




2018湖湘杯逆向（Reverse）

原创

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订阅专栏

0x00: 介绍

HighwayHash64这道题放了很久一直没有去做, 最近看了一下发现这道题比较特殊, 重新整理了一下思路, Replace这道题比较简单, 属于签到题难度。

0x01: 题目

题目1: Replace

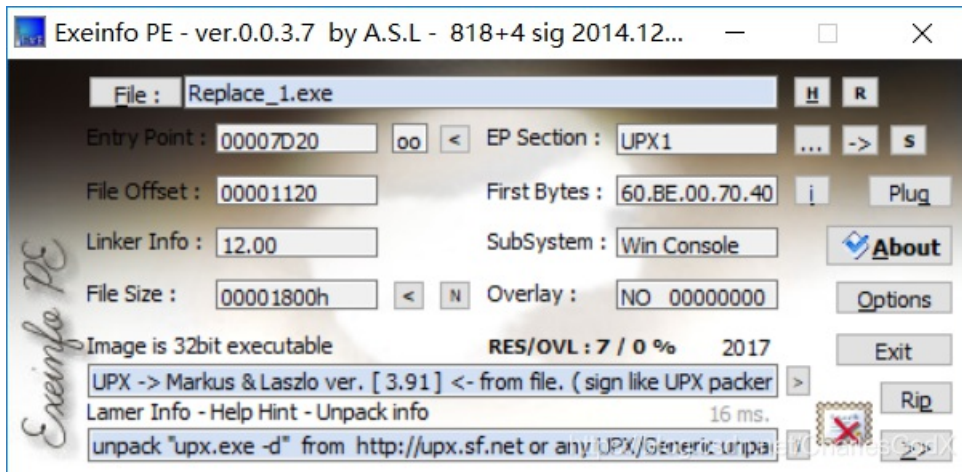
题目链接:

简单的...密码学~

http://hxb2018.oss-cn-beijing.aliyuncs.com/reserves/Replace_B21DA8B2F172C13764989DF0F99B890A.rar

解题过程:

首先查壳是UPX壳，我们直接用脱壳机就可以了（当然也可以手动）

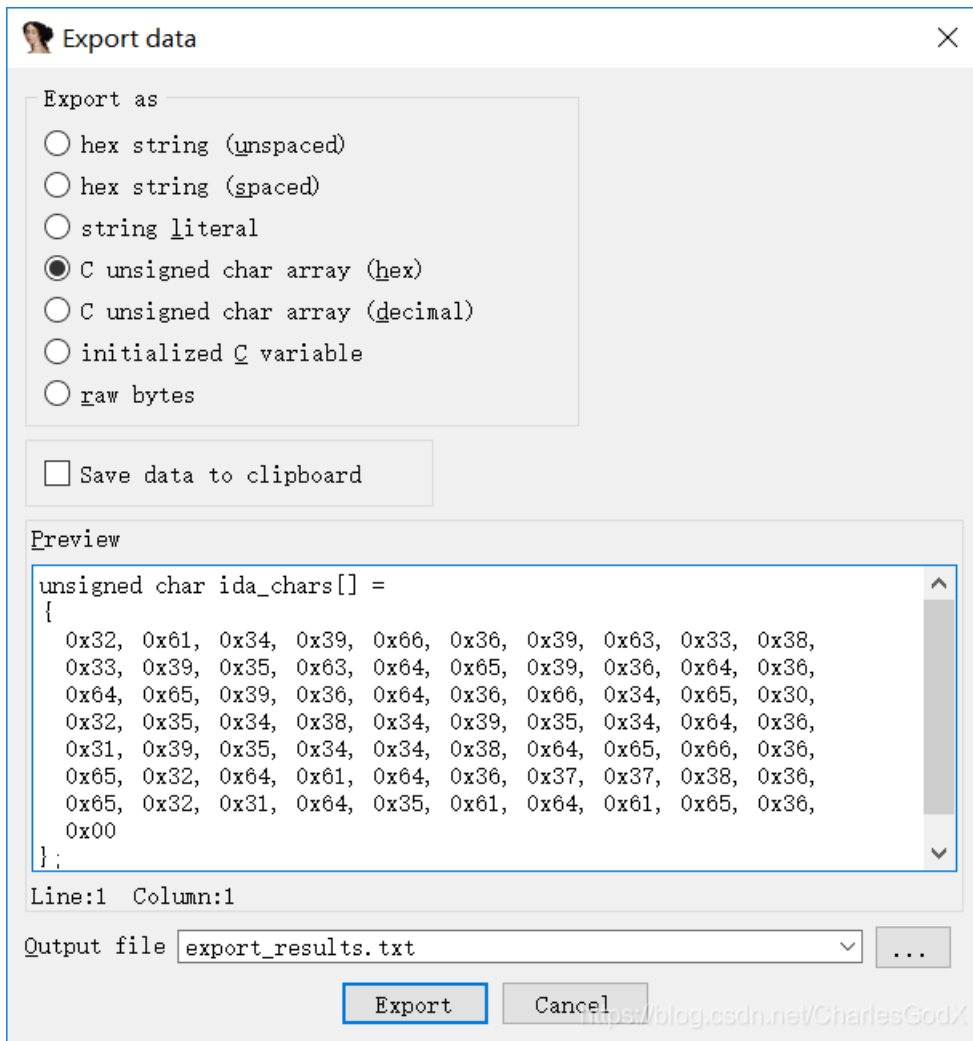


我们用IDA分析后发现，就是一个简单的算法

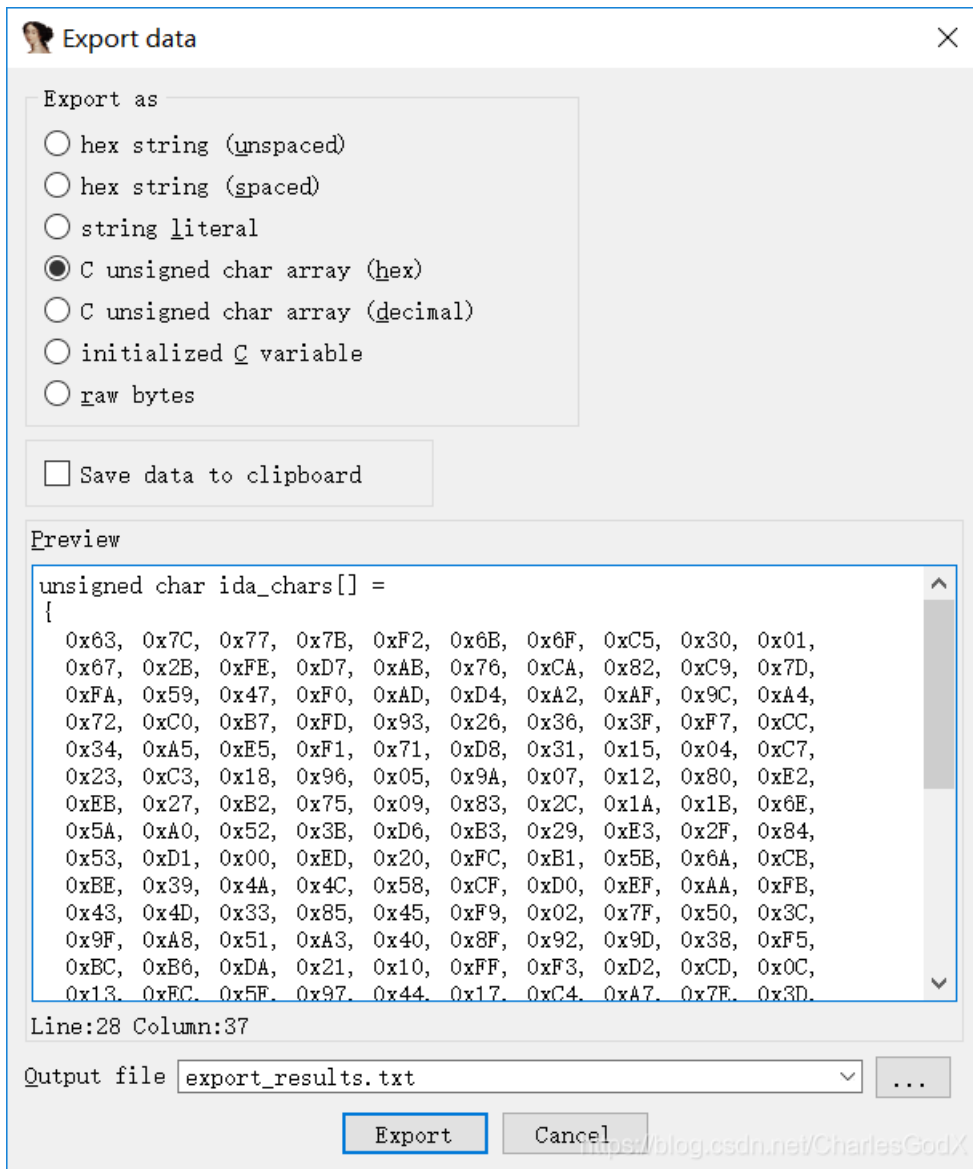
```
14 buf = a1;
15 if ( strlen != 35 )
16     return -1;
17 v4 = 0;
18 while ( 1 )
19 {
20     v5 = *(v4 + buf);
21     v6 = (v5 >> 4) % 16;
22     v7 = (16 * v5 >> 4) % 16;
23     v8 = byte_402150[2 * v4];
24     if ( v8 < 48 || v8 > 57 )
25         v9 = v8 - 87;
26     else
27         v9 = v8 - 48;
28     v10 = byte_402151[2 * v4];
29     v11 = 16 * v9;
30     if ( v10 < 48 || v10 > 57 )
31         v12 = v10 - 87;
32     else
33         v12 = v10 - 48;
34     if ( byte_4021A0[16 * v6 + v7] != ((v11 + v12) ^ 0x19) )
35         break;
36     if ( ++v4 >= 35 )
37         return 1;
38 }
39 return -1;
```

<https://blog.csdn.net/CharlesGodX>

我们提取出第一处关键数组，这里可以不包括最后一个0x00:



第二处关键数组:



提取完之后，就可以写爆破脚本了：

```
a = [
    0x63, 0x7C, 0x77, 0x7B, 0xF2, 0x6B, 0x6F, 0xC5, 0x30, 0x01,
    0x67, 0x2B, 0xFE, 0xD7, 0xAB, 0x76, 0xCA, 0x82, 0xC9, 0x7D,
    0xFA, 0x59, 0x47, 0xF0, 0xAD, 0xD4, 0xA2, 0xAF, 0x9C, 0xA4,
    0x72, 0xC0, 0xB7, 0xFD, 0x93, 0x26, 0x36, 0x3F, 0xF7, 0xCC,
    0x34, 0xA5, 0xE5, 0xF1, 0x71, 0xD8, 0x31, 0x15, 0x04, 0xC7,
    0x23, 0xC3, 0x18, 0x96, 0x05, 0x9A, 0x07, 0x12, 0x80, 0xE2,
    0xEB, 0x27, 0xB2, 0x75, 0x09, 0x83, 0x2C, 0x1A, 0x1B, 0x6E,
    0x5A, 0xA0, 0x52, 0x3B, 0xD6, 0xB3, 0x29, 0xE3, 0x2F, 0x84,
    0x53, 0xD1, 0x00, 0xED, 0x20, 0xFC, 0xB1, 0x5B, 0x6A, 0xCB,
    0xBE, 0x39, 0x4A, 0x4C, 0x58, 0xCF, 0xD0, 0xEF, 0xAA, 0xFB,
    0x43, 0x4D, 0x33, 0x85, 0x45, 0xF9, 0x02, 0x7F, 0x50, 0x3C,
    0x9F, 0xA8, 0x51, 0xA3, 0x40, 0x8F, 0x92, 0x9D, 0x38, 0xF5,
    0xBC, 0xB6, 0xDA, 0x21, 0x10, 0xFF, 0xF3, 0xD2, 0xCD, 0x0C,
    0x13, 0xEC, 0x5F, 0x97, 0x44, 0x17, 0xC4, 0xA7, 0x7E, 0x3D,
    0x64, 0x5D, 0x19, 0x73, 0x60, 0x81, 0x4F, 0xDC, 0x22, 0x2A,
    0x90, 0x88, 0x46, 0xEE, 0xB8, 0x14, 0xDE, 0x5E, 0x0B, 0xDB,
    0xE0, 0x32, 0x3A, 0x0A, 0x49, 0x06, 0x24, 0x5C, 0xC2, 0xD3,
    0xAC, 0x62, 0x91, 0x95, 0xE4, 0x79, 0xE7, 0xC8, 0x37, 0x6D,
    0x8D, 0xD5, 0x4E, 0xA9, 0x6C, 0x56, 0xF4, 0xEA, 0x65, 0x7A,
    0xAE, 0x08, 0xBA, 0x78, 0x25, 0x2E, 0x1C, 0xA6, 0xB4, 0xC6,
    0xE8, 0xDD, 0x74, 0x1F, 0x4B, 0xBD, 0x8B, 0x8A, 0x70, 0x3E,
```

```

0xB5, 0x66, 0x48, 0x03, 0xF6, 0x0E, 0x61, 0x35, 0x57, 0xB9,
0x86, 0xC1, 0x1D, 0x9E, 0xE1, 0xF8, 0x98, 0x11, 0x69, 0xD9,
0x8E, 0x94, 0x9B, 0x1E, 0x87, 0xE9, 0xCE, 0x55, 0x28, 0xDF,
0x8C, 0xA1, 0x89, 0x0D, 0xBF, 0xE6, 0x42, 0x68, 0x41, 0x99,
0x2D, 0x0F, 0xB0, 0x54, 0xBB, 0x16
]
b = [
0x32, 0x61, 0x34, 0x39, 0x66, 0x36, 0x39, 0x63, 0x33, 0x38,
0x33, 0x39, 0x35, 0x63, 0x64, 0x65, 0x39, 0x36, 0x64, 0x36,
0x64, 0x65, 0x39, 0x36, 0x64, 0x36, 0x66, 0x34, 0x65, 0x30,
0x32, 0x35, 0x34, 0x38, 0x34, 0x39, 0x35, 0x34, 0x64, 0x36,
0x31, 0x39, 0x35, 0x34, 0x34, 0x38, 0x64, 0x65, 0x66, 0x36,
0x65, 0x32, 0x64, 0x61, 0x64, 0x36, 0x37, 0x37, 0x38, 0x36,
0x65, 0x32, 0x31, 0x64, 0x35, 0x61, 0x64, 0x61, 0x65, 0x36,
0x00
]

flag = ""

v4 = 0
while(1):
    v8 = b[2 * v4]
    if (v8 < 48 | v8 > 57):
        v9 = v8 - 87
    else:
        v9 = v8 - 48
    v10 = b[2 * v4+1]
    v11 = 16 * v9
    if ( v10 < 48 | v10 > 57 ):
        v12 = v10 - 87
    else:
        v12 = v10 - 48

    for v5 in range(0,127):
        v6 = (v5 >> 4) % 16
        v7 = (16 * v5 >> 4) % 16
        if a[16 * v6 + v7] == (v11 + v12) ^ 0x19:
            flag += chr(v5)
            break

    v4 += 1

    if (v4 >= 35):
        break

print(flag)

```

运行得到flag

```

flag{This_1s_Simple_Replac3_Enc0d3}
>>>

```

题目2: HighwayHash64

题目链接:

解题过程:

拿到题运行之，提示要我们输入flag，随便输入之后程序自动退出，放入IDA中分析一下

```
Please enter flag(Note:hxb2018{digital}):
```

找到main函数，这两处有明显的比较，结合题目意思，这里应该是hash摘要计算，我们需要爆破hash的值

```
16  memset(Dst, 0, 0x104ui64);
17  sub_140001880((__int64)"Please enter flag(Note:hxb2018{digital}:", v3, v4, v5);
18  gets_s(Dst, 0x104ui64);
19  v6 = -1i64;
20  len = -1i64;
21  do
22      ++len;
23  while ( Dst[len] );
24  v13 = len;
25  if ( sub_1400017A0((__int64)&v13, 4ui64) != 0xD31580A28DD8E6C4i64 )
26      exit(1);
27  v8 = (unsigned int)(v13 - 1);
28  if ( (unsigned int)v8 >= 0x104 )
29      {
30      _report_rangecheckfailure();
31      JUMPOUT*( _QWORD *)&byte_1400019E1;
32      }
33  Dst[v8] = 0;
34  do
35      ++v6;
36  while ( v15[v6] );
37  if ( sub_1400017A0((__int64)v15, (unsigned int)v6) != 0xE3BE26AF8730545Ai64 )
38      exit(1);
39  sub_140001880((__int64)"successful!\nplease entry any key exit...", v9, v10, v11);
40  fgetchar();
41  return 0;
```

判断flag长度

判断flag内容

<https://blog.csdn.net/CharlesGodX>

我们逐个分析，可以分析得出第一个函数负责计算flag长度，第二个函数负责计算flag括号中的内容，既然是自定义的hash函数，我们可以考虑新写一个exe程序调用这个程序的函数，既然是调用那我们肯定要将exe文件修改为dll文件，工具使用010Editor，修改的方法主要是：

- IMAGE_FILE_HEADER->Characteristics(文件属性)->2102h (DLL文件一般是2102h)
- IMAGE_OPTIONAL_HEADER->AddressOfEntryPoint (程序执行入口RVA) ->0000h

我们先通过MS-DOS头部找到PE文件头，然后在通过偏移找到上面的两处位置，如下图：

0000h:	4D 5A 90 00	03 00 00 00	04 00 00 00	FF FF 00 00	MZ.....ÿÿ..
0010h:	B8 00 00 00	00 00 00 00	40 00 00 00	00 00 00 00	,.....@.....
0020h:	00 00 00 00	00 00 00 00	00 00 00 00	00 00 00 00
0030h:	00 00 00 00	00 00 00 00	00 00 00 00	08 01 00 00
0040h:	0E 1F BA 0E	00 B4 09 CD	21 B8 01 4C	CD 21 54 68	..°..'í!,.Lí!Th
0050h:	69 73 20 70	72 6F 67 72	61 6D 20 63	61 6E 6E 6F	is program cannot
0060h:	74 20 62 65	20 72 75 6E	20 69 6E 20	44 4F 53 20	be run in DOS
0070h:	6D 6F 64 65	2E 0D 0D 0A	24 00 00 00	00 00 00 00	mode....\$.....
0080h:	1F FB C5 75	5B 9A AB 26	5B 9A AB 26	5B 9A AB 26	.úÁu[š«&[š«&[š«&
0090h:	EF 06 5A 26	5F 9A AB 26	EF 06 58 26	26 9A AB 26	ï.Z&_š«&ï.X&&š«&
00A0h:	EF 06 59 26	56 9A AB 26	52 E2 38 26	58 9A AB 26	ï.Y&Vš«&Râ&Xš«&
00B0h:	5B 9A AA 26	0C 9A AB 26	60 C4 A8 27	5C 9A AB 26	[š^&.š«&`Ä''\š«&
00C0h:	60 C4 AE 27	47 9A AB 26	60 C4 AF 27	49 9A AB 26	`Ä@'Gš«&`Ä' Iš«&
00D0h:	CC C4 A3 27	58 9A AB 26	C9 C4 54 26	5A 9A AB 26	ÏÄf'Xš«&ÉÄT&Zš«&
00E0h:	CC C4 A9 27	5A 9A AB 26	52 59 63 68	5B 9A AB 26	ÏÄ@'Zš«&Rich[š«&
00F0h:	00 00 00 00	00 00 00 00	00 00 00 00	00 00 00 00
0100h:	00 00 00 00	00 00 00 00	50 45 00 00	64 86 07 00PE..d†..
0110h:	2B 92 ED 5B	00 00 00 00	00 00 00 00	F0 00 02 21	+ 'í[.....ö..!

第一处修改:

00F0h:	00 00 00 00	00 00 00 00	00 00 00 00	00 00 00 00
0100h:	00 00 00 00	00 00 00 00	50 45 00 00	64 86 07 00PE..d†..
0110h:	2B 92 ED 5B	00 00 00 00	00 00 00 00	F0 00 02 21	+ 'í[.....ö..!

第二处修改:

0110h:	2B 92 ED 5B	00 00 00 00	00 00 00 00	F0 00 02 21	+ 'í[.....ö..!
0120h:	0B 02 0E 00	00 14 01 00	00 D6 00 00	00 00 00 00ö.....
0130h:	00 00 00 00	00 10 00 00	00 00 00 40	01 00 00 00@.....

修改之后我们将后缀改为.dll然后就可以实现链接了, 根据汇编我们可以看出是__fastcall的调用约定, 因为是用rcx和rdx传的参数, 编译的时候也需要用64位编译, 因为我们调用的dll是64位

```

• .text:000000014000193D BA 04 00 00 00 mov     edx, 4
• .text:0000000140001942 89 44 24 20 mov     [rsp+158h+var_138], eax
• .text:0000000140001946 48 8D 4C 24 20 lea    rcx, [rsp+158h+var_138]
• .text:000000014000194B E8 50 FE FF FF call   sub_1400017A0

```

然后我们就可以开始实现链接了:

```

typedef __int64(__fastcall *f)(__int64 buff, unsigned __int64 len);

int main()
{
    HINSTANCE hdll;
    hdll = LoadLibrary(TEXT("F:\\reverse.dll"));
    if (hdll == NULL)
    {
        printf("Load dll Error: %d\n", GetLastError());
        return 0;
    }
    printf("Dll base is %llx\n", hdll);

    f func = ((f)((char*)hdll + 0x17A0));
}

```

第一处是需要我们爆破长度，长度最后减去hxb2018{}这9位：

```
do
    ++len;
while ( Dst[len] );
v13 = len;
if ( sub_1400017A0((__int64)&v13, 4ui64) != 0xD31580A28DD8E6C4i64 )
    exit(1);
```

代码如下：

```
int i;
unsigned long long result;
for (i = 0; i < 50; i++)
{
    result = func((long long)&i, 4);
    if (result == 0xD31580A28DD8E6C4)
    {
        printf("Len is %d\n", i - 9);
    }
}
```

运行可以看到长度为10

```
Dll base is 7fffa6e20000
Len is 10
```

第二次是需要我们爆破内容，通过sprintf快速制作10个字节的十进制数，然后穷举：

```
unsigned long long j;
unsigned long long result2;
char buff[20];
for (j = 0; j < 1000000000; j++)
{
    sprintf_s(buff, "%0.10llu", j);
    if (j % 100000 == 0)
    {
        printf("%0.10llu\n", j);
    }
    result2 = func((long long)buff, 10);
    if (result2 == 0xE3BE26AF8730545A)
    {
        printf("flag is %lld\n", j);
        return 0;
    }
}
```

最后爆破一段时间得到flag


```
1530900000
1531000000
1531100000
1531200000
1531300000
1531400000
1531500000
1531600000
1531700000
1531800000
1531900000
1532000000
1532100000
1532200000
1532300000
1532400000
1532500000
1532600000
flag is 1532649708
```

最后贴一个总代码

```

#include "stdafx.h"
#include<Windows.h>

typedef __int64(__fastcall *f)(__int64 buff, unsigned __int64 len);

int main()
{
    HINSTANCE hdll;
    hdll = LoadLibrary(TEXT("F:\\reverse.dll"));
    if (hdll == NULL)
    {
        printf("Load dll Error: %d\n", GetLastError());
        return 0;
    }
    printf("Dll base is %llx\n", hdll);

    f func = ((f)((char*)hdll + 0x17A0));

    int i;
    unsigned long long result;
    for (i = 0; i<50; i++)
    {
        result = func((long long)&i, 4);
        if (result == 0xD31580A28DD8E6C4)
        {
            printf("Len is %d\n", i - 9);
        }
    }

    unsigned long long j;
    unsigned long long result2;
    char buff[20];
    for (j = 0; j < 1000000000; j++)
    {
        sprintf_s(buff, "%0.10llu", j);
        if (j % 100000 == 0)
        {
            printf("%0.10llu\n", j);
        }
        result2 = func((long long)buff, 10);
        if (result2 == 0xE3BE26AF8730545A)
        {
            printf("flag is %lld\n", j);
            return 0;
        }
    }
    return 0;
}

```

Tips: 每个人的flag在复现的时候可能不同哦

0x02: 总结

HighwayHash64这道题可以帮助我们熟悉PE文件结构，是一道很不错的题目，还有一道题目More efficient than JS我没有记录，大概是关于WebAssembly逆向的，可以使用idawasm插件在IDA中进行分析，也可以使用wasmdc生成伪c代码分析，就是环境搭建比较麻烦。

参考的wp:

<https://www.anquanke.com/post/id/164604#h2-6>

<https://impakho.com/post/hxb-2018-writeup>