484261470172644587063010919601667672518341287987046343227
7629916669130494040403733295593655306104176367902352484367520262917943100467697540593925707162162616635533550262
7188087462545994562865784091878951710157969919101238045298255195192783889104831338133309025301604489729260960839
9020824327454856123825300278947492073076000110404809329568059303332781882125530089342341219226581441854613401555
7579236219461780344469127987669565138930308525189944897421753947
Q_E_D = 1007720792222981345861161568507428178554081277169628919292598687466725726023339189580755826717524936182
5951828633612277270333018303722110505829865349079433788509849907358382183253279830951353838317523342953346734839
0389323225198805294950484802068148590902907221150968539067980432831310376368202773212266320112670699737501054831
6462865851422814192375722227139756468435550247318556885738341087118744061495400782537743497081580630557549328126
7578612370076828804844532619988098371750453882549810378930487368219105305036680682580260265867426844084457795549
9368404019114913934477160428428662847012289516655310680119638600315228284298935201
Ciphertext = 40855937355228438525361161524441274634175356845950884889338630813182607485910094677909779126550263
3041947960009043847754950009434240703963344358101265361653325654173367970366117733827283446871752530810475866028
3868502742829262155791451462902432479427577252201312646492699062014040641299948572875038587686811509173542557755
5027394033416643032644774339644654011686716639760512353355719065795222201167219831780961308225780478482467294410
8285434884122587644464948152387661857284544166918988594625320834372137931048237591473176136378814197875819207451
51430394526712790608442960106537539121880514269830696341737507717448946962021
'''
```

114.[MRCTF2020]Easy_RSA

[查看题目](#)

```
import sympy
from gmpy2 import gcd, invert
from random import randint
from Crypto.Util.number import getPrime, isPrime, getRandomNBitInteger, bytes_to_long, long_to_bytes
import base64

from zlib import *
flag = b"MRCTF{XXXX}"
base = 65537

def gen_prime(N):
    A = 0
    while 1:
        A = getPrime(N)
        if A % 8 == 5:
            break
    return A

def gen_p():
    p = getPrime(1024)
    q = getPrime(1024)
    assert (p < q)
    n = p * q
    print("P_n = ", n)
    F_n = (p - 1) * (q - 1)
```

```

print("P_F_n = ", F_n)
factor2 = 2021 * p + 2020 * q
if factor2 < 0:
    factor2 = (-1) * factor2
return sympy.nextprime(factor2)

def gen_q():
    p = getPrime(1024)
    q = getPrime(1024)
    assert (p < q)
    n = p * q
    print("Q_n = ", n)
    e = getRandomNBitInteger(53)
    F_n = (p - 1) * (q - 1)
    while gcd(e, F_n) != 1:
        e = getRandomNBitInteger(53)
    d = invert(e, F_n)
    print("Q_E_D = ", e * d)
    factor2 = 2021 * p - 2020 * q
    if factor2 < 0:
        factor2 = (-1) * factor2
    return sympy.nextprime(factor2)

if __name__ == "__main__":
    _E = base
    _P = gen_p()
    _Q = gen_q()
    assert (gcd(_E, (_P - 1) * (_Q - 1)) == 1)
    _M = bytes_to_long(flag)
    _C = pow(_M, _E, _P * _Q)
    print("Ciphertext = ", _C)
'''
P_n = 140573321395373957012384636448279482040305765285585432834059669335099444446812575211087693039996799553714
7454621319605138680293634309296520251950411123857226982307219903981220810030193936508032851857870407676914748492
2508482686658959347725753762078590928561862163337382463252361958145933210306431342748775024336556028267742021320
8916817625436604684840186868658910731107573941540248335525588636715374910899570386483289737906923560147784203338
9670559525271151411747807282888019850618766792402026060012471724306742087636398053899410192943797866870912865258
7073901337310278665778299513763593234951137512120572797739181693
P_F_n = 1405733213953739570123846364482794820403057652855854328340596693350994444468125752110876930399967995537
1474546213196051386802936343092965202519504111238572269823072199039812208100301939365080328518578704076769147484
9225084826866589593477257537620785909285618621633373824632523619581459332103064313427487750240994273639673211101
2756203987901861608292693556795137818528088242690306459837666810661669462354007405721043279030957101877828172371
0994930151635857933293394780142192586806292968028305922173313521186946635709194350912242693822450297748434301924
950358561859804256788098033426537956252964976682327991427626735740
Q_n = 207142983381604497495453607436880188428772740545408520964594852839368023412713637661579761125250340043199
3805403493488086095696658505168448366253578062167331677484261470172644587063010919601667672518341287987046343227
7629916669130494040403733295593655306104176367902352484367520262917943100467697540593925707162162616635533550262
7188087462545994562865784091878951710157969919101238045298255195192783889104831338133309025301604489729260960839
902082432745485612382530027894749207307600011040480932956805930332781882125530089342341219226581441854613401555
7579236219461780344469127987669565138930308525189944897421753947
Q_E_D = 1007720792222981345861161568507428178554081277169628919292598687466725726023339189580755826717524936182
5951828633612277270333018303722110505829865349079433788509849907358382183253279830951353838317523342953346734839
0389323225198805294950484802068148590902907221150968539067980432831310376368202773212266320112670699737501054831
6462865851422814192375722227139756468435550247318556885738341087118744061495400782537743497081580630557549328126
7578612370076828804844532619988098371750453882549810378930487368219105305036680682580260265867426844084457795549
9368404019114913934477160428428662847012289516655310680119638600315228284298935201
Ciphertext = 40855937355228438525361161524441274634175356845950884889338630813182607485910094677909779126550263

```

```
3041947960009043847754950009434240703963344358101265361653325654173367970366117733827283446871752530810475866028
3868502742829262155791451462902432479427577252201312646492699062014040641299948572875038587686811509173542557755
502739403341664303264477433964465401168671663976051235335571906579522201167219831780961308225780478482467294410
8285434884122587644464948152387661857284544166918988594625320834372137931048237591473176136378814197875819207451
51430394526712790608442960106537539121880514269830696341737507717448946962021
'''
```

脚本如下

```
P_n = 140573321395373957012384636448279482040305765285585432834059669335099444446812575211087693039996799553714
7454621319605138680293634309296520251950411123857226982307219903981220810030193936508032851857870407676914748492
2508482686658959347725753762078590928561862163337382463252361958145933210306431342748775024336556028267742021320
8916817625436604684840186868658910731107573941540248335525588636715374910899570386483289737906923560147784203338
9670559525271151411747807282888019850618766792402026060012471724306742087636398053899410192943797866870912865258
7073901337310278665778299513763593234951137512120572797739181693
P_F_n = 1405733213953739570123846364482794820403057652855854328340596693350994444468125752110876930399967995537
1474546213196051386802936343092965202519504111238572269823072199039812208100301939365080328518578704076769147484
9225084826866589593477257537620785909285618621633373824632523619581459332103064313427487750240994273639673211101
2756203987901861608292693556795137818528088242690306459837666810661669462354007405721043279030957101877828172371
0994930151635857933293394780142192586806292968028305922173313521186946635709194350912242693822450297748434301924
950358561859804256788098033426537956252964976682327991427626735740
Q_n = 20714298338160449749545360743688018842877274054540852096459485283936802341271363766157976112525034004319938
0540349348808609569665850516844836625357806216733167748426147017264458706301091960166767251834128798704634322776
2991666913049404040373329559365530610417636790235248436752026291794310046769754059392570716216261663553355026271
8808746254599456286578409187895171015796991910123804529825519519278388910483133813330902530160448972926096083990
2082432745485612382530027894749207307600011040480932956805930333278188212553008934234121922658144185461340155575
79236219461780344469127987669565138930308525189944897421753947
Q_E_D = -100772079222298134586116156850742817855408127716962891929259868746672572602333918958075582671752493618259
5182863361227727033301830372211050582986534907943378850984990735838218325327983095135383831752334295334673483903
8932322519880529495048480206814859090290722115096853906798043283131037636820277321226632011267069973750105483164
628658514228141923757222713975646843555024731855688573834108711874406149540078253774349708158063055754932812675
7861237007682880484453261998809837175045388254981037893048736821910530503668068258026026586742684408445779554993
68404019114913934477160428428662847012289516655310680119638600315228284298935201
Ciphertext = 40855937355228438525361161524441274634175356845950884889338630813182607485910094677909779126550263
3041947960009043847754950009434240703963344358101265361653325654173367970366117733827283446871752530810475866028
3868502742829262155791451462902432479427577252201312646492699062014040641299948572875038587686811509173542557755
502739403341664303264477433964465401168671663976051235335571906579522201167219831780961308225780478482467294410
8285434884122587644464948152387661857284544166918988594625320834372137931048237591473176136378814197875819207451
51430394526712790608442960106537539121880514269830696341737507717448946962021
import gmpy2
import sympy
base = 65537
def getp(n, phi):
    p_q = n + 1 - phi
    p_q_2 = p_q ** 2 - 4 * n
    p_q_2 = gmpy2.iroot(p_q_2, 2)[0]
    q = (p_q_2 + p_q) // 2
    p = p_q - q
    factor2 = 2021 * p + 2020 * q
    if factor2 < 0:
        factor2 = (-1) * factor2
    P = sympy.nextprime(factor2)
    return P
def getq(n, phi):
    p_q = n + 1 - phi
    p_q_2 = p_q ** 2 - 4 * n
    p_q_2 = gmpy2.iroot(p_q_2, 2)[0]
    q = (p_q_2 + p_q) // 2
    p = p_q - q
```

```

factor2 = 2021 * p - 2020 * q
if factor2 < 0:
    factor2 = (-1) * factor2
Q = sympy.nextprime(factor2)
return Q
p=getp(P_n,P_F_n)
K=((Q_E_D-1)//Q_n)+1
phi=(Q_E_D-1)//K
q=getq(Q_n,phi)
e=base
d=gmpy2.invert(e,(p-1)*(q-1))
M=gmpy2.powmod(Ciphertext,d,p*q)
import binascii
print(binascii.unhexlify(hex(M)[2:]))

#b'MRCTF{Ju3t@_31mp13_que3t10n}'

```

115.[NPUCTF2020]EzRSA

[查看题目](#)

```

from gmpy2 import lcm , powmod , invert , gcd , mpz
from Crypto.Util.number import getPrime
from sympy import nextprime
from random import randint
p = getPrime(1024)
q = getPrime(1024)
n = p * q
gift = lcm(p - 1 , q - 1)
e = 54722
flag = b'NPUCTF{*****}'
m = int.from_bytes(flag , 'big')
c = powmod(m , e , n)
print('n: ' , n)
print('gift: ' , gift)
print('c: ' , c)

#n: 17083941230213489700426636484487738282426471494607098847295335339638177583685457921198569105417734668692072
7277591393582076672487039524366801831533276061474219323658899833472820464391561766857651436206371073478704019469
4650162053166557366806834908041080799658229750588994620505287900202893612531531225647058362291364631977912555969
1270916064588684997382451412747432722966919513413709987353038375477178385125453567111965259721484997156799355617
6421315690958103040771310535884830572443407427518049354940876873634169213140415470931185657676096670338595831252
75322077617576783247853718516166743858265291135353895239981121
#gift: 21354926537766862125533295605609672853033089368258873559119169174547721979606822401498211381772168335865
0909096989241977595840608799405458502289416595076842774154573624791841025580489452208572064295257963841848380024
3368312702566458196708508543635051350999572787188236243275631609875253617015664414032058822919469443284453403064
0762327650242484355433265974188517515863085145401245713091527875597129502093578255768961322780451121779102660197
4101399510657948486876825108445333841711548351513286959471216205236208341416395468130625913705758103665744189742
8432575924018950961141822554251369262248368899977337886190114104
#c: 37389606391947379576676841435650055035962764516179224746697455292999293955079714353111815783872233234293232
8692737057695507861833575750816126358516412604754541302882987326934292409233929895763507973644685183741435775731
2525158356579607212496060244403765822636515347192211817658170822313646743520831977673861869637519843133863288550
0583594294550526763231967282804085086145279530572147791654503565778203788104675270063772961941026713603020599018
9797733972829234513282718422715506132632858564001991632884737229575447283231825863605466309147580123505065740185
7262960415898483713074139212596685365780269667500271108538319

```

由于gift较大，直接遍历

得到p,q后，变形 $c = me \pmod n \Rightarrow c = (m^2)e/2 \pmod n$ (因为e和φ(n)不互质)

得到m² 后尝试直接开平方看是否默认最小值为解

```
gift=21354926537766862125533295605609672853033089368258873559119169174547721979606822401498211381772168335865090
9096989241977595840608799405458502289416595076842774154573624791841025580489452208572064295257963841848380024336
8312702566458196708508543635051350999572787188236243275631609875253617015664414032058822919469443284453403064076
2327650242484355433265974188517515863085145401245713091527875597129502093578255768961322780451121779102660197410
1399510657948486876825108445333841711548351513286959471216205236208341416395468130625913705758103665744189742843
2575924018950961141822554251369262248368899977337886190114104
import gmpy2
e=54722
c=37389606391947379576676841435650055035962764516179224746697455292999293955079714353111815783872233234293232869
2737057695507861833575750816126358516412604754541302882987326934292409233929895763507973644685183741435775731252
5158356579607212496060244403765822636515347192211817658170822313646743520831977673861869637519843133863288550058
3594294550526763231967282804085086145279530572147791654503565778203788104675270063772961941026713603020599018979
7733972829234513282718422715506132632858564001991632884737229575447283231825863605466309147580123505065740185726
2960415898483713074139212596685365780269667500271108538319
n=17083941230213489700426636484487738282426471494607098847295335339638177583685457921198569105417734668692072727
7591393582076672487039524366801831533276061474219323658899833472820464391561766857651436206371073478704019469465
0162053166557366806834908041080799658229750588994620505287900202893612531531225647058362291364631977912555969127
0916064588684997382451412747432722966919513413709987353038375477178385125453567111965259721484997156799355617642
1315690958103040771310535884830572443407427518049354940876873634169213140415470931185657676096670338595831252753
22077617576783247853718516166743858265291135353895239981121
for i in range(100):
    phi=gift*i
    pq=n+1-phi
    p_q=gmpy2.iroot(pq**2-4*n,2)[0]
    if gmpy2.is_prime((pq+p_q)//2):
        p=(pq+p_q)//2
        q=pq-p
        break
new_e=e//2
new_d=gmpy2.invert(new_e,(p-1)*(q-1))
m_2=gmpy2.powmod(c,new_d,n)
m=gmpy2.iroot(m_2,2)[0]
import binascii
print(binascii.unhexlify(hex(m)[2:]))
```

运行得到 `b'NPUCTF{diff1cult_rsa_1s_e@sy}'`

116.[AFCTF2018]花开藏宝地

查看题目 一个zip压缩包一个txt

第80804238007977405688648566160504278593148666302626415149704905628622876270862865768337953835725801963142685182510812938072115996355782396318303927020705623120652014080032809421180400984242061592520733710243483947230962631945045134540159517488288781666622635328316972979183761952842010806304748313326215619695085380586052550443025074501971925005072999275628549710915357400946408857号藏宝图

我把我的宝藏都藏在了那里！

那个神秘的地方！

于是我把藏宝图分成了5份，交给五位贤者让他们帮我妥善保管，并且只要搜集3份就可以获得宝藏的地址。

第一位贤者将藏宝图放进时空门中说道：

“那么口令就是我的生日吧，那可是个好数字呢。”

第二位贤者将藏宝图放进宝箱，【小】声念着自己的名字锁上了宝箱。

第三位贤者将藏宝图施上咒语丢进大海：“只要【大】声喊出那句咒语就可以把水驱逐！”

第四位贤者找了个破锁锁上了宝箱，狡黠地笑着：“谁知道它是坏的呢？”

第五位贤者给藏宝图裹上了隐身衣，放入了一个匣子里

据说，只有拥有【智慧】与【力量】就可以获得宝藏了呢！~

你是这样的勇者吗？

那么我们就直接爆破第一个zip的密码

过程不写了

- secret1是爆破生日数字：19260817
- secret2爆破英文字母：alice
- secret4伪加密
- secret5 ntfs隐写

得到四组x和m，题目说只要3组就行了

应该是e=3，然后共模攻击。即明文m不变，给出改变的c和n。可是如果是这样，为什么还需要提示一个大数？

而且实际运算以后，并无法得到flag

百度了一下，原来这是标准的门限秘密共享方案(threshold secret sharing scheme),简称门限方案：

<http://www.matrix67.com/blog/archives/1261>

假设公司董事会共五个人，每个人保存密钥的一部分。要求三个人在场就可以拿到密钥打开保险箱，而且保险箱打开后，无法知道到底是哪三个人提供的密钥。

和这题提示正好对应：

于是我把藏宝图分成了5份，交给五位贤者让他们帮我妥善保管，并且只要搜集3份就可以获得宝藏的地址。

门限加密有多种方案，举个例子，三个平面能确定一个点。而有无数平面通过同一个点。

题目标题为花开，即Asm nth-Bloom方案

利用的就是中国剩余数定理，可以给出很多组n和c，满足 $m \bmod n = c$ ，然后根据其中的几组就可以找到解，但要注意crt的有多解，只要是 $m+kn_1n_2\dots n_n$ 的都满足。

这题同理：

```

from Crypto.Util.number import *

z=80804238007977405688648566160504278593148666302626415149704905628622876270862865768337953835725801963142685182
5108129380721159963557823963183039270207056231206520140800328094211804009842420615925207337102434839472309626319
4504513454015951748828878166662263532831697297918376195284201080630474831332621561969508538058605255044302507450
1971925005072999275628549710915357400946408857

x5 = 23050206438294728234366015979161193669652080797036113946960345868931128604151676787590354926386195074077870
5012699983268093626403307298415066249636346303539570207577050391796770068203937723627361951969413683246596072925
6926703654909708478252695810044839642614919176807590917916537595142131887784019686764332847537810067382937524401
86858616315727565803777032119737689210471541053061940547213
m5 = 3470515596224631445396699500966581634256464114357976919737015137257015751008104461758494240000007585507043
0240507732735393411493866540572679626172742301366146501862670272443070970511943485865887494229487420503750457974
2628020537220939051262353403802618285935084556216673099463617055306679574847319291518755274894784493611986483106
84702574627199321092927111137398333029697068474762820822249
x4 = 10045977991352054009806540742062995481667792642335676952475907263221910615584945012518520555749113835776049
4272691949199099803239098119602186117878931534968435982565071570831032814288620974807498206233914826253433847572
7034076787129650983201225497595795663163722209596108145739456980839095750053032532056532442385423002664605597906
06278310650849881421791081944960157781855164700773081375247
m4 = 3470515596224631445396699500966581634256464114357976919737015137257015751008104461758494240000007585507043
0240507732735393411493866540572679626172742301366146501862670272443070970511943485865887494229487420503750457974
2628020537220939051262353403802618285935084556216673099463617055306679574847319291518755274894784493611986483106
84702574627199321092927111137398333029697068474762820820091
x2 = 15201268127068234005169062792458623270255246081003032226782740177130490746980259186191292128183389061318631
7787813611372838066924894691892444503039545946728621696590087591246339208248647926966446848123290344911662916758
0391348174047205124658178672552774767173534395052432475681261933615580429403522040933812604024007394290502805262
12446967632582771424597203000629197487733610187359662268583
m2 = 3470515596224631445396699500966581634256464114357976919737015137257015751008104461758494240000007585507043
0240507732735393411493866540572679626172742301366146501862670272443070970511943485865887494229487420503750457974
2628020537220939051262353403802618285935084556216673099463617055306679574847319291518755274894784493611986483106
84702574627199321092927111137398333029697068474762820818553
x2 = 15201268127068234005169062792458623270255246081003032226782740177130490746980259186191292128183389061318631
7787813611372838066924894691892444503039545946728621696590087591246339208248647926966446848123290344911662916758
0391348174047205124658178672552774767173534395052432475681261933615580429403522040933812604024007394290502805262
12446967632582771424597203000629197487733610187359662268583
m2 = 3470515596224631445396699500966581634256464114357976919737015137257015751008104461758494240000007585507043
0240507732735393411493866540572679626172742301366146501862670272443070970511943485865887494229487420503750457974
2628020537220939051262353403802618285935084556216673099463617055306679574847319291518755274894784493611986483106
84702574627199321092927111137398333029697068474762820818553

x1 = 30534513391139521857379090350829623865914780227403179664301753901164880280876316290233564419564852537551894
1848430114497150082025133000033835083076541927530829557051524161069423494451667848236452337271862085346869364976
9890471805321675607964700675499153907732712079015378472138824799973255752786729176484178687590771509990448910992
06133296336190476413164240995177077671480352739572539631359
m1 = 3470515596224631445396699500966581634256464114357976919737015137257015751008104461758494240000007585507043
0240507732735393411493866540572679626172742301366146501862670272443070970511943485865887494229487420503750457974
2628020537220939051262353403802618285935084556216673099463617055306679574847319291518755274894784493611986483106
84702574627199321092927111137398333029697068474762820813413
c=[x1,x2,x4]
n=[m1,m2,m4]
a=crt(c,n)
#通过这一步来去除掉kn1n2n3的干扰，实际上这题目中给你的就是n1*n2*n3*n4*n5.
r=a%z
print(long_to_bytes(r))
#b"A treasure map is a map that marks the location of buried treasure, a lost mine, a valuable secret or a hidden locale. So flag is afctf{1sn't_s0_int3Resting}."

```

查看题目

```
p = getPrime(1024)
q = getPrime(1024)
N = p * q
g, r1, r2 = [getRandomRange(1, N) for _ in range(3)]
g1 = pow(g, r1 * (p-1), N)
g2 = pow(g, r2 * (q-1), N)
def encrypt(m):
    s1, s2 = [getRandomRange(1, N) for _ in range(2)]
    c1 = (m * pow(g1, s1, N)) % N
    c2 = (m * pow(g2, s2, N)) % N
    return (c1, c2)
def decrypt(c1, c2):
    xp = c1 % p
    xq = c2 % q
    # Chinese Remainder Theorem
    m = (xp*inverse(q, p)*q + xq*inverse(p, q)*p) % N
    return m
```

注意到后面给了一个解密方程，所以只需要求出 p 、 q ，放进去就可以解密。

考虑encrypt中给了 $c1/c2$ ，那按逻辑来说，就是通过 $g1/g2$ 的值，求 p/q 。带入解密，得到flag:

$$g^{r1(p-1)} \equiv g1 \pmod N$$

$$g1 = kN + g^{r1(p-1)}$$

两边都 $\pmod p$

$$g1 \pmod p = g^{r1(p-1)} \pmod p$$

$$\text{费马小定理有: } a^{p-1} \equiv 1 \pmod p$$

$$\text{故 } g1 \pmod p = g^{r1(p-1)} \pmod p = (g^{(p-1)} \pmod p)^{r1} \pmod p = 1^{r1} \pmod p = 1$$

$$g1 \pmod p = 1$$

$$\text{同理 } g2 \pmod q = 1$$

所以

```
g1 - 1 = x*p
gcd(n, g1-1) == p
g2 - 1 = y*q
gcd(n, g2-1) == q
#事实上好像gcd(n, g1-1)并不一定等于p, 有可能是k*p。所以先求出来, 至少公因数会是n的因子
#本题正好求出是p和q
```

所以通过最大公约数可以求得 pq 。

再看encrypt:

```
def encrypt(m):
    s1, s2 = [getRandomRange(1, N) for _ in range(2)]
    c1 = (m * pow(g1, s1, N)) % N
    c2 = (m * pow(g2, s2, N)) % N
    return (c1, c2)
```

直接套用decrypt可以得到flag:


```

from Crypto.Util.number import *
p=gcd((g1-1),N)
q=gcd((g2-1),N)
phi=(p-1)*(q-1)
xp = c1 % p
xq = c2 % q
m = (xp*inverse_mod(q, p)*q + xq*inverse_mod(p, q)*p) % N
long_to_bytes(m)
#b'flag{1CE9514E-12AF-49BE-B002-6A3D7E6078FA}\x00\x0bob\x0f8'

```

118.[BJDCTF2020]伏羲六十四卦

查看题目

```

from secret import flag

def encrypt5():
    enc=''
    for i in flag:
        enc+=chr((a*(ord(i)-97)+b)%26+97)
    return(enc)

def encrypt4():
    temp=''
    offset=5
    for i in range(len(enc)):
        temp+=chr(ord(enc[i])-offset-i)
    return(temp)

```

这是什么，怎么看起来像是再算64卦!!!

密文:升随临损巽睽萃小过讼艮颐小过震蛊屯未济中孚艮困恒晋升损蛊萃未济巽解艮贲未济观豫损蛊晋噬嗑晋旅解大畜困未济随蒙升解睽未济井困未济旅萃未济震蒙未济师涣归妹大有

嗯? 为什么还有个b呢?

b=7

flag: 请按照格式BJD{}

脚本如下

```

def decrypt4(enc):
    temp=''
    offset=5
    for i in range(len(enc)):
        temp+=chr(ord(enc[i])+offset+i)
    return temp
def decrypt5(flag):
    for a in range(1,200):
        enc = ''
        for i in flag:
            for k in range(200):
                if (ord(i) - 97 - 7+26*k)%a==0:
                    enc+= chr((ord(i) - 97 - 7 + 26 * k) // a + 97)
                    break
        print(enc)

s='升随临损巽睽颐萃小过讼艮颐小过震蛊屯未济中孚艮困恒晋升损蛊萃蛊未济巽解艮贲未济观豫损蛊晋噬嗑晋旅解大畜困未济随蒙升解睽未济井困未济旅萃未济震蒙未济师涣归妹大有'
dic={'坤': '000000', '剥': '000001', '比': '000010', '观': '000011', '豫': '000100', '晋': '000101', '萃': '000110', '否': '000111', '谦': '001000', '艮': '001001', '蹇': '001010', '渐': '001011', '小过': '001100', '旅': '001101', '咸': '001110', '遁': '001111', '师': '010000', '蒙': '010001', '坎': '010010', '涣': '010011', '解': '010100', '未济': '010101', '困': '010110', '讼': '010111', '升': '011000', '蛊': '011001', '井': '011010', '巽': '011011', '恒': '011100', '鼎': '011101', '大过': '011110', '姤': '011111', '复': '100000', '颐': '100001', '屯': '100010', '益': '100011', '震': '100100', '噬嗑': '100101', '随': '100110', '无妄': '100111', '明夷': '101000', '贲': '101001', '既济': '101010', '家人': '101011', '丰': '101100', '离': '101101', '革': '101110', '同人': '101111', '临': '110000', '损': '110001', '节': '110010', '中孚': '110011', '归妹': '110100', '睽': '110101', '兑': '110110', '履': '110111', '泰': '111000', '大畜': '111001', '需': '111010', '小畜': '111011', '大壮': '111100', '大有': '111101', '夬': '111110', '乾': '111111'}
li=[]
k=0
for i in range(len(s)):
    if k ==1:
        k=0
        continue
    try:
        li.append(dic[s[i]])
    except:
        t=''
        t=t+s[i]+s[i+1]
        li.append(dic[t])
        k=1
ss=''.join(li)
print(ss)
enc=''
for i in range(0,len(ss),8):
    enc+=chr(eval('0b'+ss[i:i+8]))
import base64
print(enc)
x=base64.b64decode(enc).decode()
print(x)
x=decrypt4(x)
x=decrypt5(x)

```

bjdcongratulationsongettingtheflag

运行得到

flag包裹提交即可

119.[ACTF新生赛2020]crypto-aes

查看题目 91144196586662942563895769614300232343026691029427747065707381728622849079757 b'\x8c-\xcd\xde\xa7\xe9\x7f.b\x8aKs\xf1\xba\xc75\xc4d\x13\x07\xac\xa4&\xd6\x91\xfe\xf3\x14\x10|\xf8p'

```
from Cryptodome.Cipher import AES
import os
import gmpy2
from flag import FLAG
from Cryptodome.Util.number import *

def main():
    key=os.urandom(2)*16
    iv=os.urandom(16)
    print(bytes_to_long(key)^bytes_to_long(iv))
    aes=AES.new(key,AES.MODE_CBC,iv)
    enc_flag = aes.encrypt(FLAG)
    print(enc_flag)
if __name__=="__main__":
    main()
```

key是两个字节不断重复得到的，因此结合输出的与iv向量的异或值很容易的到key和iv

脚本如下

```
key_iv=91144196586662942563895769614300232343026691029427747065707381728622849079757
flag_encrypt=b'\x8c-\xcd\xde\xa7\xe9\x7f.b\x8aKs\xf1\xba\xc75\xc4d\x13\x07\xac\xa4&\xd6\x91\xfe\xf3\x14\x10|\xf8p'
#print(hex(key_iv))
key=hex(key_iv)[2:6]*16
iv=key_iv^eval('0x'+key)
import Crypto.Util.number
iv=Crypto.Util.number.long_to_bytes(iv)
key=Crypto.Util.number.long_to_bytes(eval('0x'+key))
import Crypto.Cipher.AES
decrypt=Crypto.Cipher.AES.new(key,Crypto.Cipher.AES.MODE_CBC,iv)
print(decrypt.decrypt(flag_encrypt))
```

运行得到 b'actf{w0w_y0u_can_so1v3_AES_now!}'

120.[NCTF2019]Sore

查看题目

```
91144196586662942563895769614300232343026691029427747065707381728622849079757
b'\x8c-\xcd\xde\xa7\xe9\x7f.b\x8aKs\xf1\xba\xc75\xc4d\x13\x07\xac\xa4&\xd6\x91\xfe\xf3\x14\x10|\xf8p'
```

